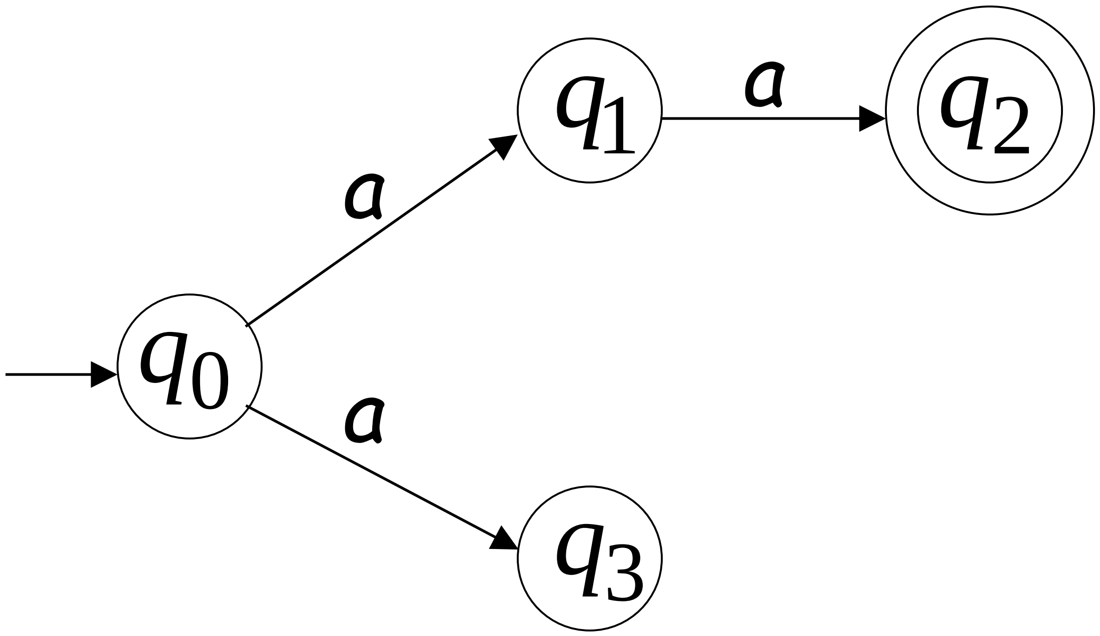
Non Deterministic Automata

Linz: Nondeterministic Finite Accepters, page 51

Nondeterministic Finite Accepter (NFA)

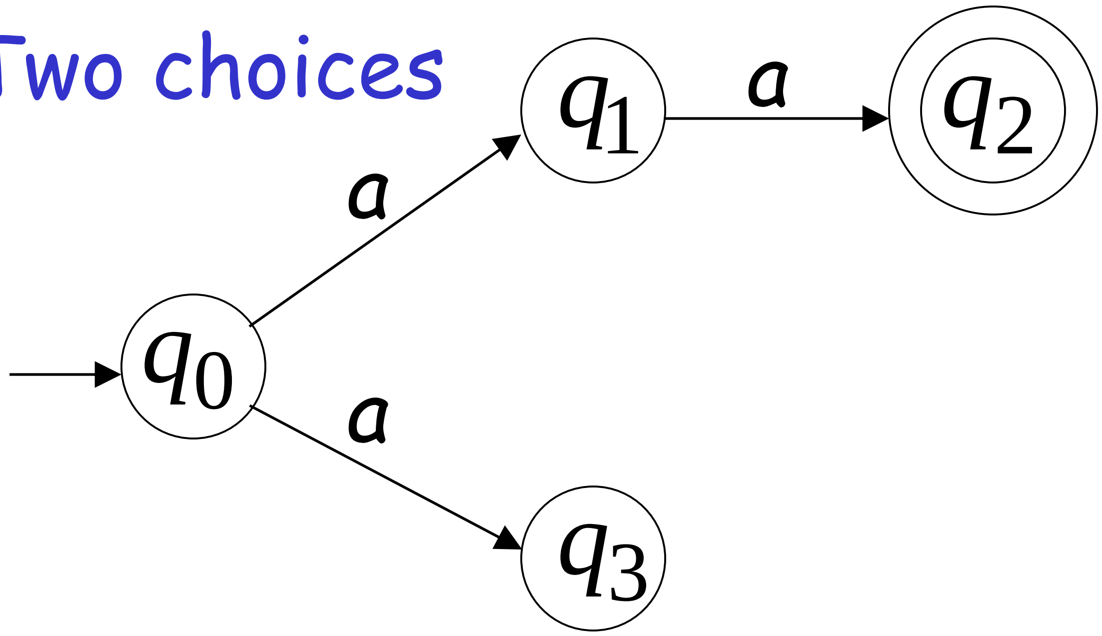
Alphabet = $\{a\}$



Nondeterministic Finite Accepter (NFA)

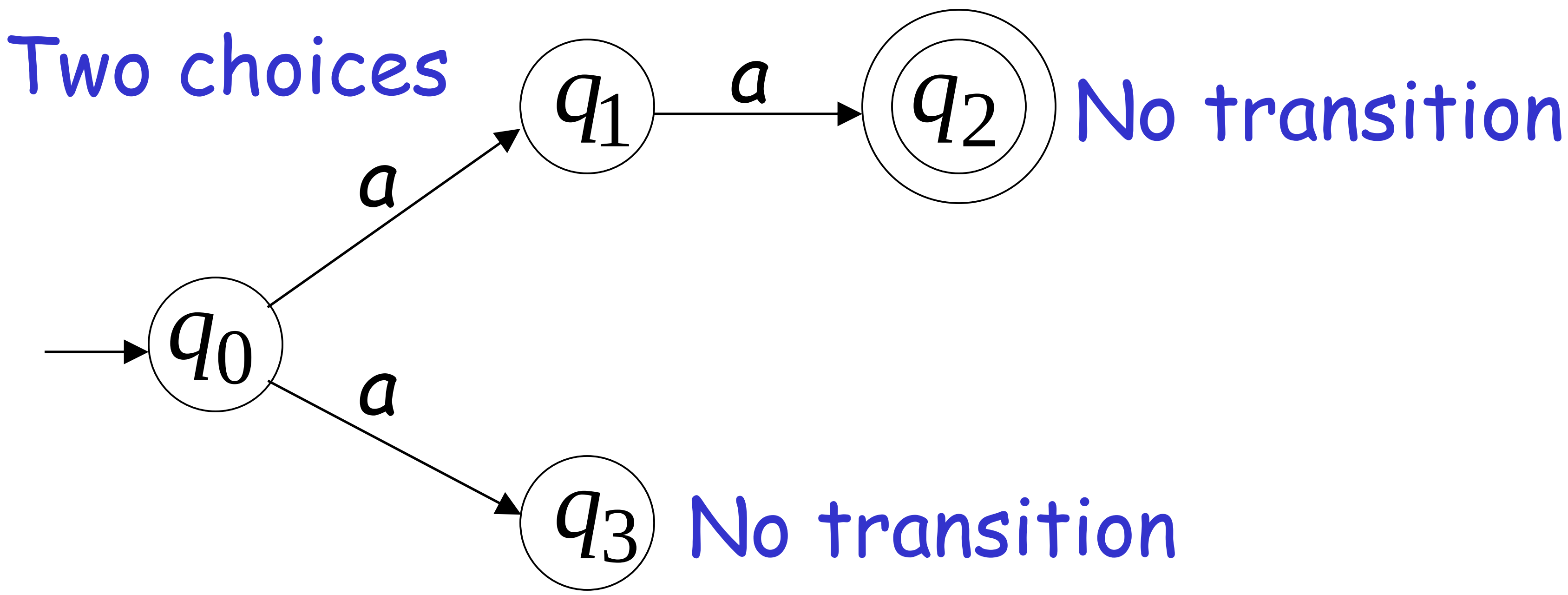
Alphabet = $\{a\}$

Two choices

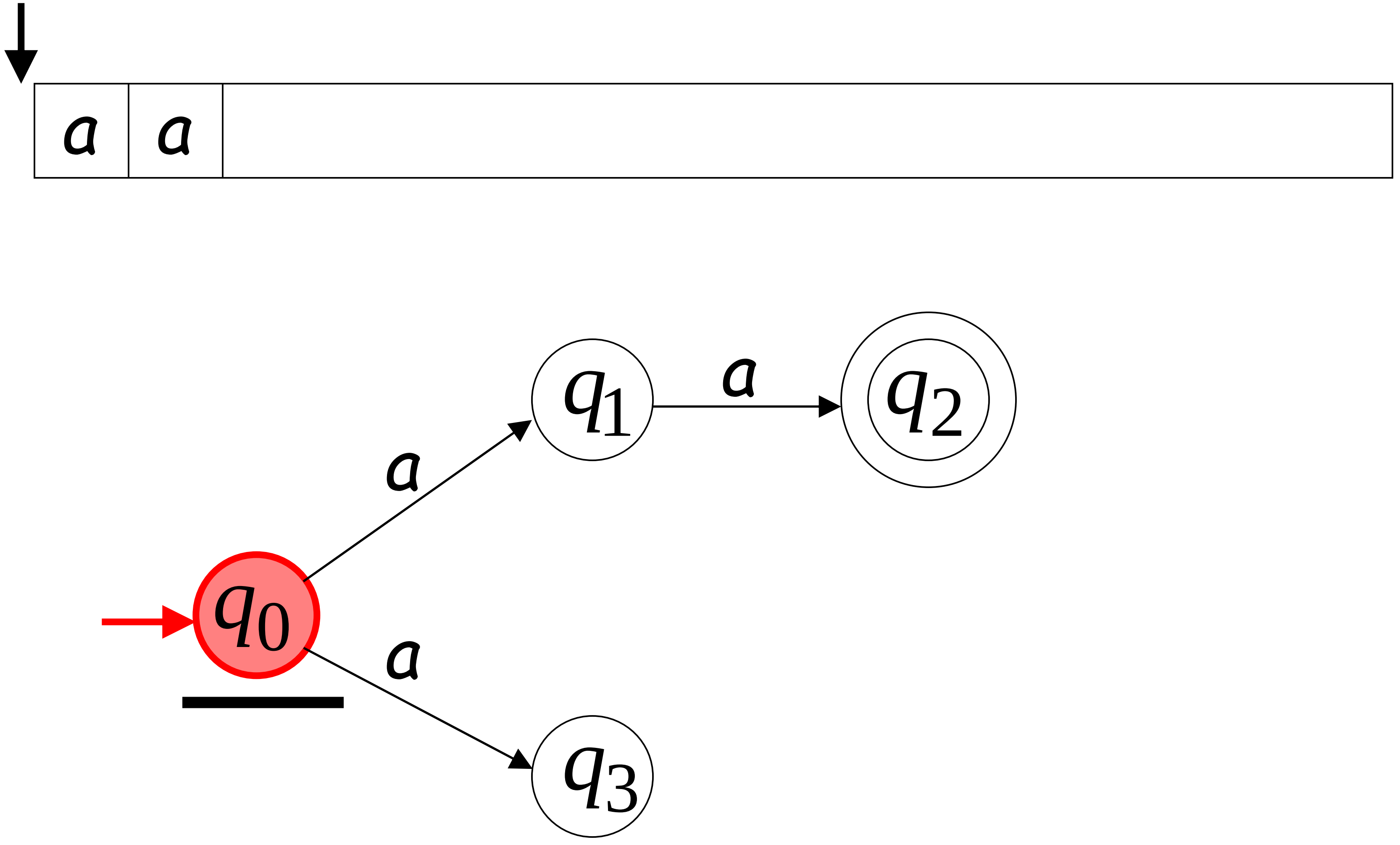


Nondeterministic Finite Accepter (NFA)

