

Fall 2003

ICS 275A – Constraint Networks

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Figure 2.1: The 4-queens constraint network. The network has four variables, all with domains  $D_i = \{1, 2, 3, 4\}$ . (a) The labeled chess board. (b) The constraints between variables.



 $R_{12} = \{(1,3), (1,4), (2,4), (3,1), (4,1), (4,2)\}$   $R_{13} = \{(1,2), (1,4), (2,1), (2,3), (3,2), (3,4), (4,1), (4,3)\}$   $R_{14} = \{(1,2), (1,3), (2,1), (2,3), (2,4), (3,1), (3,2), (3,4), (4,2), (4,2), (4,3)\}$   $R_{23} = \{(1,3), (1,4), (2,4), (3,1), (4,1), (4,2)\}$   $R_{24} = \{(1,2), (1,4), (2,1), (2,3), (3,2), (3,4), (4,1), (4,3)\}$   $R_{34} = \{(1,3), (1,4), (2,4), (3,1), (4,1), (4,2)\}$ 

(a)

(b)

Figure 2.2: Not all consistent instantiations are part of a solution: (a) A consistent instantiation that is not part of a solution. (b) The placement of the queens corresponding to the solution (2, 4, 1, 3). (c) The placement of the queens corresponding to the solution (3, 1, 4, 2).







Figure 2.3: Constraint graphs of (a) the crossword puzzle and (b) the 4-queens problem.



Figure 2.4: The constraint graph and constraint relations of the scheduling problem example.





#### Figure 2.7: Scene labeling constraint network

$$R_{21} = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \end{bmatrix} R_{31} = \begin{bmatrix} 0 & 1 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \end{bmatrix} R_{51} = \begin{bmatrix} 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$$R_{24} = R_{37} = R_{56} = \begin{bmatrix} 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 \end{bmatrix}$$

$$R_{26} = R_{34} = R_{57} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 \end{bmatrix}$$
Fork: 
$$Arrow: + Arrow: +$$

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Figure 2.8: The interaction graph of theory  $\varphi = \{(\neg C), (A \lor B \lor A)\}$  $C), (\neg A \lor B \lor E), (\neg B \lor C \lor D)\}.$ В С E



**Properties of binary constraint networks:** 

Figure 2.10: (a) A graph  $\Re$  to be colored by two colors, (b) an equivalent representation  $\Re$ ' having a newly inferred constraint between  $x_1$  and  $x_3$ .



## **Relations vs networks**

- Can we represent the relations
- x1, x2, x3 = (0,0,0)(0,1,1)(1,0,1)(1,1,0)
- X1, x2, x3, x4 = (1,0,0,0)(0,1,0,0)(0,0,1,0)(0,0,0,1)

# **Relations vs networks**

- Can we represent the relations
- x1, x2, x3 = (0,0,0)(0,1,1)(1,0,1)(1,1,0)
- X1, x2, x3, x4 = (1, 0, 0, 0)(0, 1, 0, 0)(0, 0, 1, 0)(0, 0, 0, 1)
- Most relations cannot be represented by networks:
- Number of relations 2<sup>^</sup>(n<sup>k</sup>)
- Number of networks: 2<sup>((k<sup>2</sup>)(n<sup>2</sup>))</sup>

# The minimal and projection networks

- The projection network of a relation is obtained by projecting it onto each pair of its variables (yielding a binary network).
- Relation = {(1,1,2)(1,2,2)(1,2,1)}

What is the projection network?

- What is the relationship between a relation and its projection network?
- {(1,1,2)(1,2,2)(2,1,3)(2,2,2)}, solve its projection network?

## **Projection network (continued)**

 Theorem: Every relation is included in the set of solutions of its projection network.

• **Theorem**: The projection network is the tightest upper bound binary networks representation of the relation.

#### **Projection network**

**Theorem 2.3.8** For every relation  $\rho$ ,  $\rho \subseteq sol(P(\rho))$ .

**Theorem 2.3.9** The projection network  $P(\rho)$  is the tightest upper bound network representation of  $\rho$ ; there is no binary network  $\mathcal{R}'$ , s.t.  $\rho \subseteq sol(\mathcal{R}') \subset sol(P(\rho))$ .

#### **The Minimal Network** (partial order between networks)

**Definition 2.3.10** Given two binary networks,  $\mathcal{R}'$  and  $\mathcal{R}$ , on the same set of variables  $x_1, ..., x_n, \mathcal{R}'$  is at least as tight as  $\mathcal{R}$  iff for every i and  $j, \mathcal{R}'_{ij} \subseteq \mathcal{R}_{ij}$ .

**Definition 2.3.14** Let  $\{\mathcal{R}_1, ..., \mathcal{R}_l\}$  be the set of all networks equivalent to  $\mathcal{R}_0$  and let  $\rho = sol(\mathcal{R}_0)$ . Then the minimal network M of  $\mathcal{R}_0$  is defined by  $M(\mathcal{R}_0) = \bigcap_{i=1}^l \mathcal{R}_i$ .

**Theorem 2.3.15** For every binary network  $\mathcal{R}$  s.t.  $\rho = sol(\mathcal{R}), M(\rho) = P(\rho)$ .

Figure 2.11: The 4-queens constraint network: (a) The constraint graph. (b) The minimal binary constraints. (c) The minimal unary constraints (the domains).



$$M_{12} = \{(2,4), (3,1)\}$$

$$M_{13} = \{(2,1), (3,4)\}$$

$$M_{14} = \{(2,3), (3,2)\}$$

$$M_{23} = \{(1,4), (4,1)\}$$

$$M_{24} = \{(1,2), (4,3)\}$$

$$M_{34} = \{(1,3), (4,2)\}$$

$$D_{1} = \{1,3\}$$

$$D_{2} = \{1,4\}$$

$$D_{3} = \{1,4\}$$

$$D_{4} = \{1,3\}$$

(a)

(b) (c)

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### **Minimal network**

- The minimal network is perfectly explicit for binary and unary constraints:
  - Every pair of values permitted by the minimal constraint is in a solution.
- Binary-decomposable networks:
  - A network whose all projections are binary decomposable
  - The minimal network repesenst fully binarydecomposable networks.
  - Ex: (x,y,x,t) = {(a,a,a,a)(a,b,b,b,)(b,b,a,c)} is binary representable but what about its projection on x,y,z?