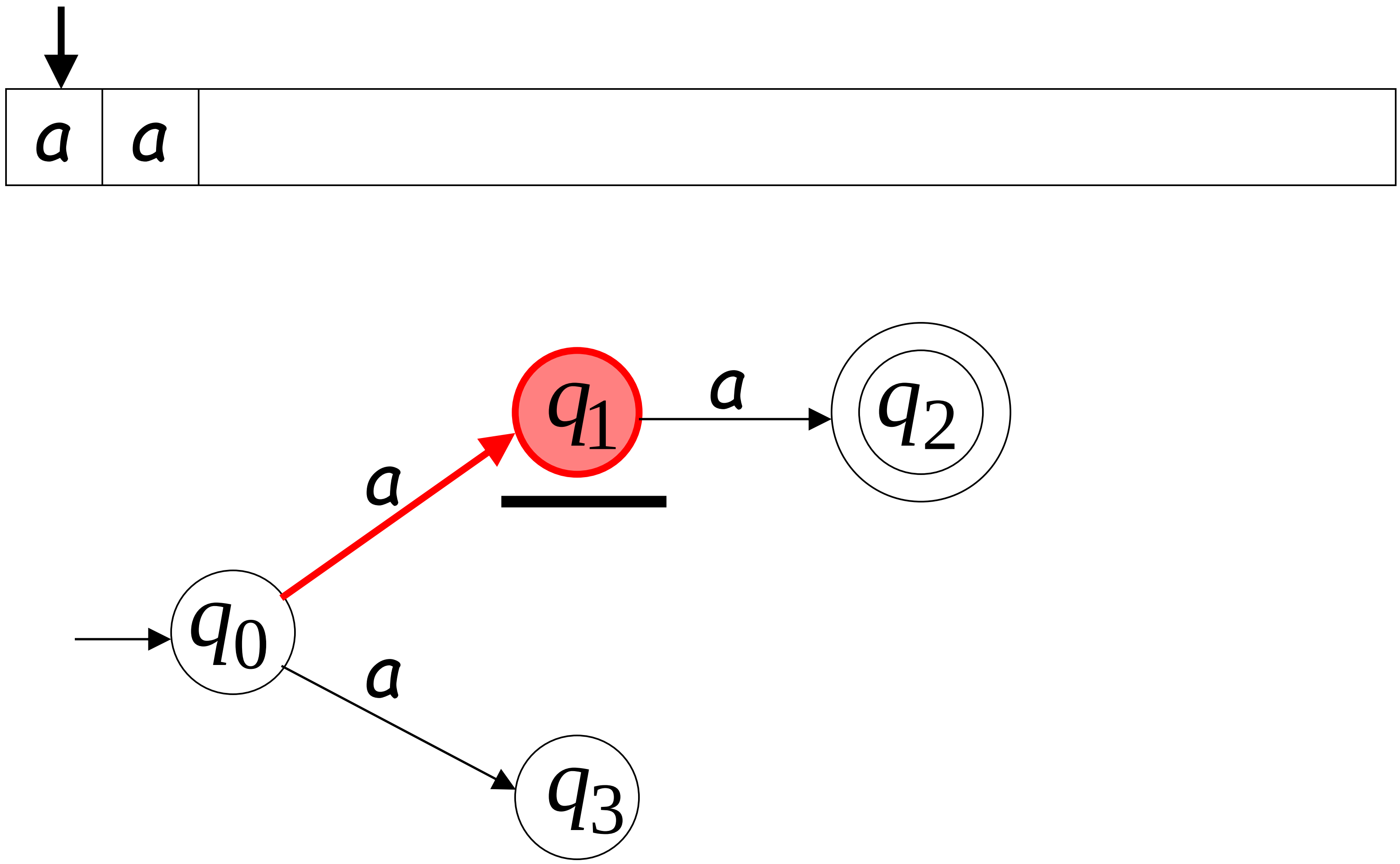
Alphabet = $\{a\}$



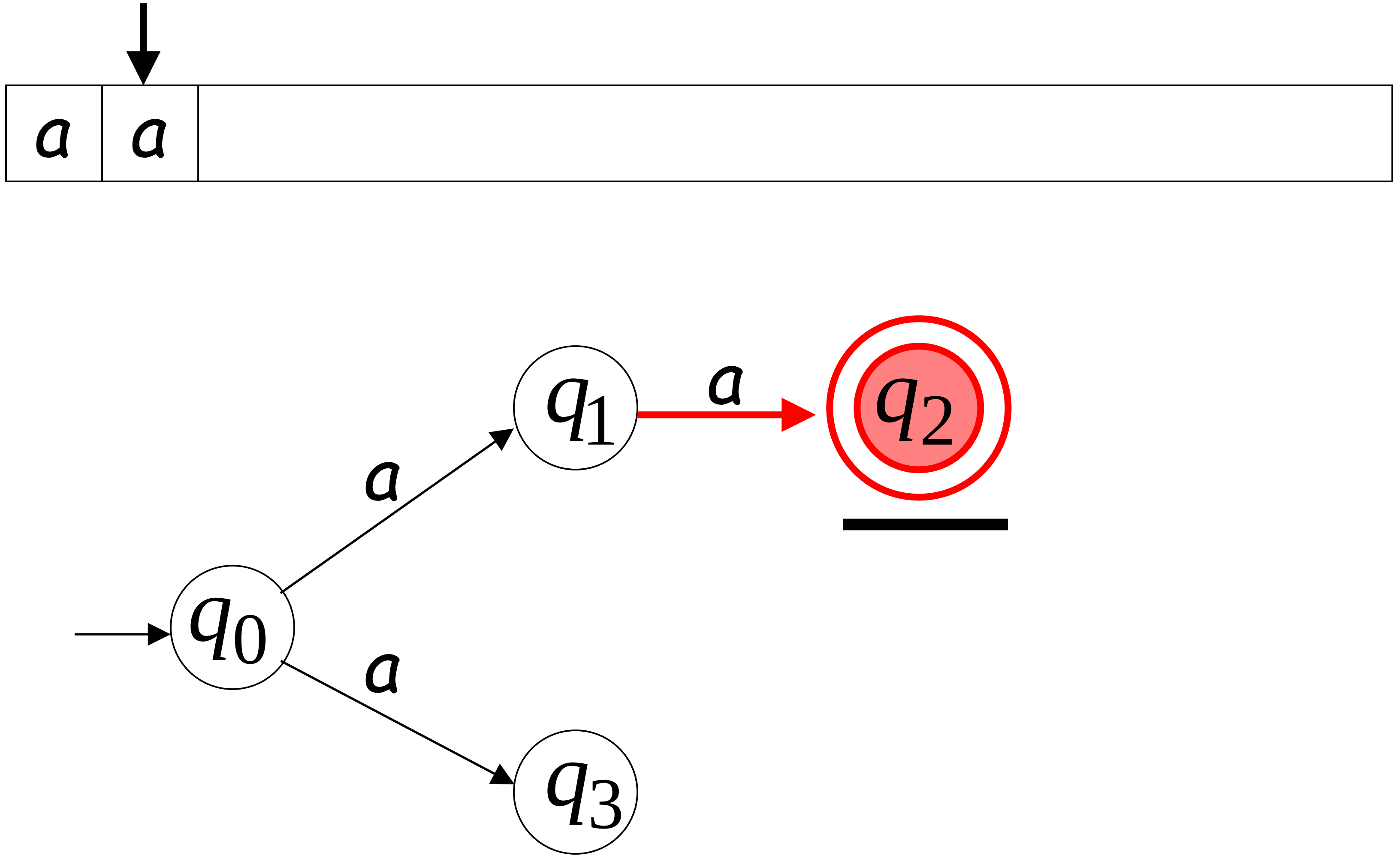
First Choice



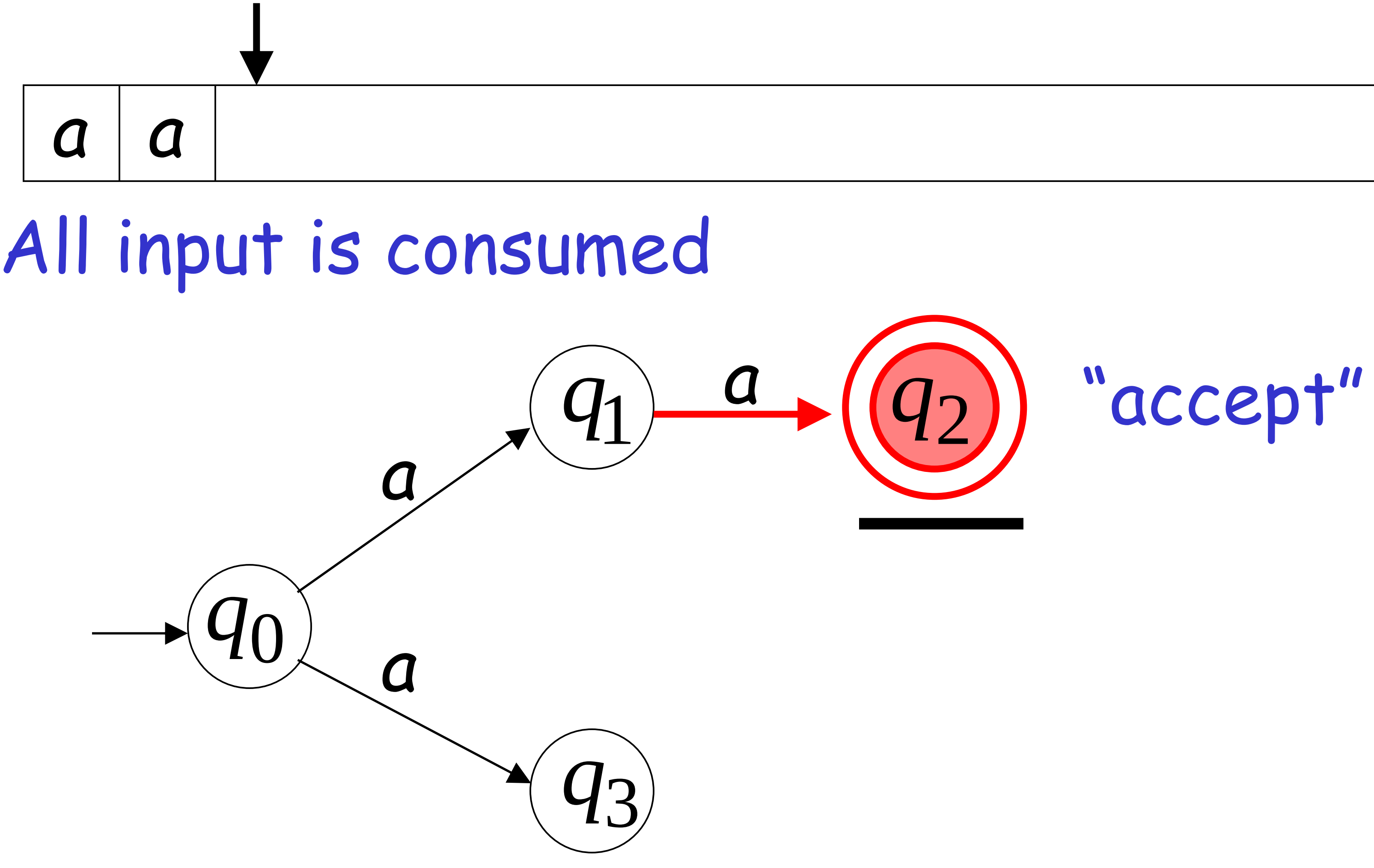
First Choice



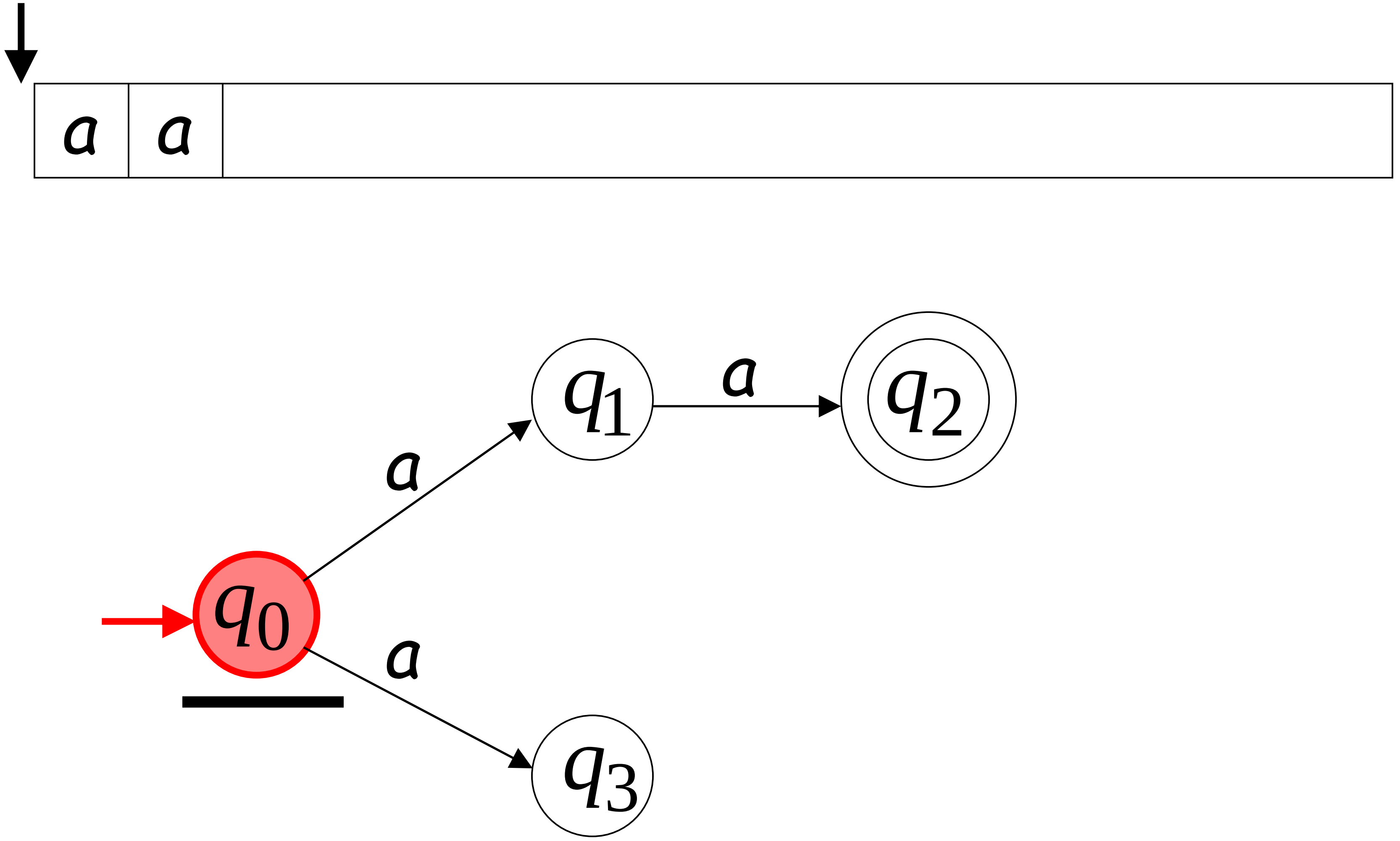
First Choice



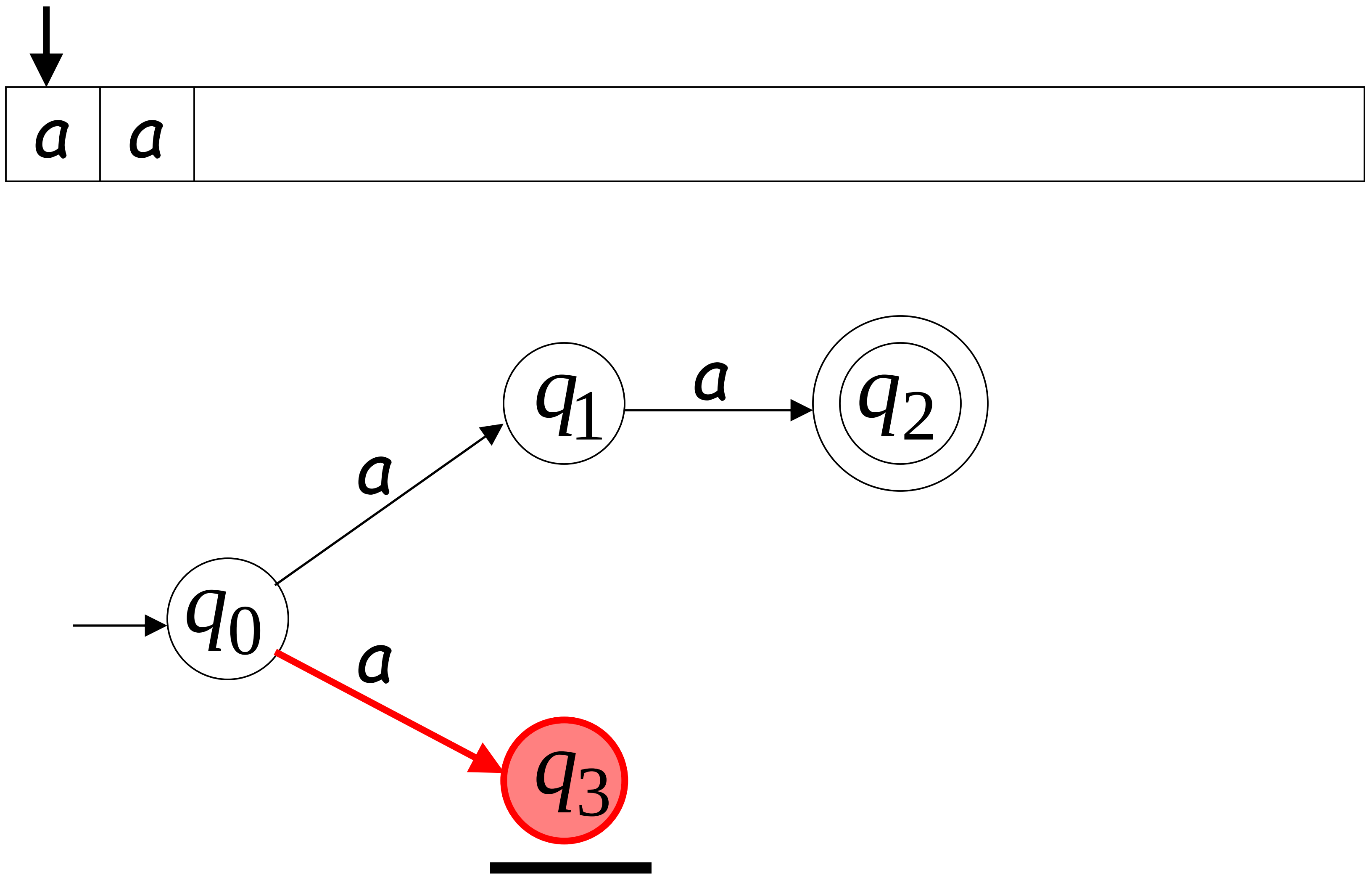
First Choice



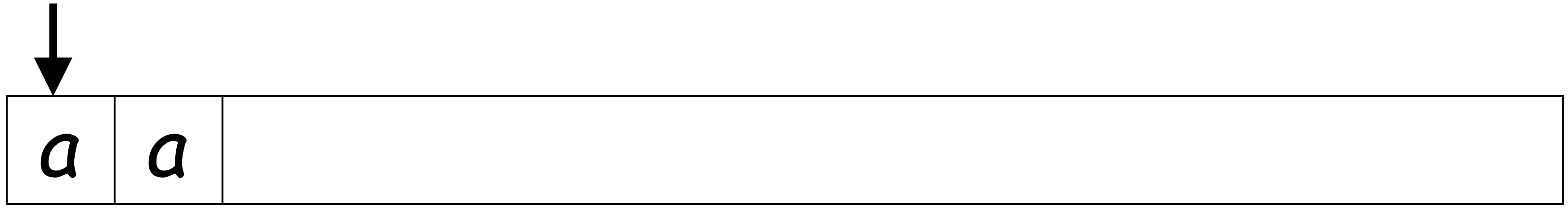
Second Choice

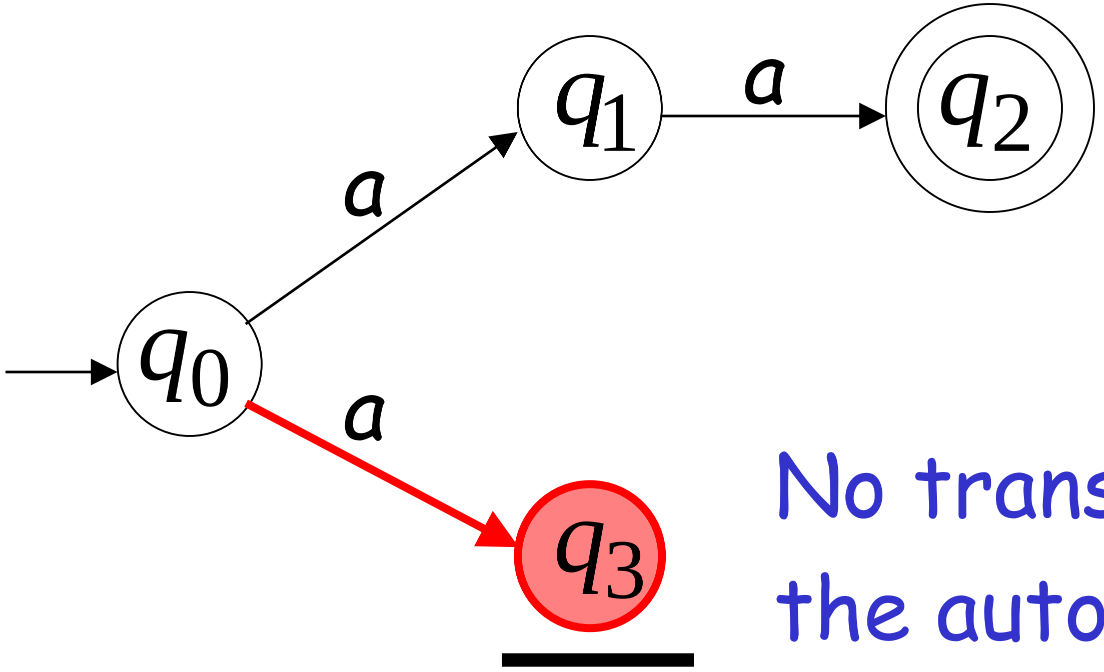


Second Choice



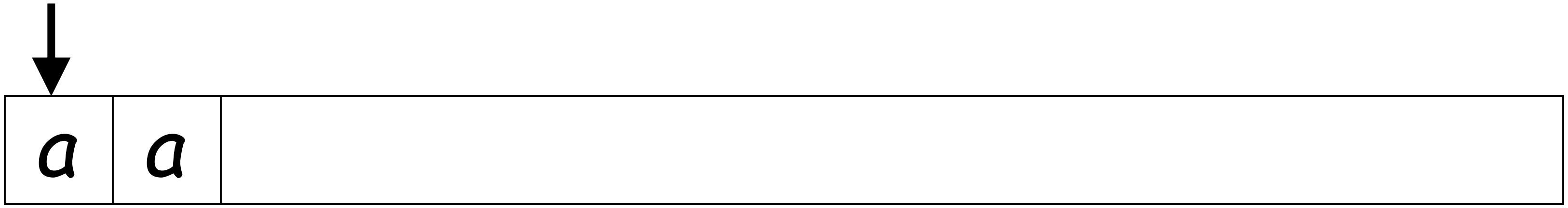
Second Choice



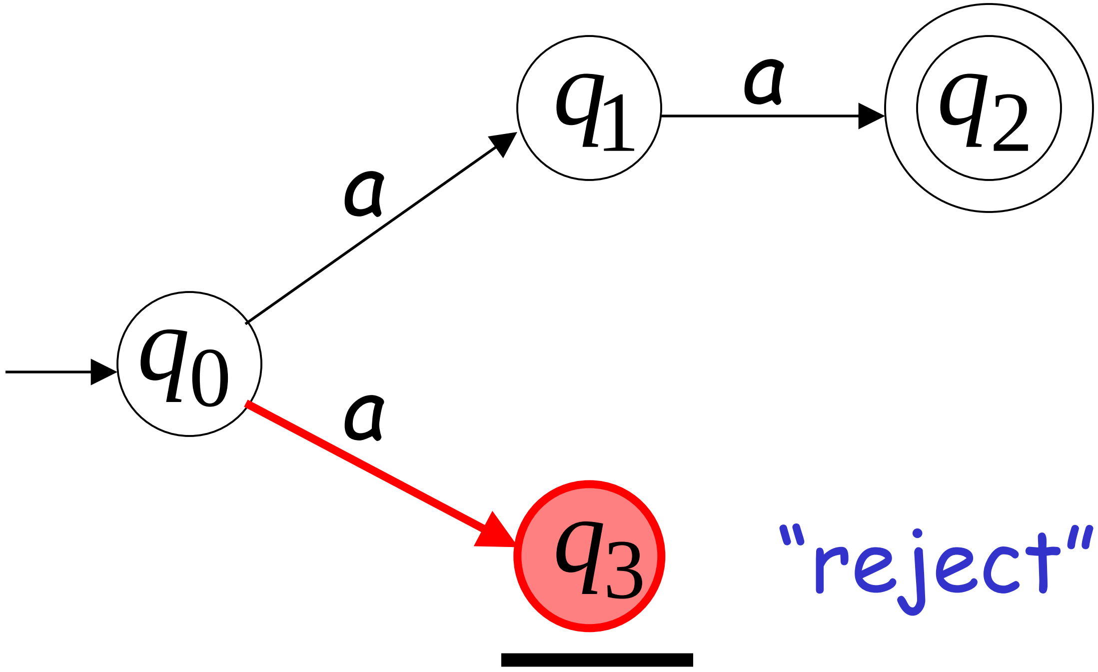


No transition:
the automaton hangs

Second Choice



Input cannot be consumed



"reject"

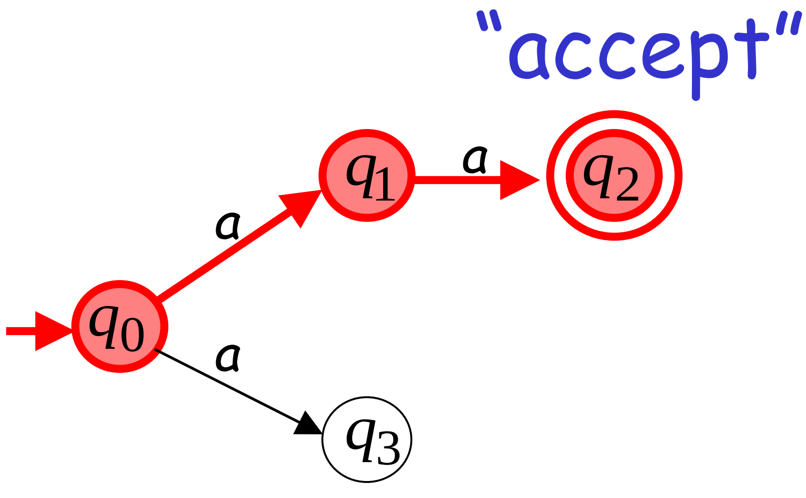
An NFA accepts a string:

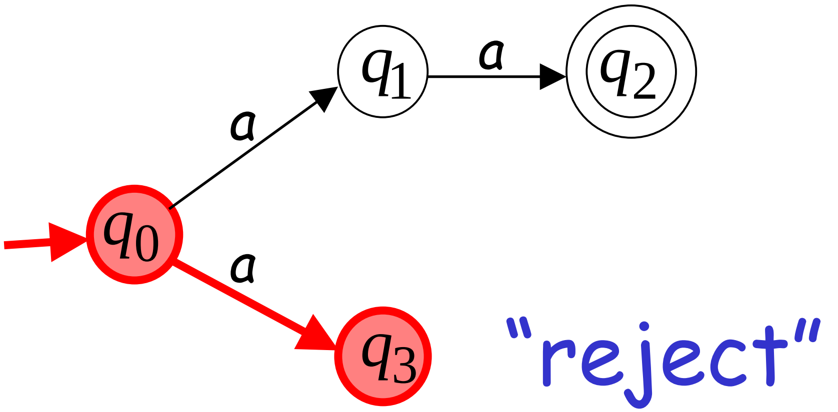
when there is a computation of the NFA that accepts the string

- All the input is consumed and the automaton is in a final state

Example

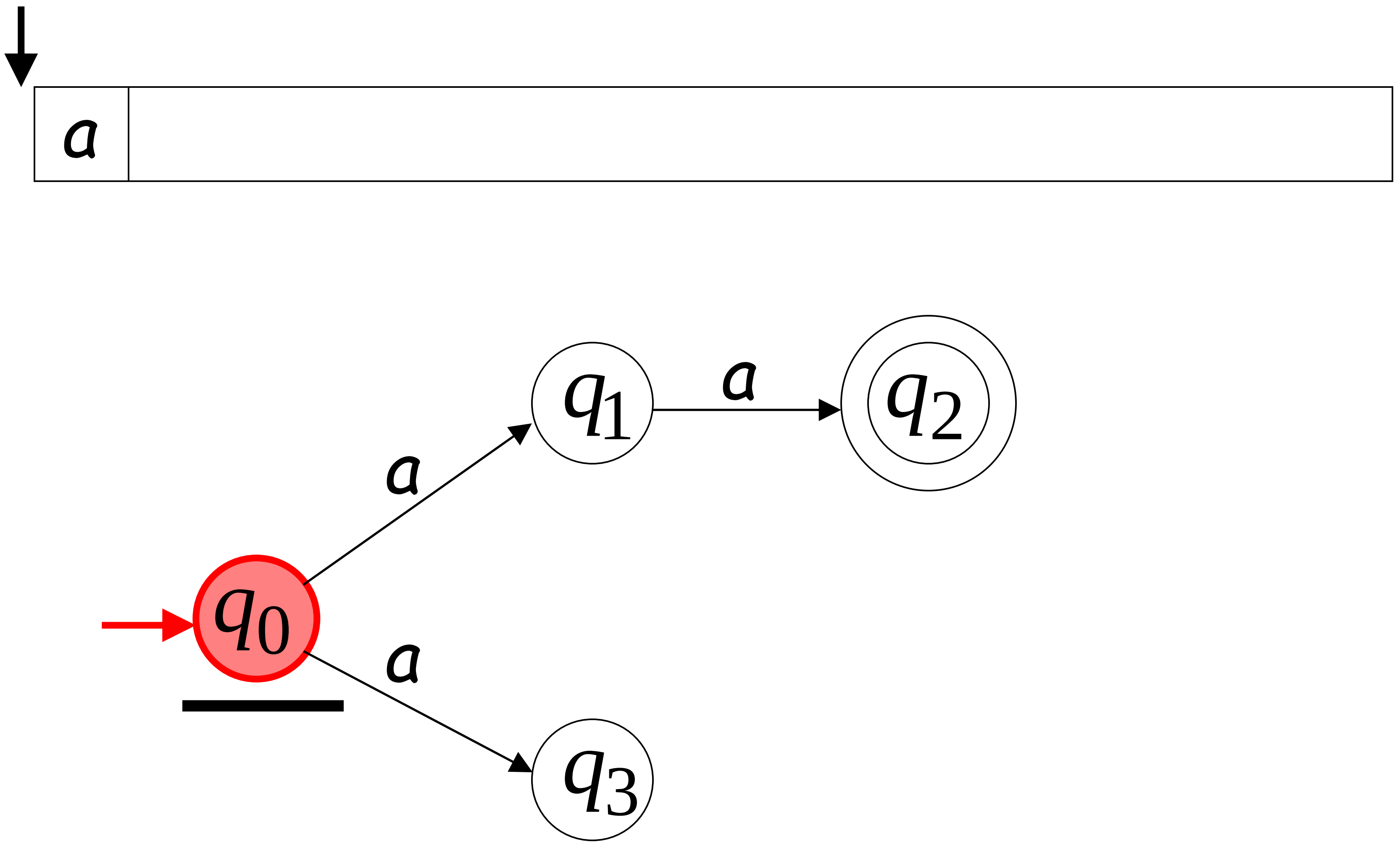
aa is accepted by the NFA:



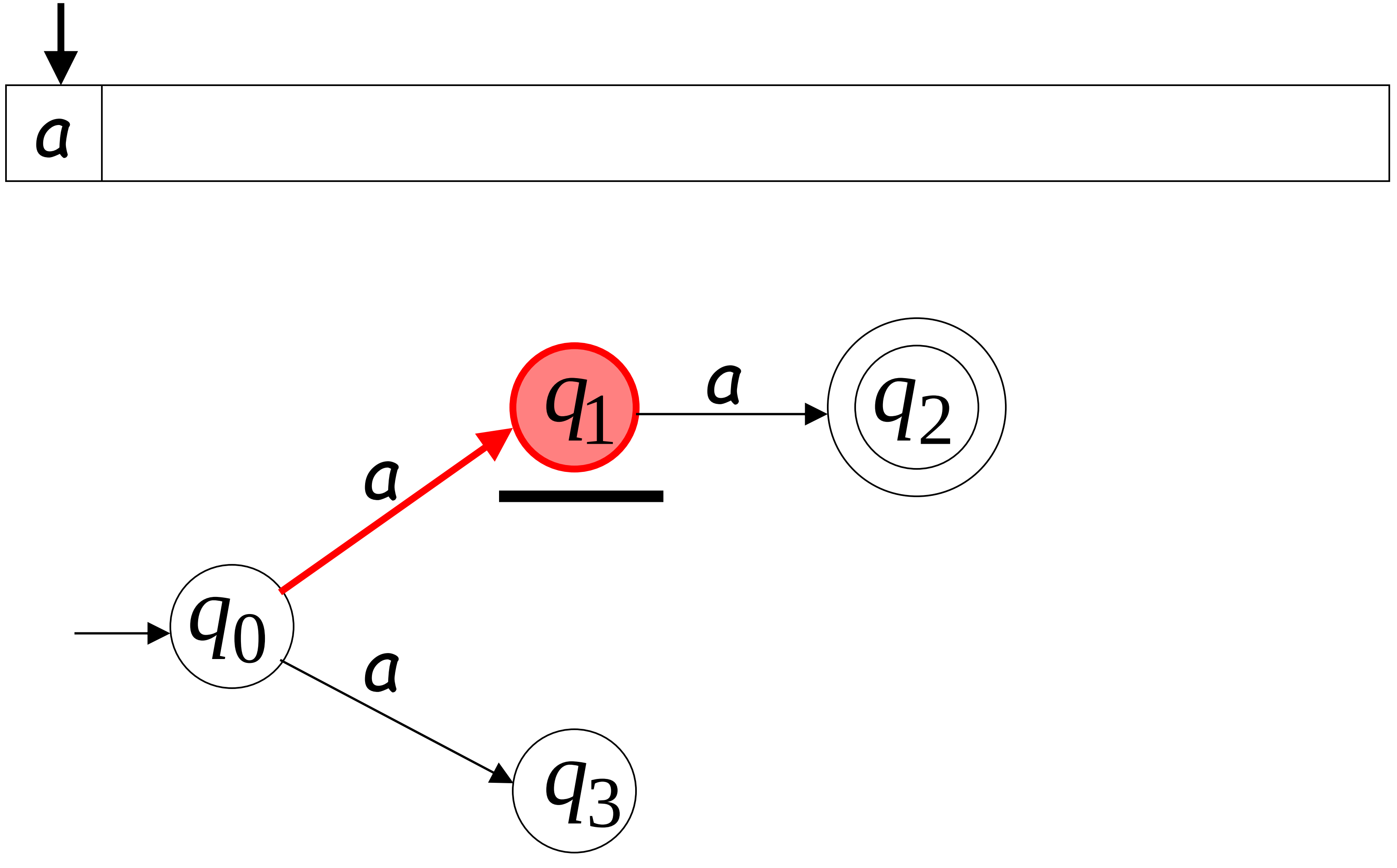


because this computation accepts aa

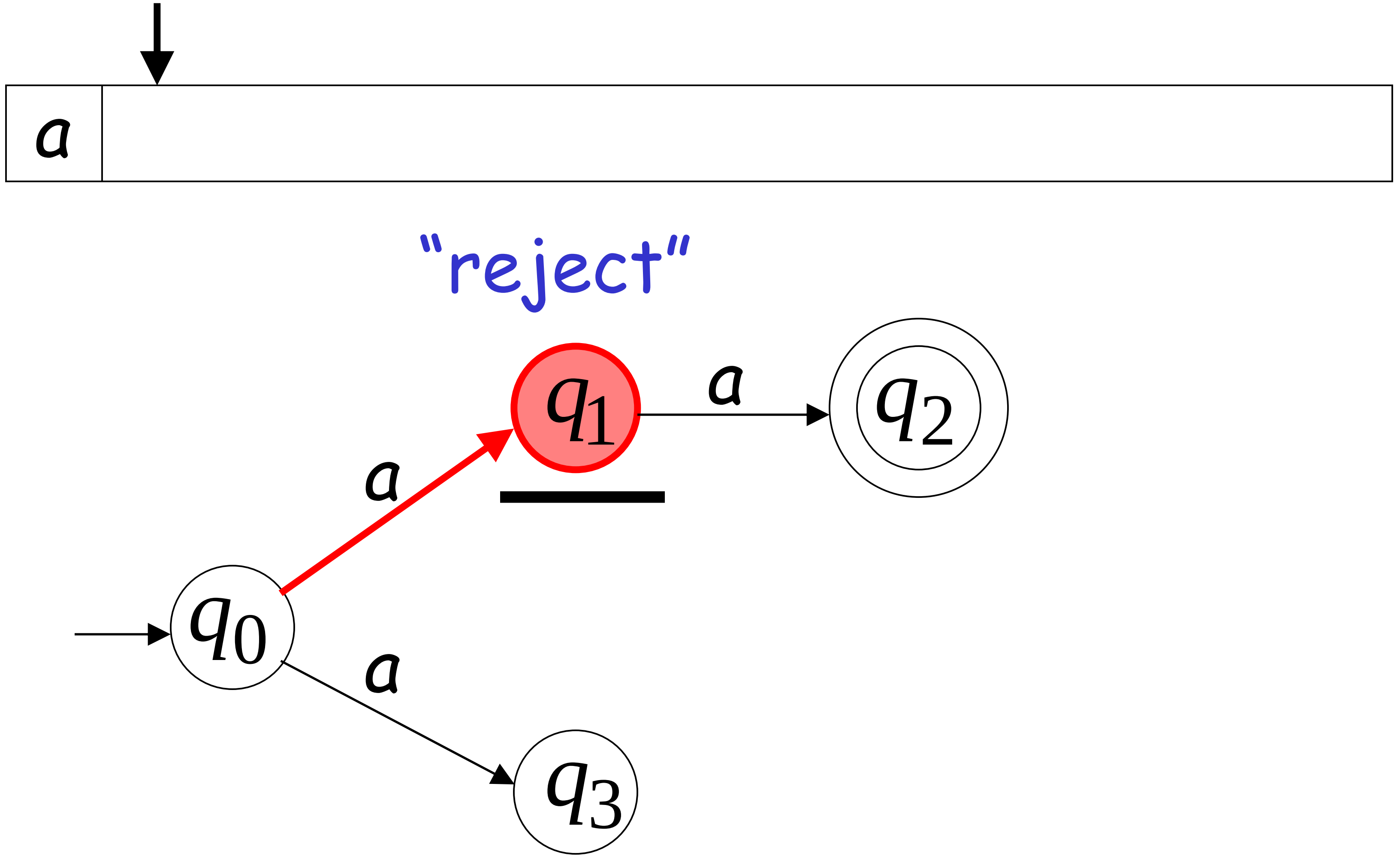
Rejection example



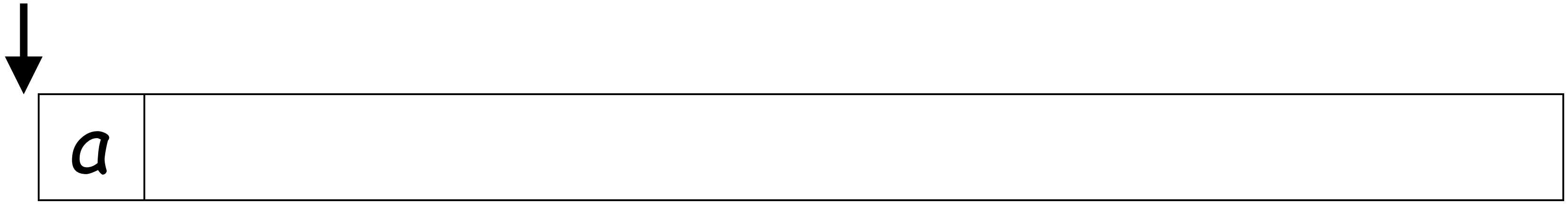
First Choice

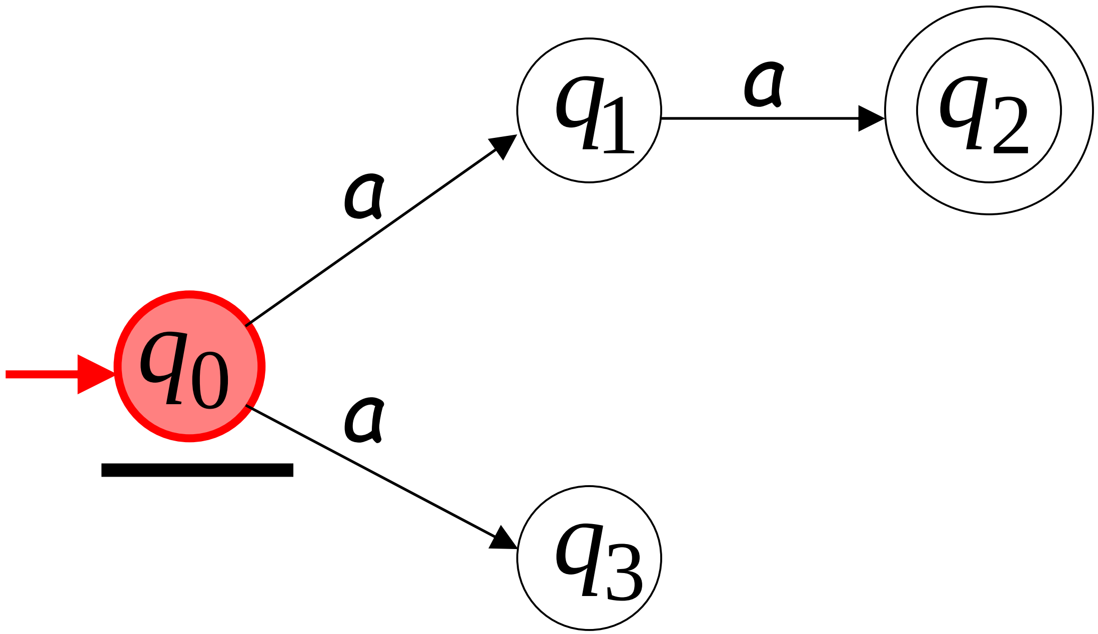


First Choice

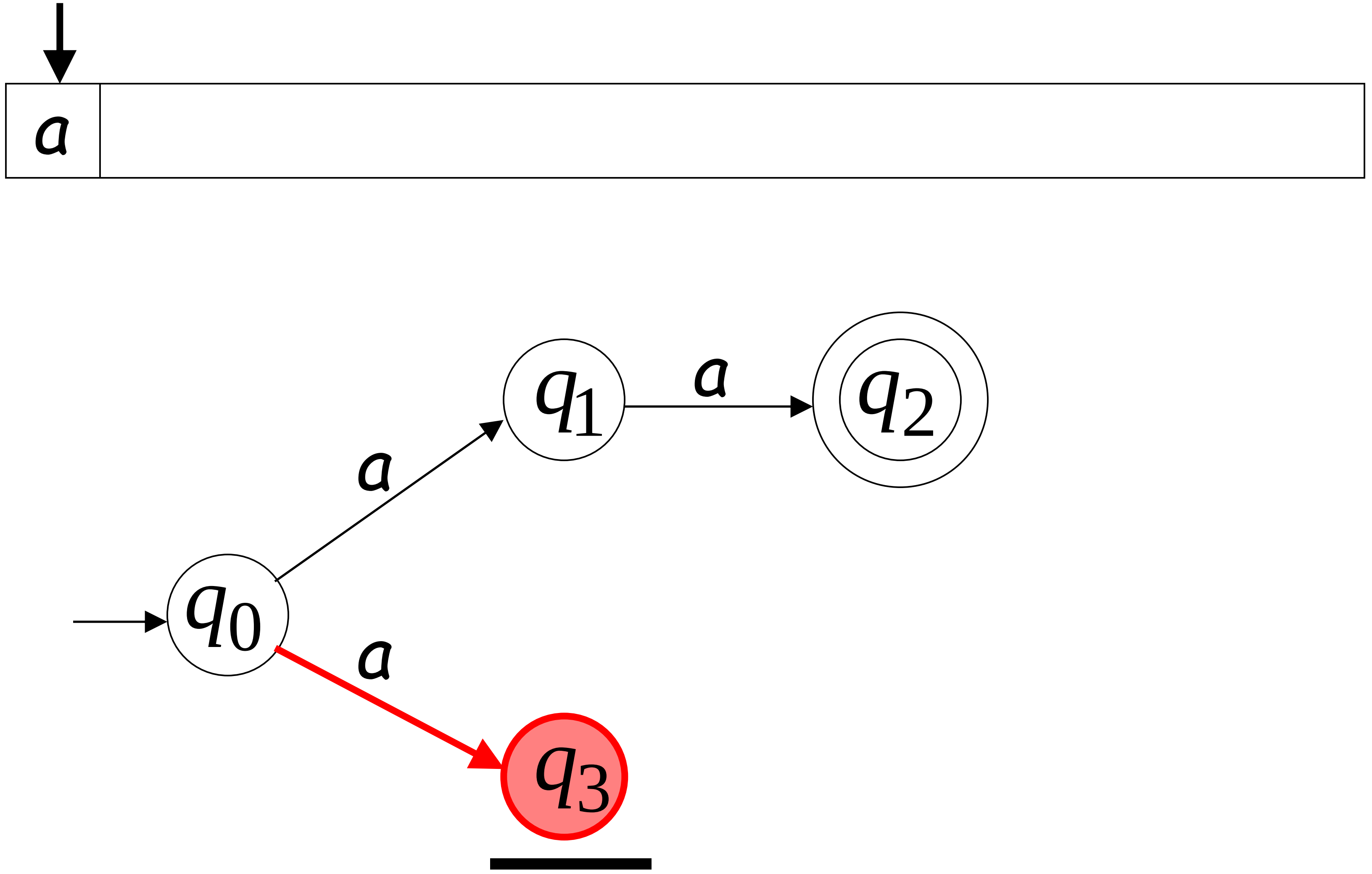


Second Choice

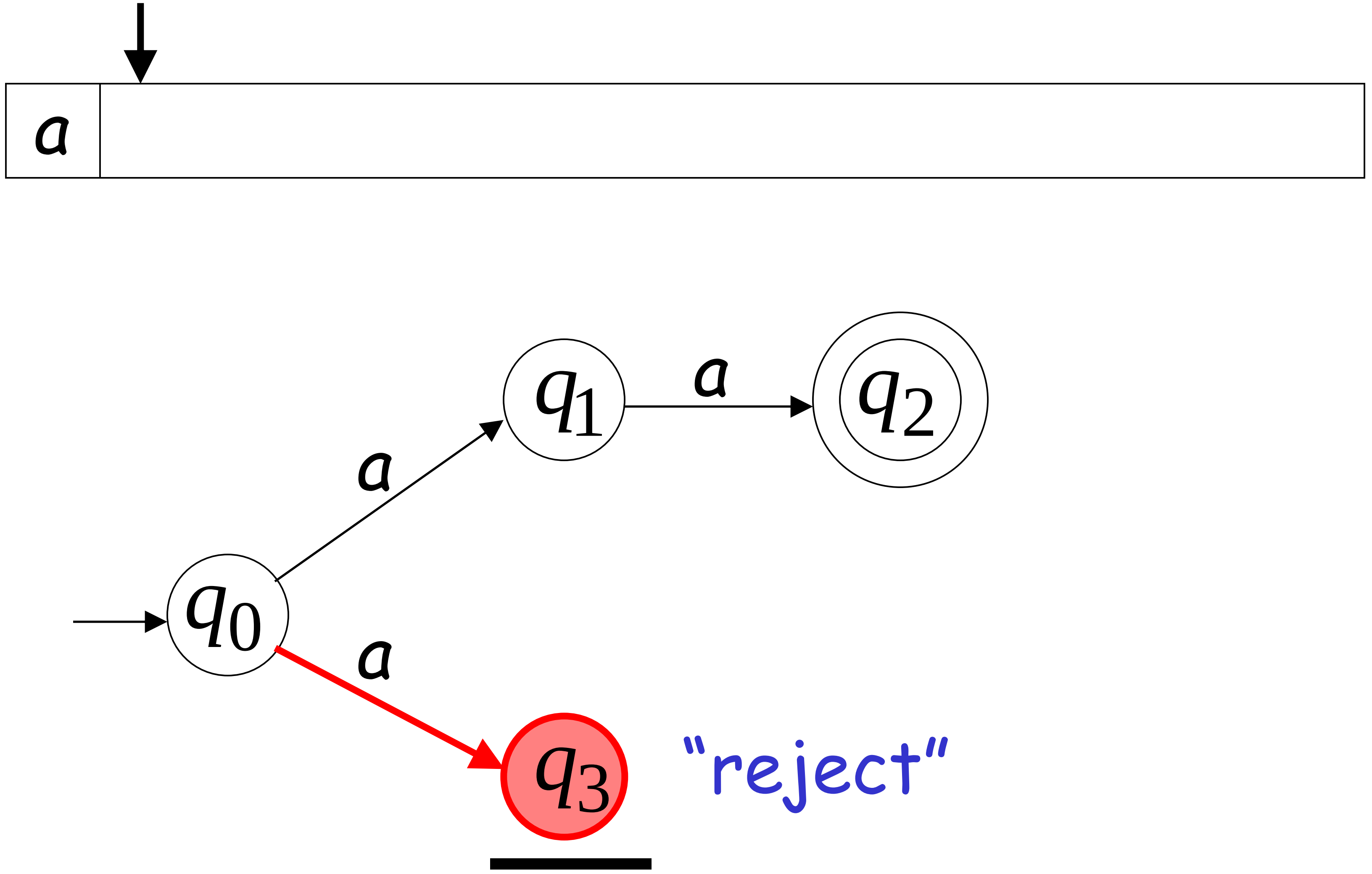




Second Choice



Second Choice



An NFA rejects a string:

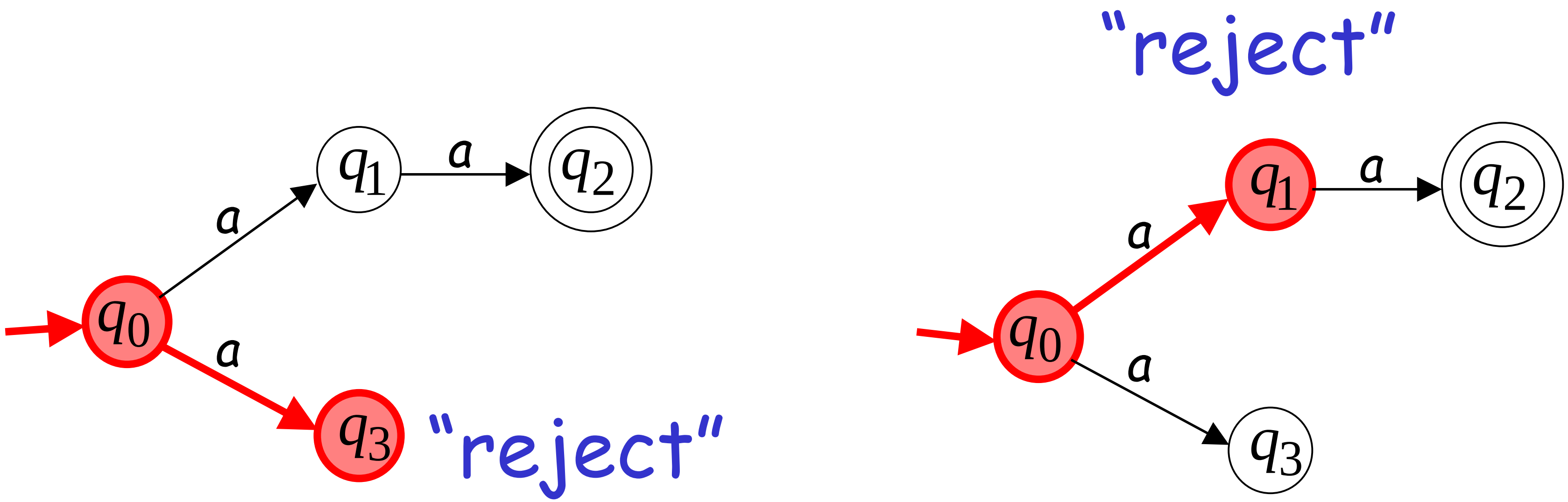
when there is no computation of the NFA that accepts the string

- All the input is consumed and the automaton is in a non final state

- The input cannot be consumed

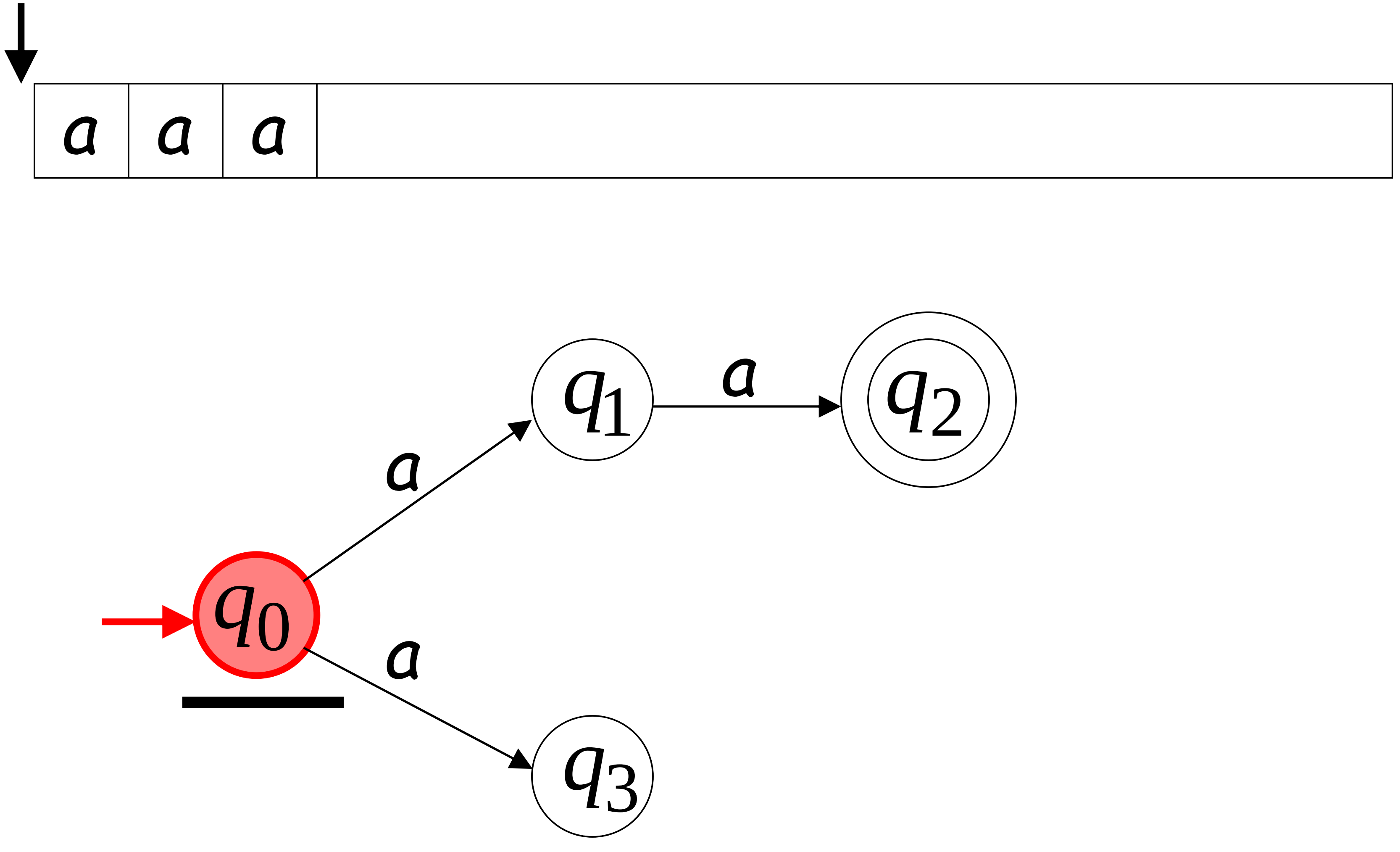
Example

a is rejected by the NFA:

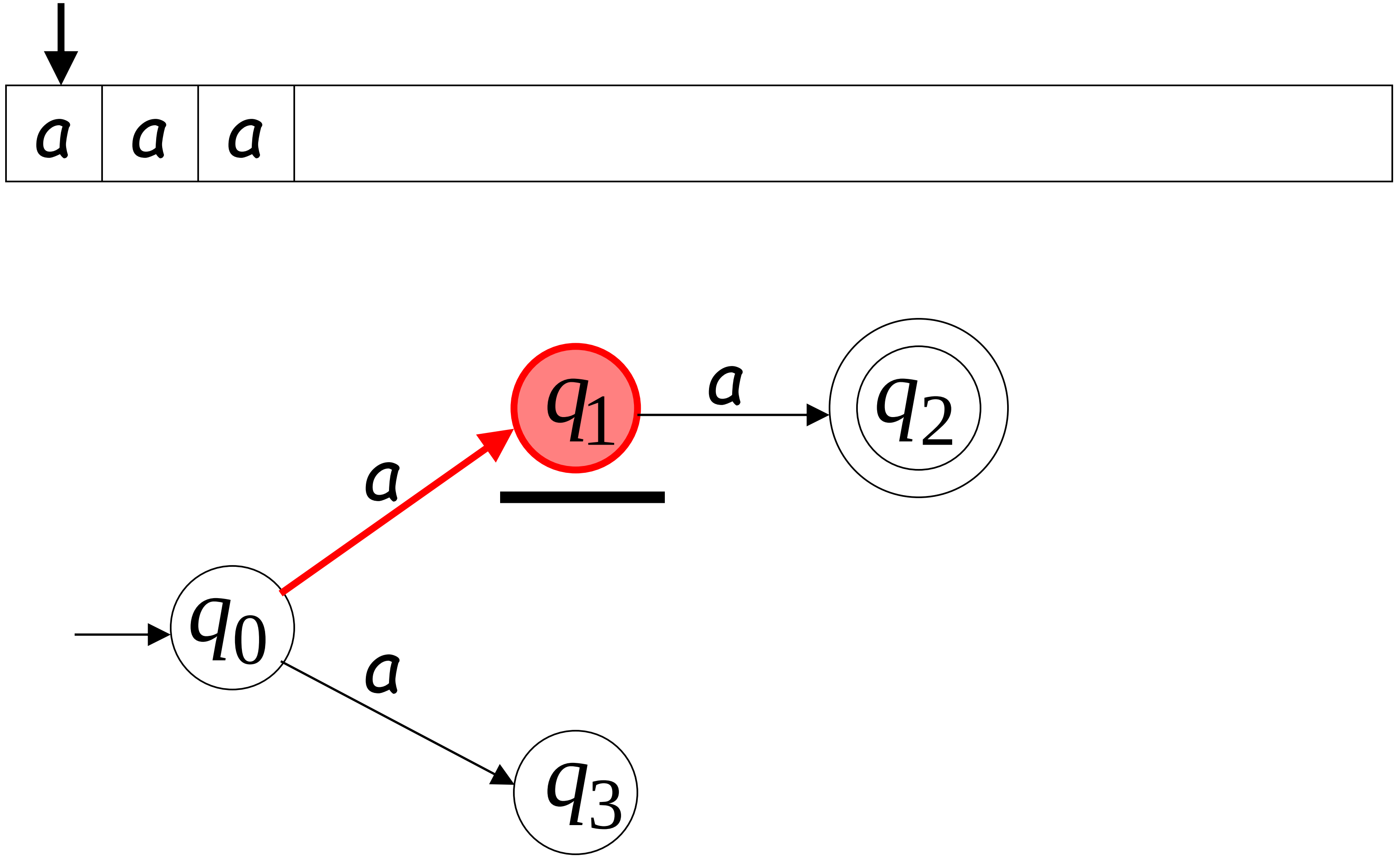


All possible computations lead to rejection

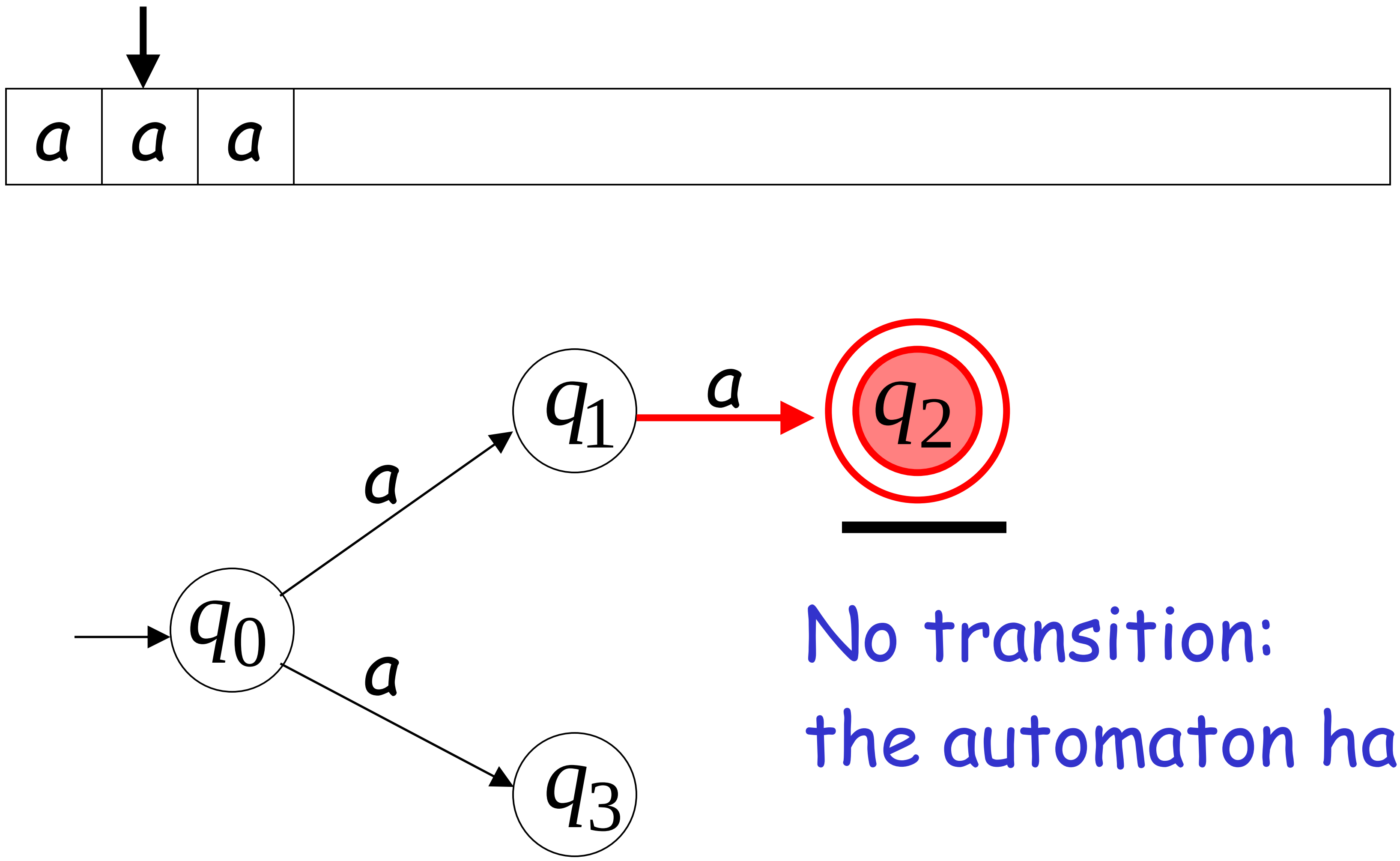
Rejection example



First Choice

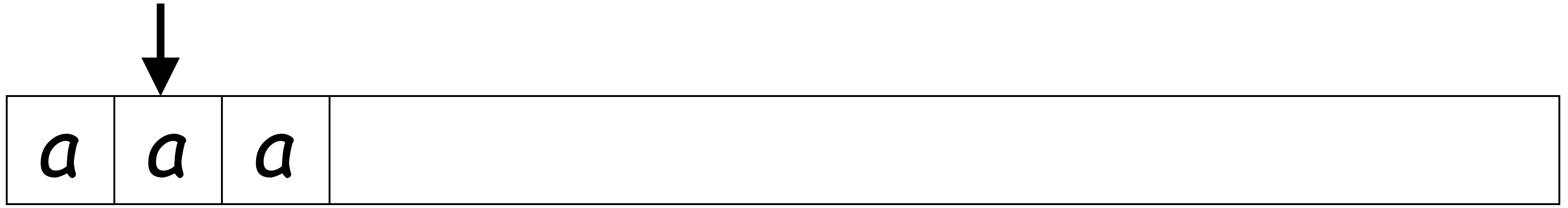


First Choice

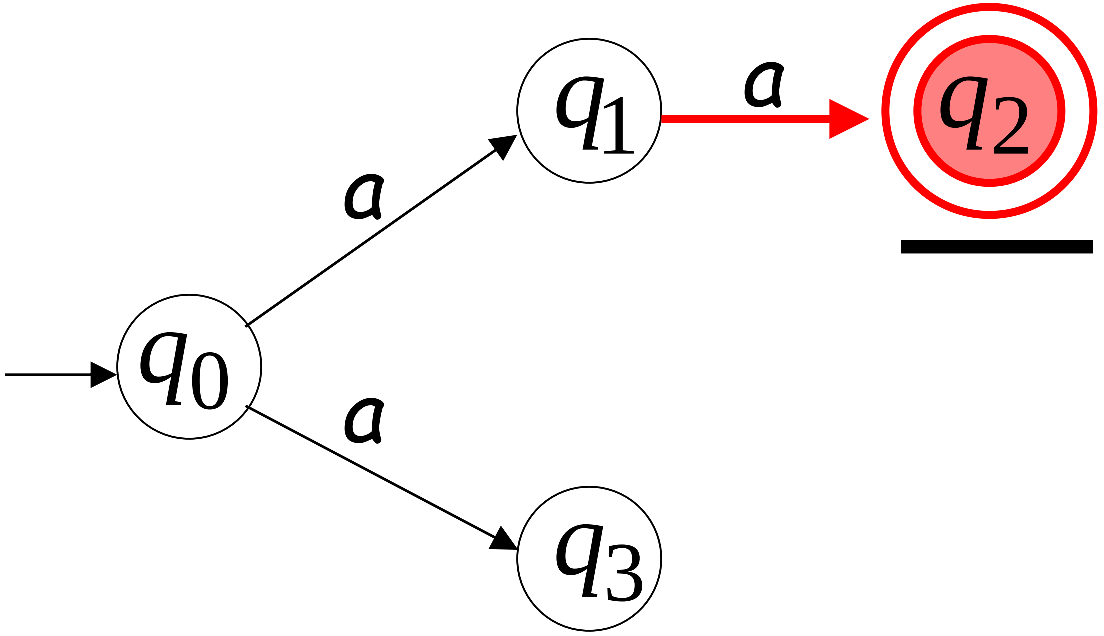


No transition:
the automaton hangs

First Choice

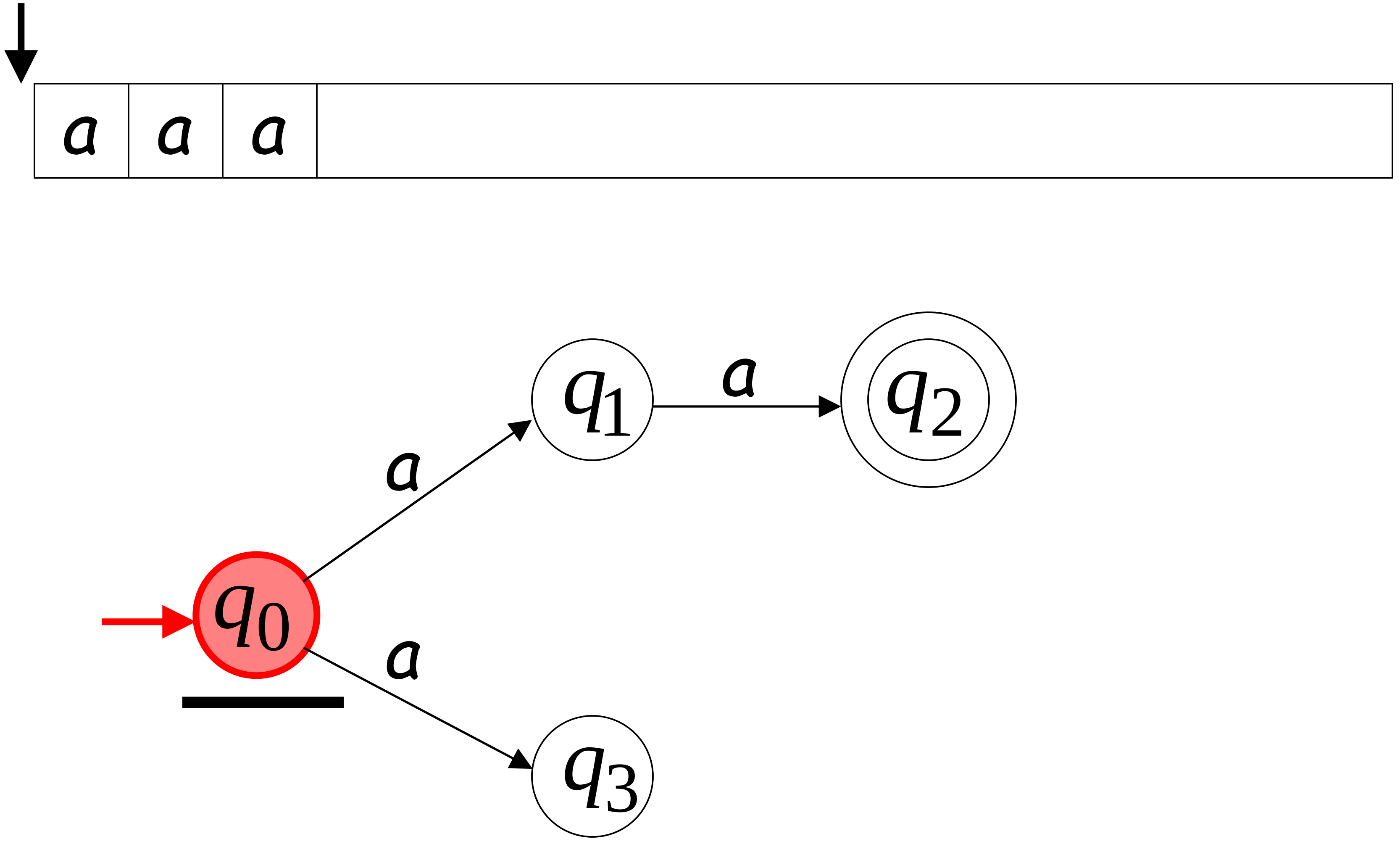


Input cannot be consumed

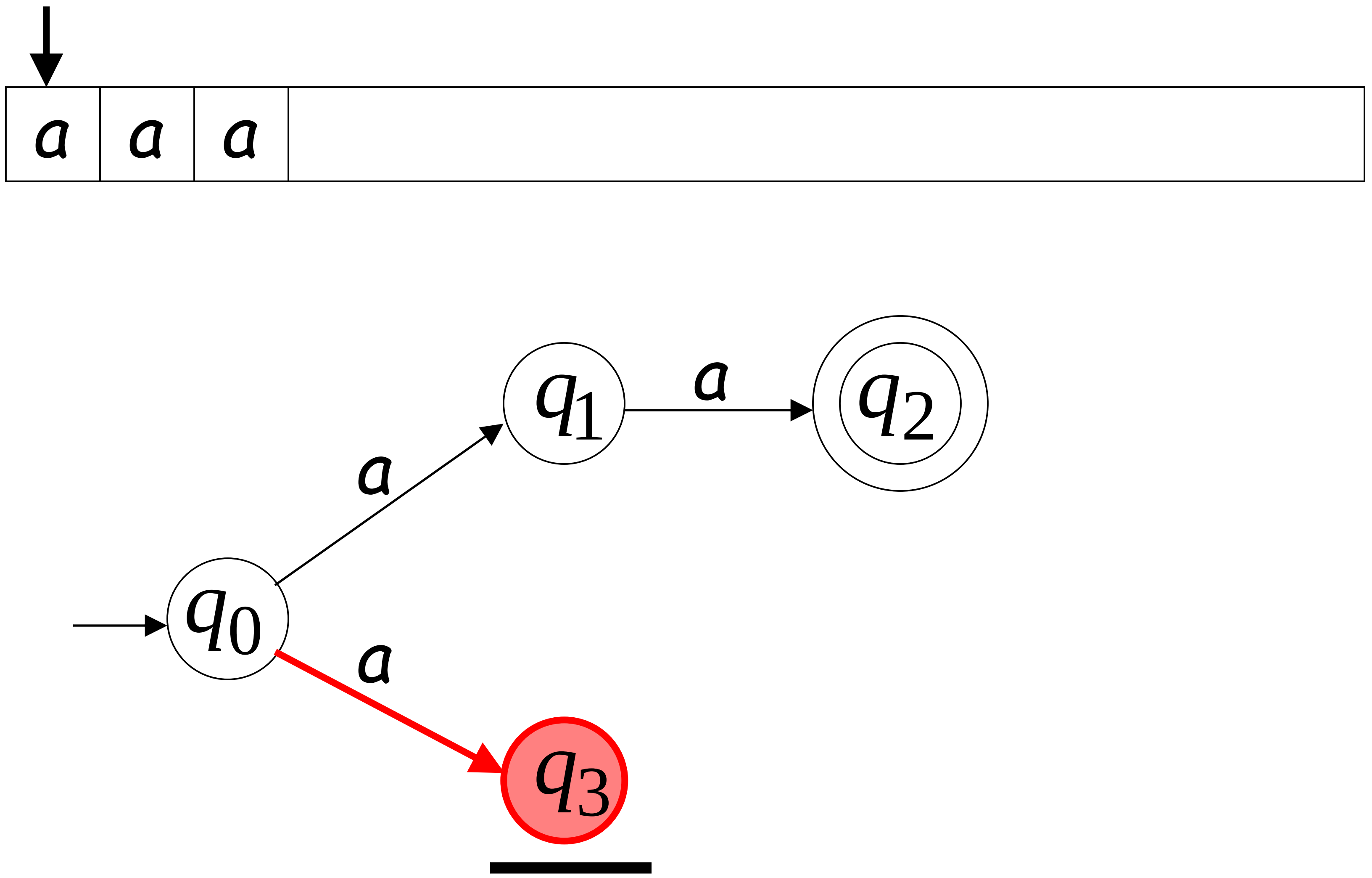


"reject"

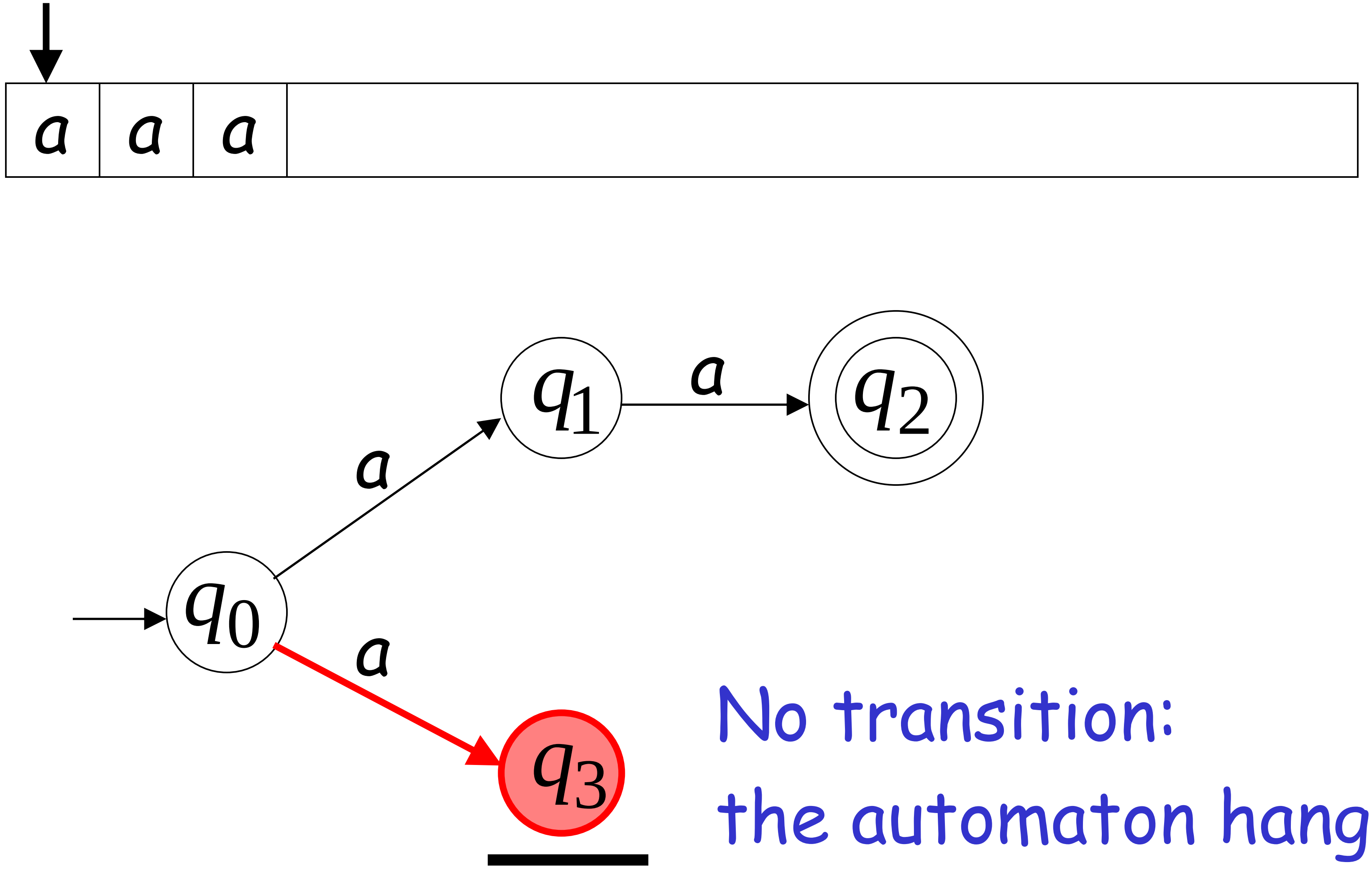
Second Choice



Second Choice

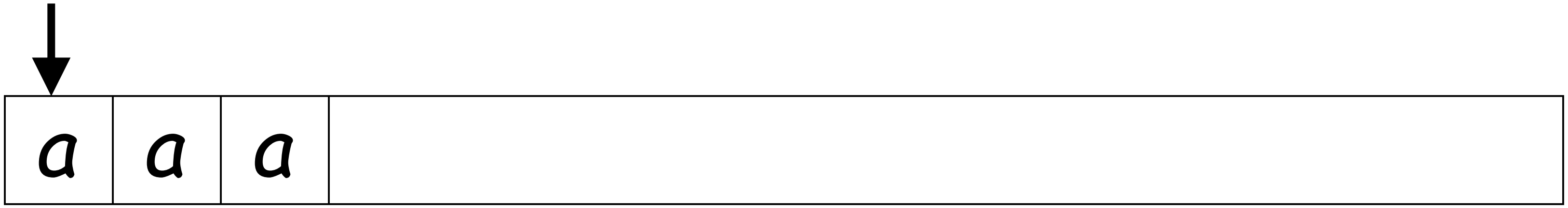


Second Choice

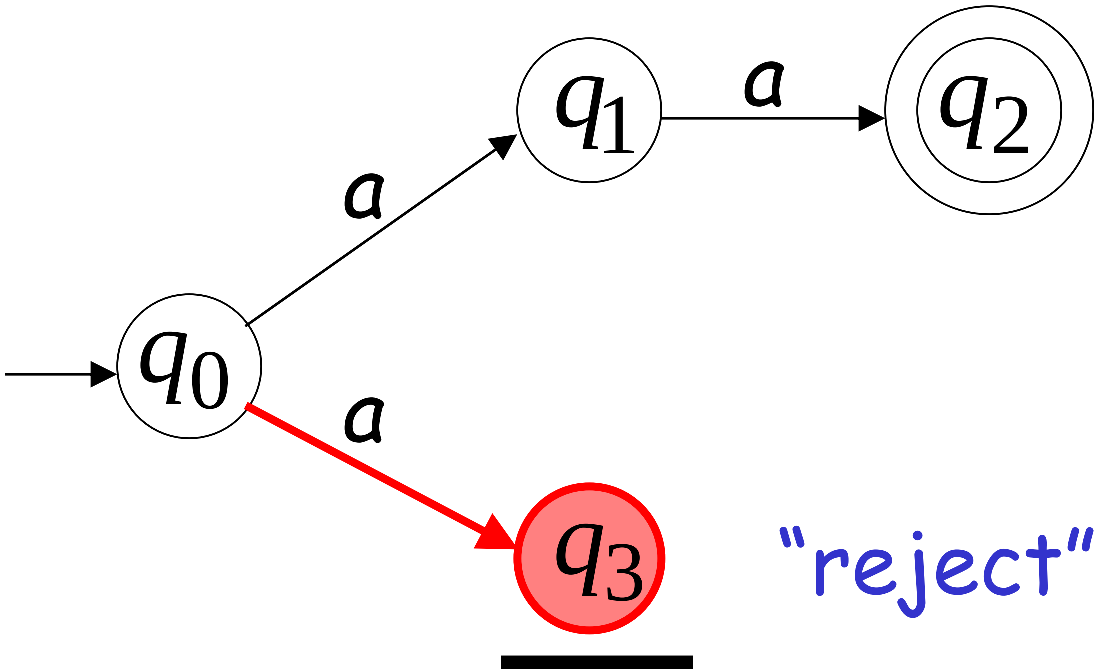


No transition:
the automaton hangs

Second Choice

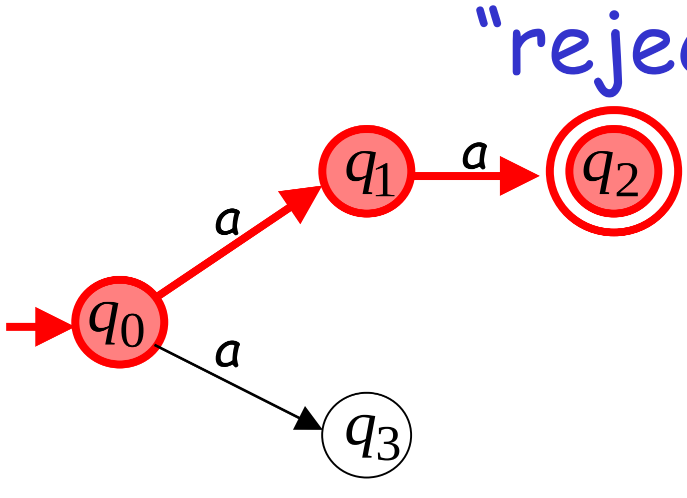


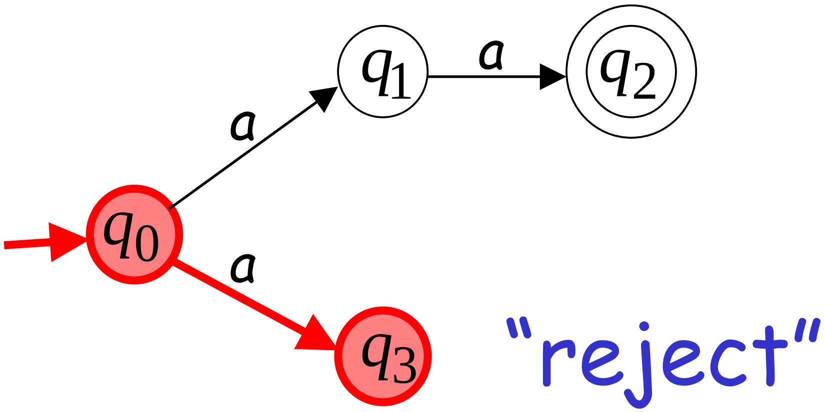
Input cannot be consumed



"reject"

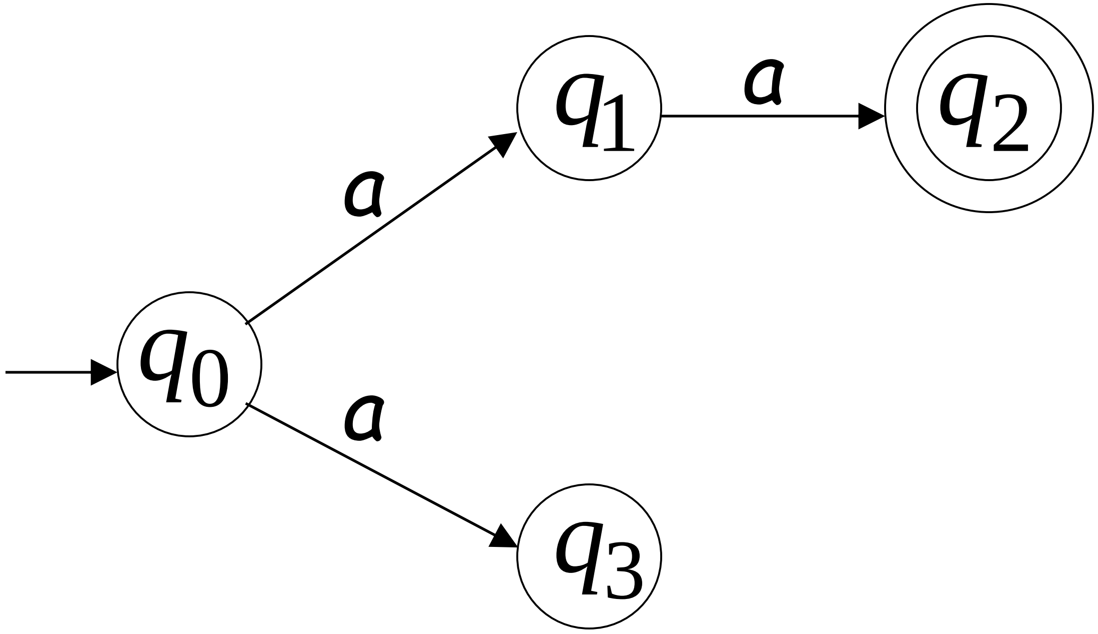
aaa is rejected by the NFA:



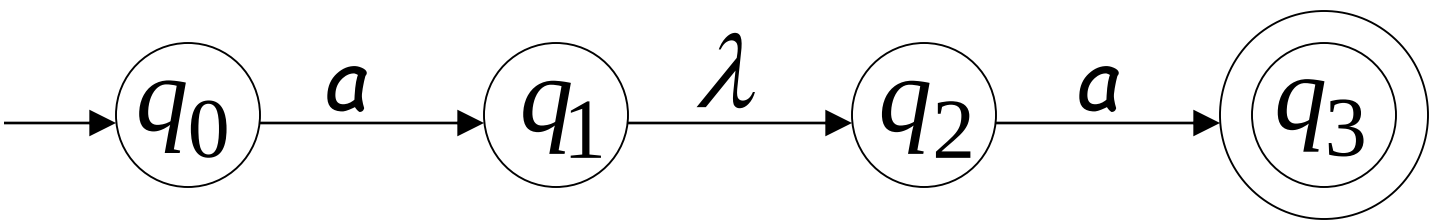


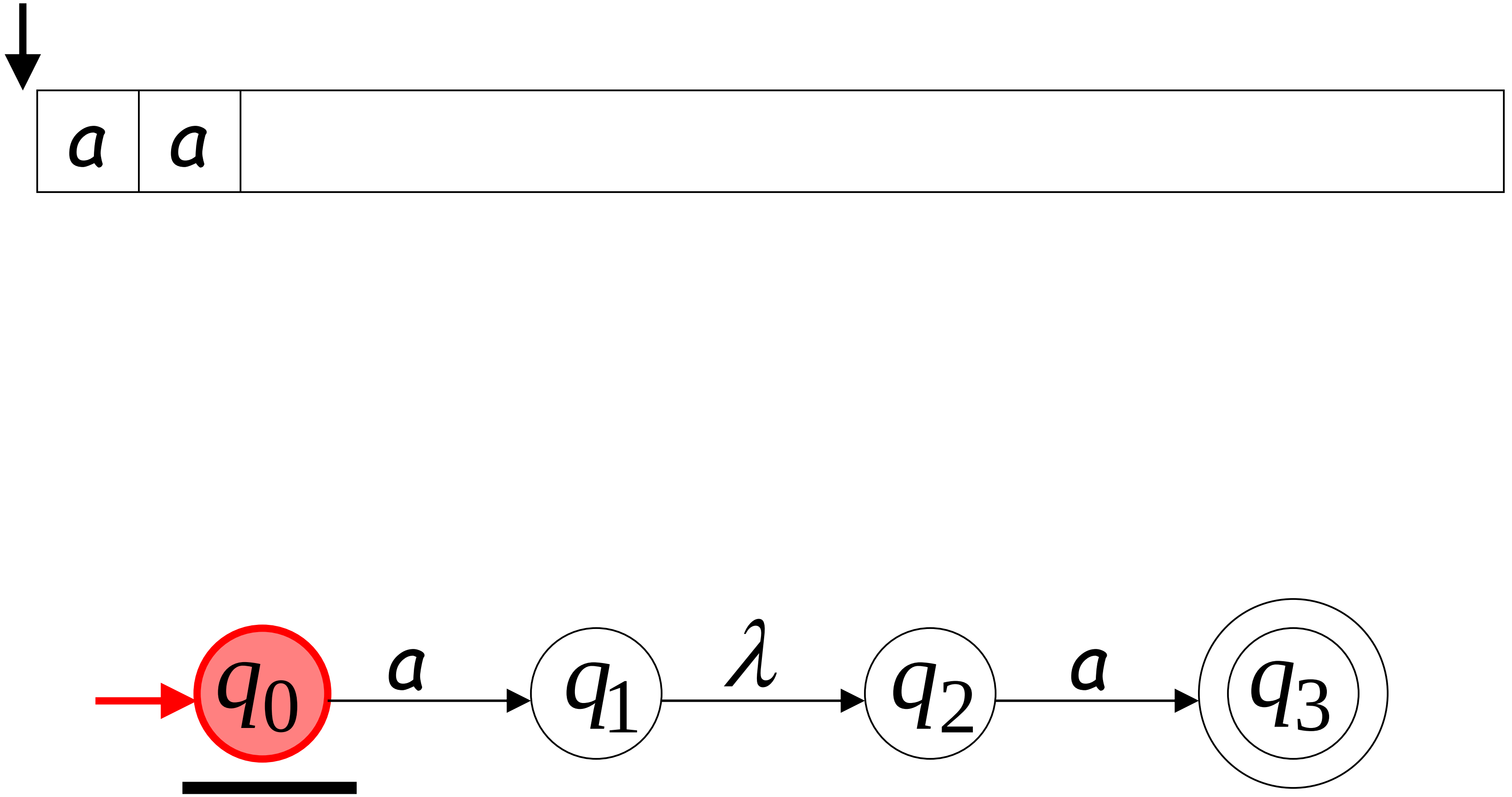
All possible computations lead to rejection

Language accepted: $L = \{aa\}$

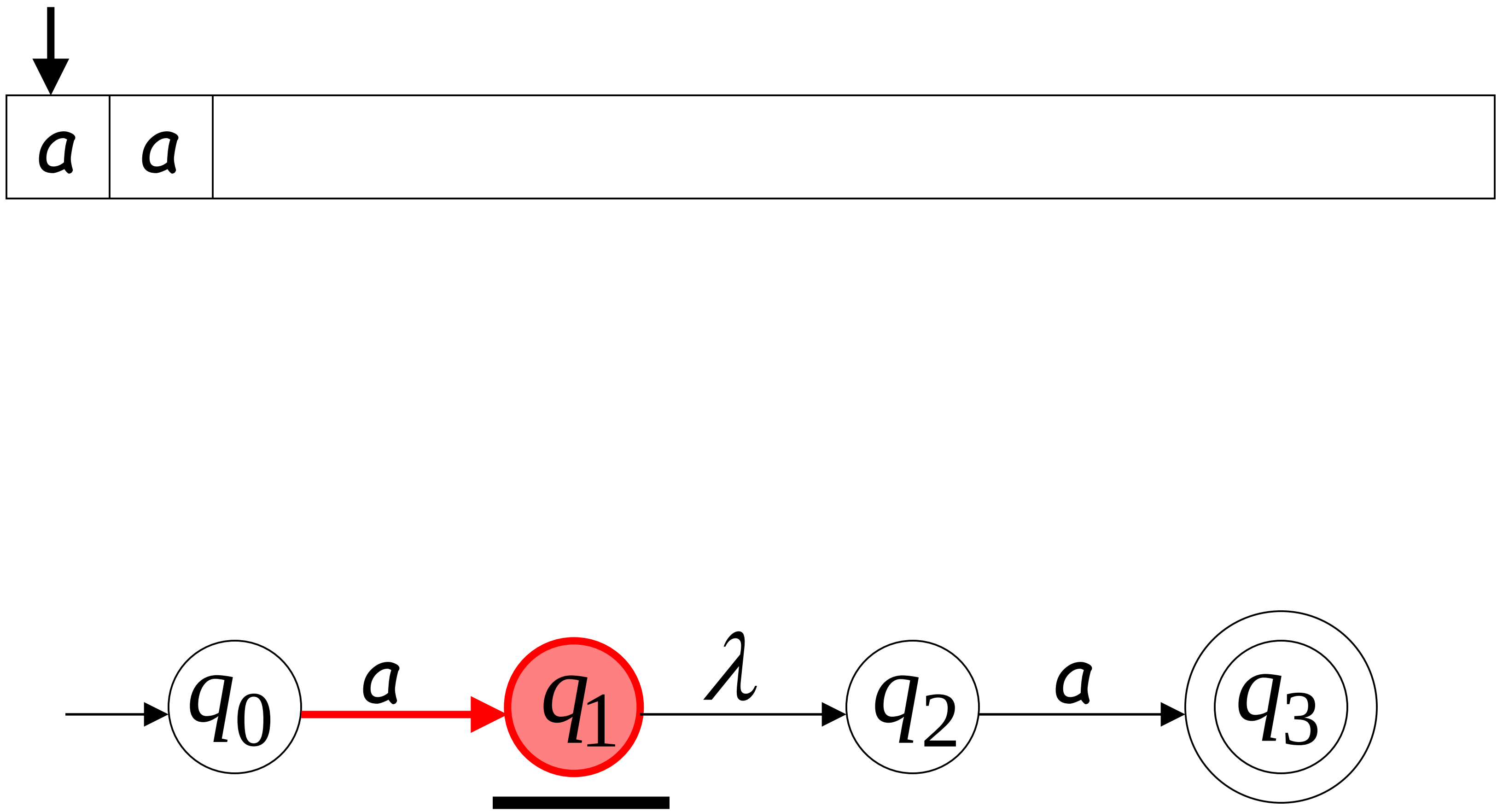


Lambda Transitions

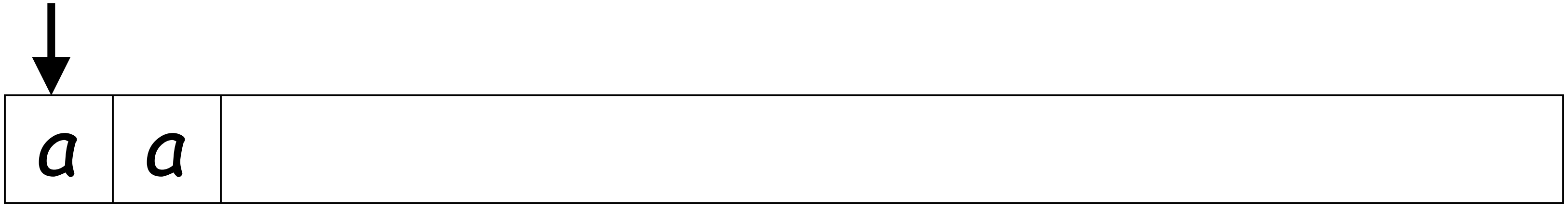


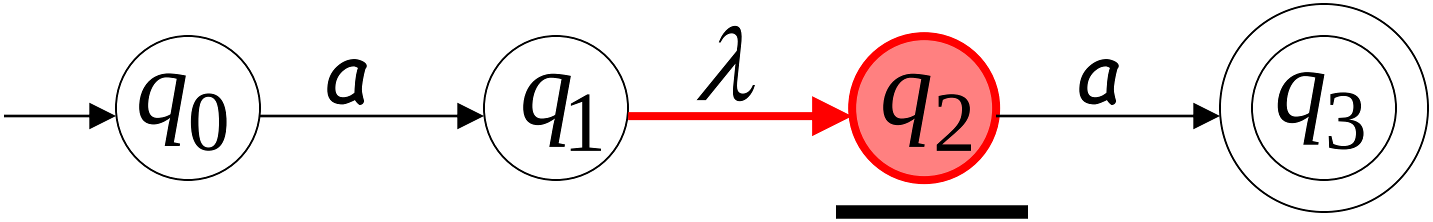


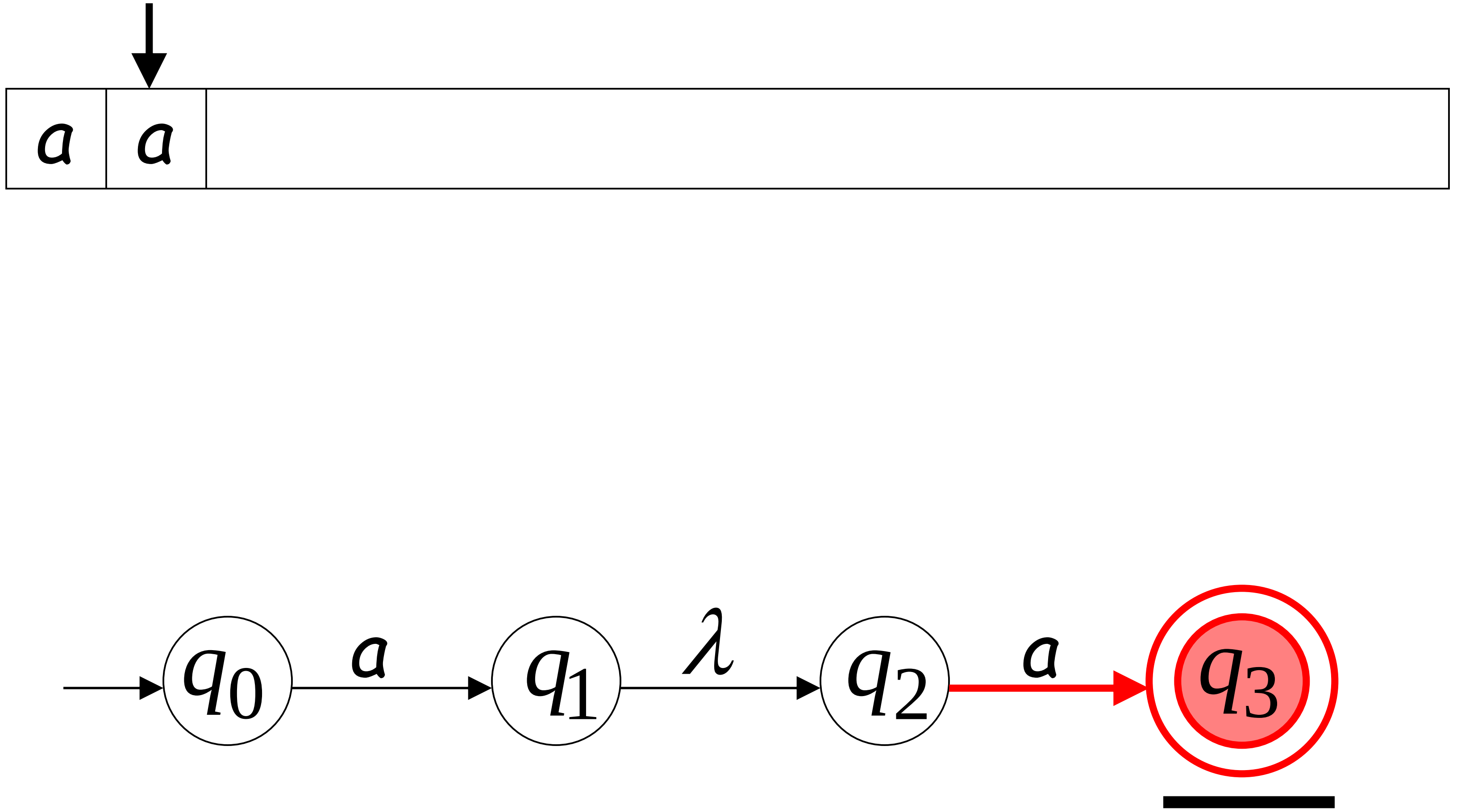
a
a
q_0
a
q_1
λ
q_2
a
q_3


a
a
q_0
a
q_1
λ
q_2
a
q_3

(read head doesn't move)

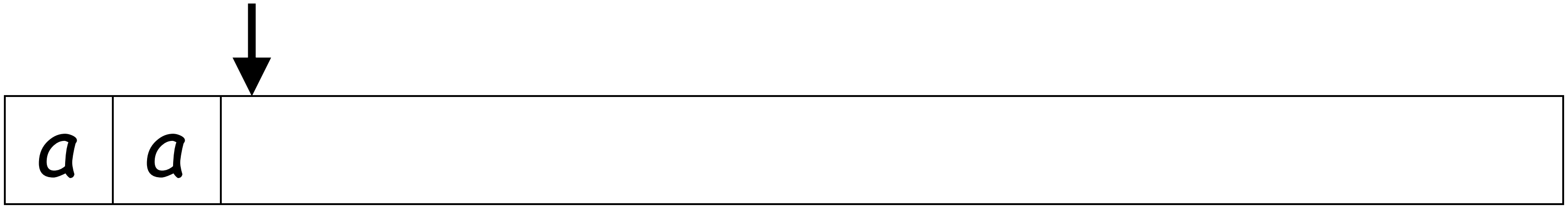


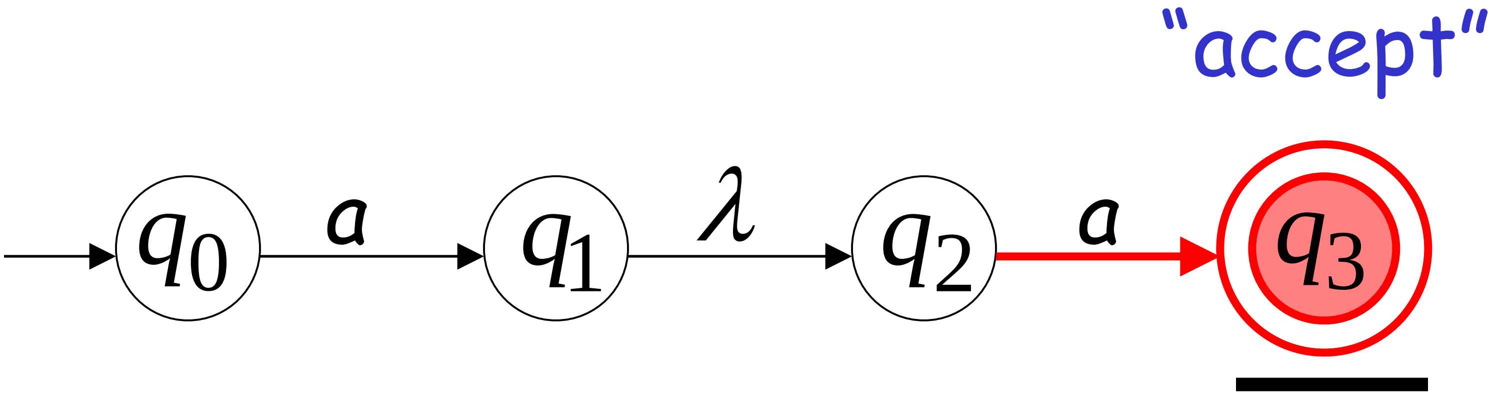




a
a
q_0
a
q_1
λ
q_2
a
q_3

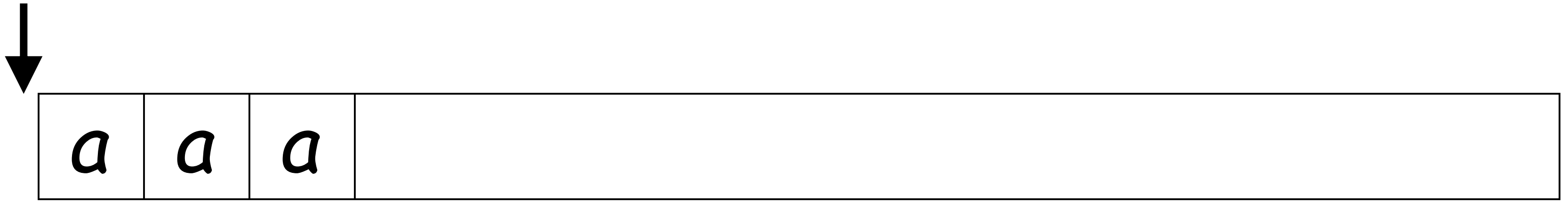
all input is consumed

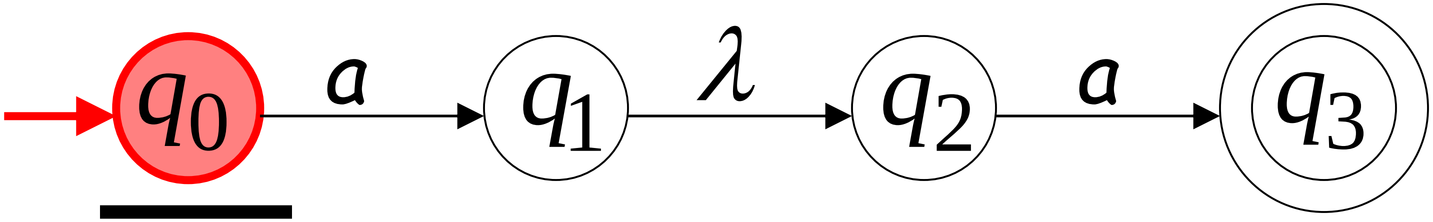


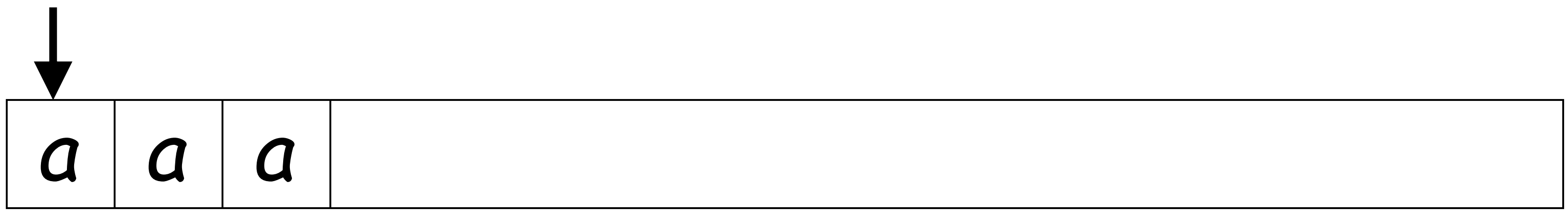


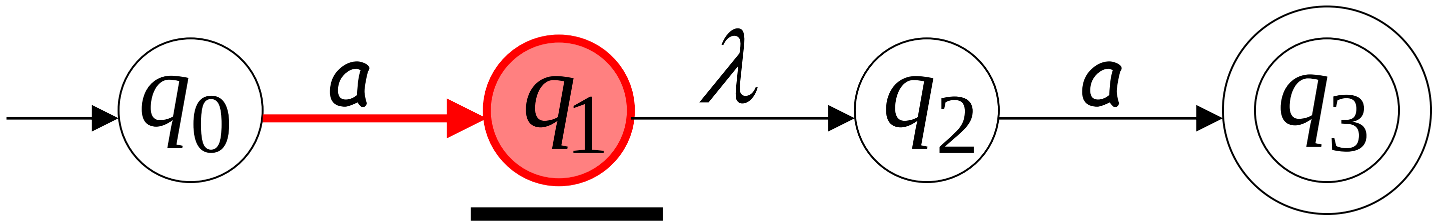
String aa is accepted

Rejection Example

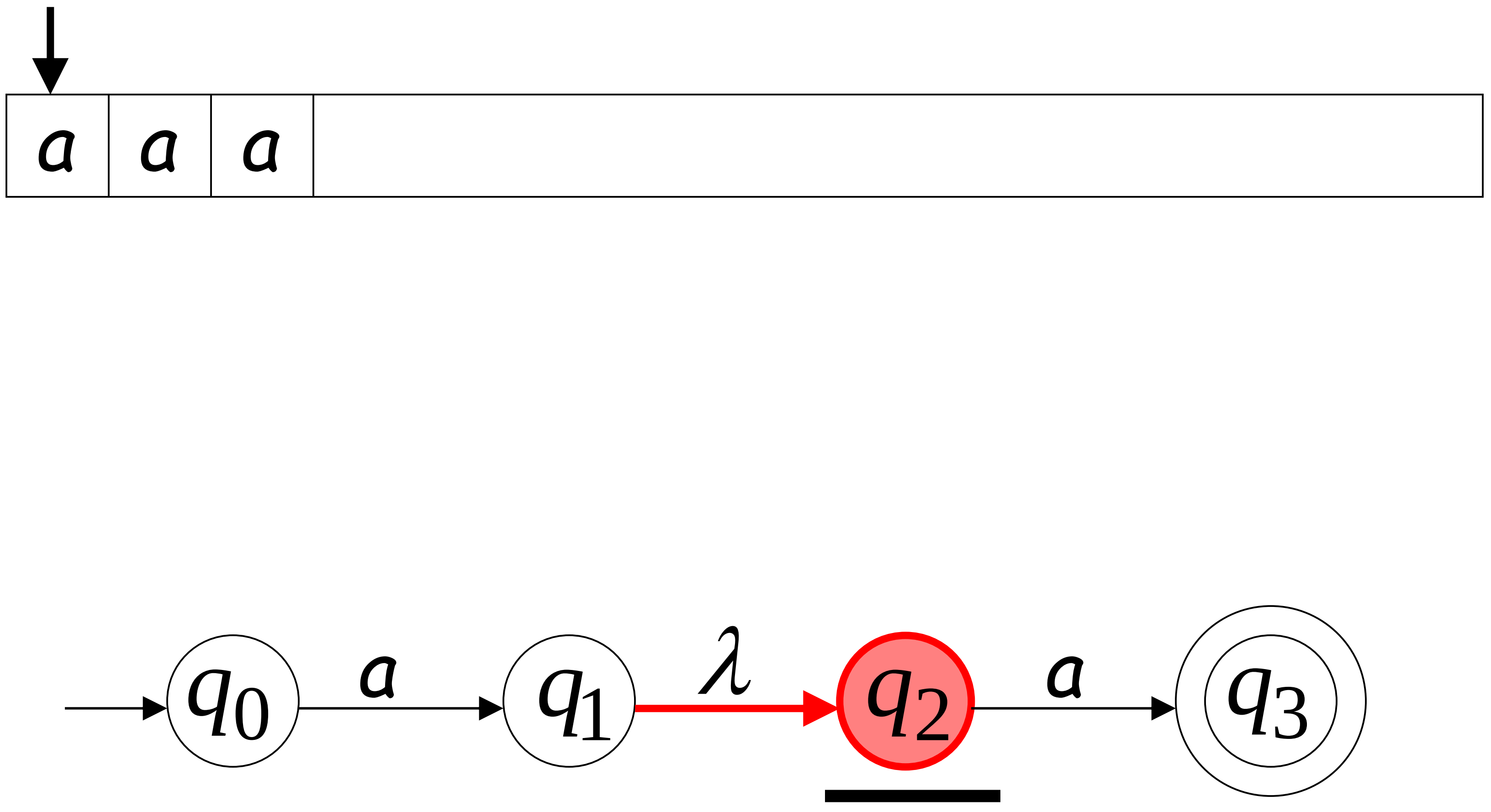


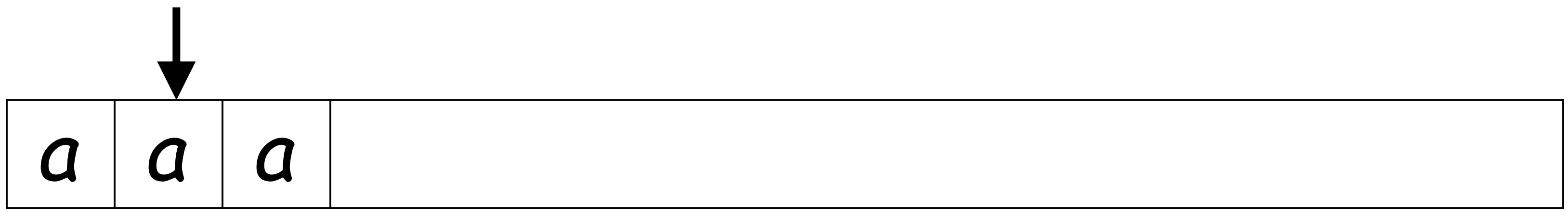


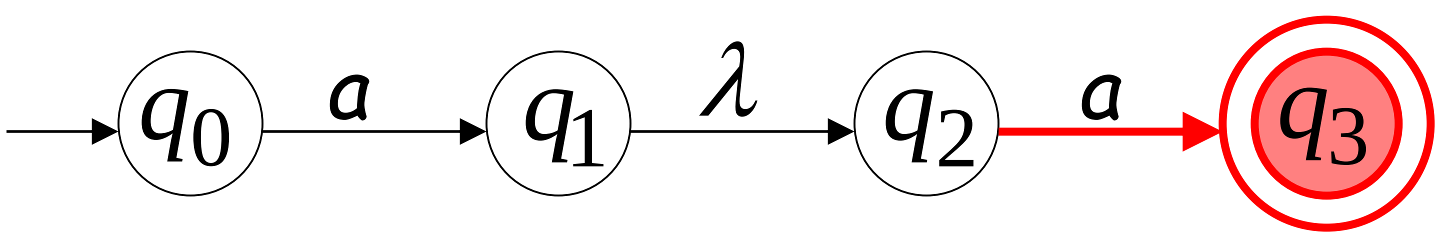

a
a
a


q_0
a
q_1
λ
q_2
a
q_3

(read head doesn't move)



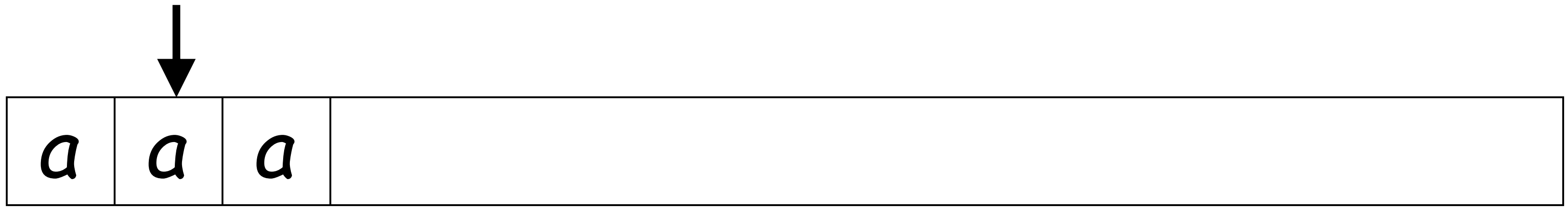


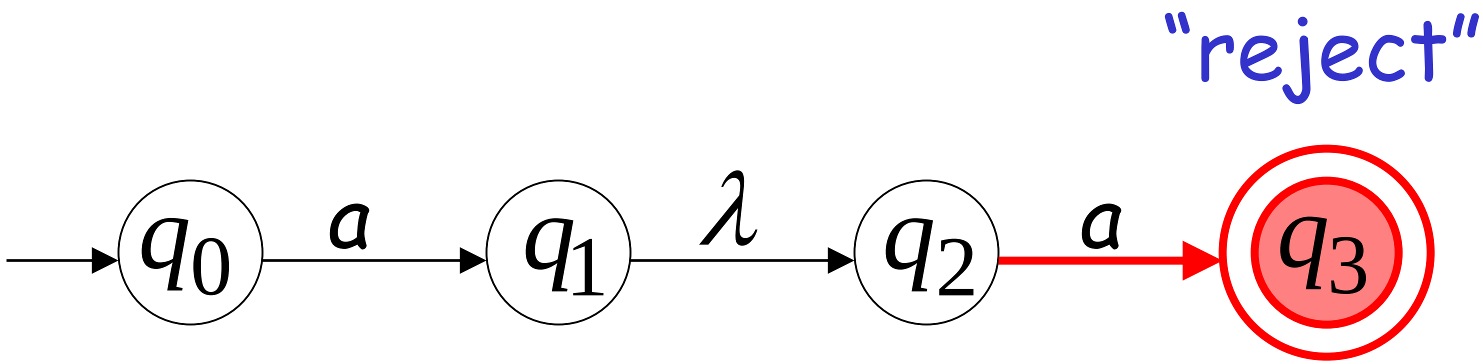


No transition:

the automaton hangs

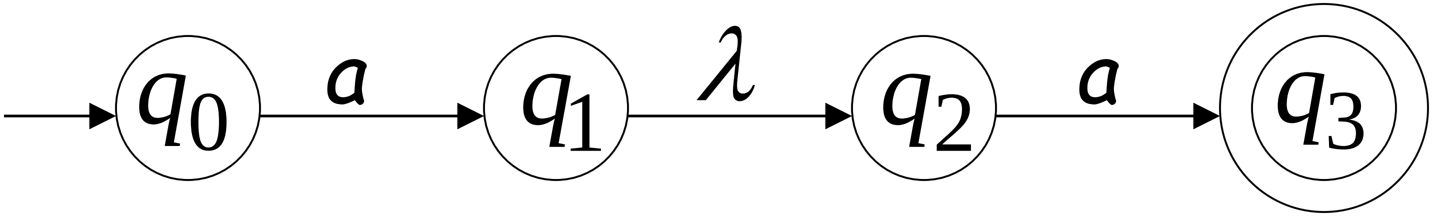
Input cannot be consumed



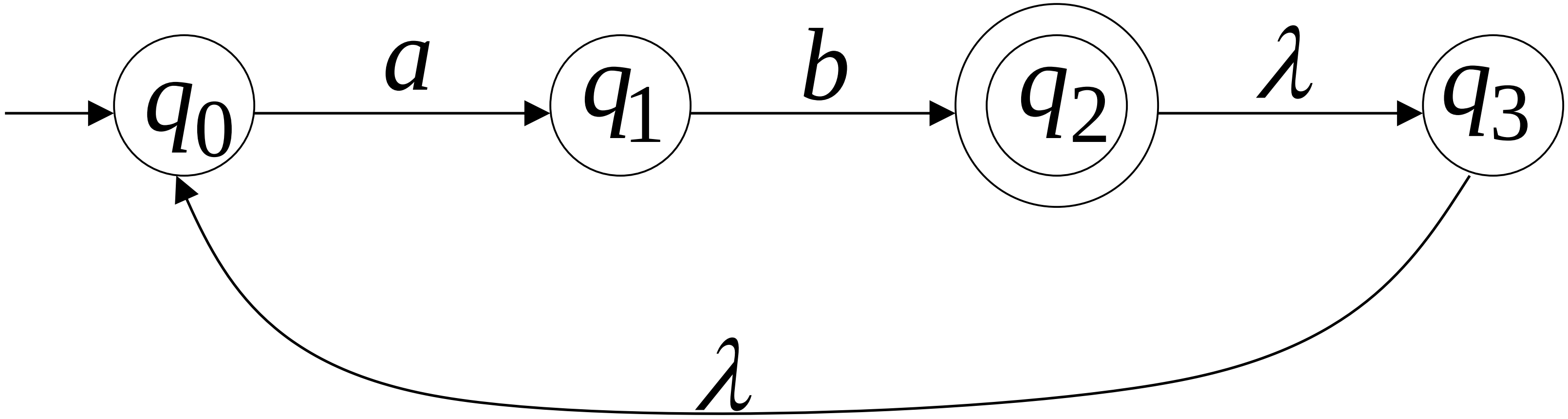


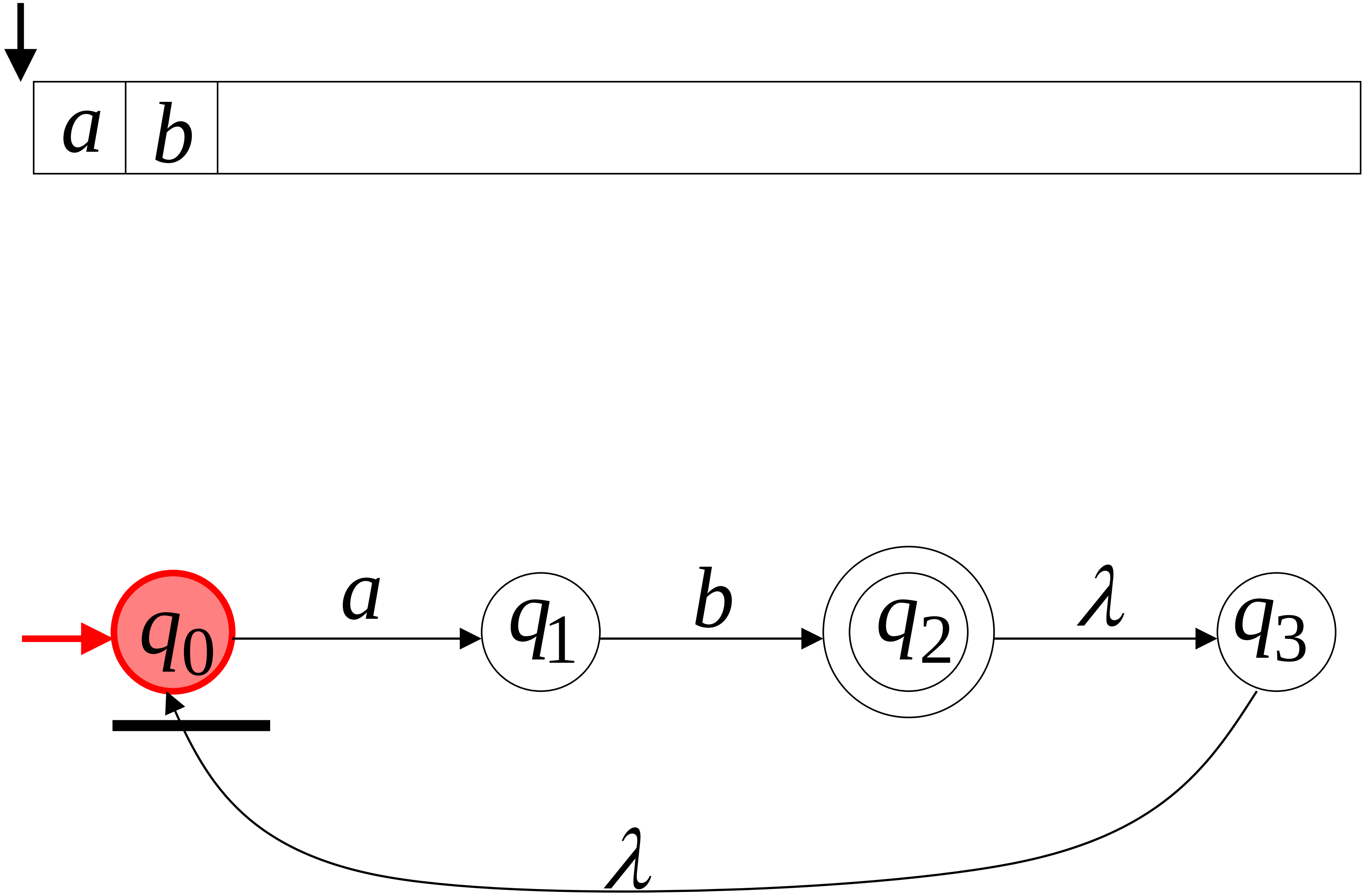
String aaa is rejected

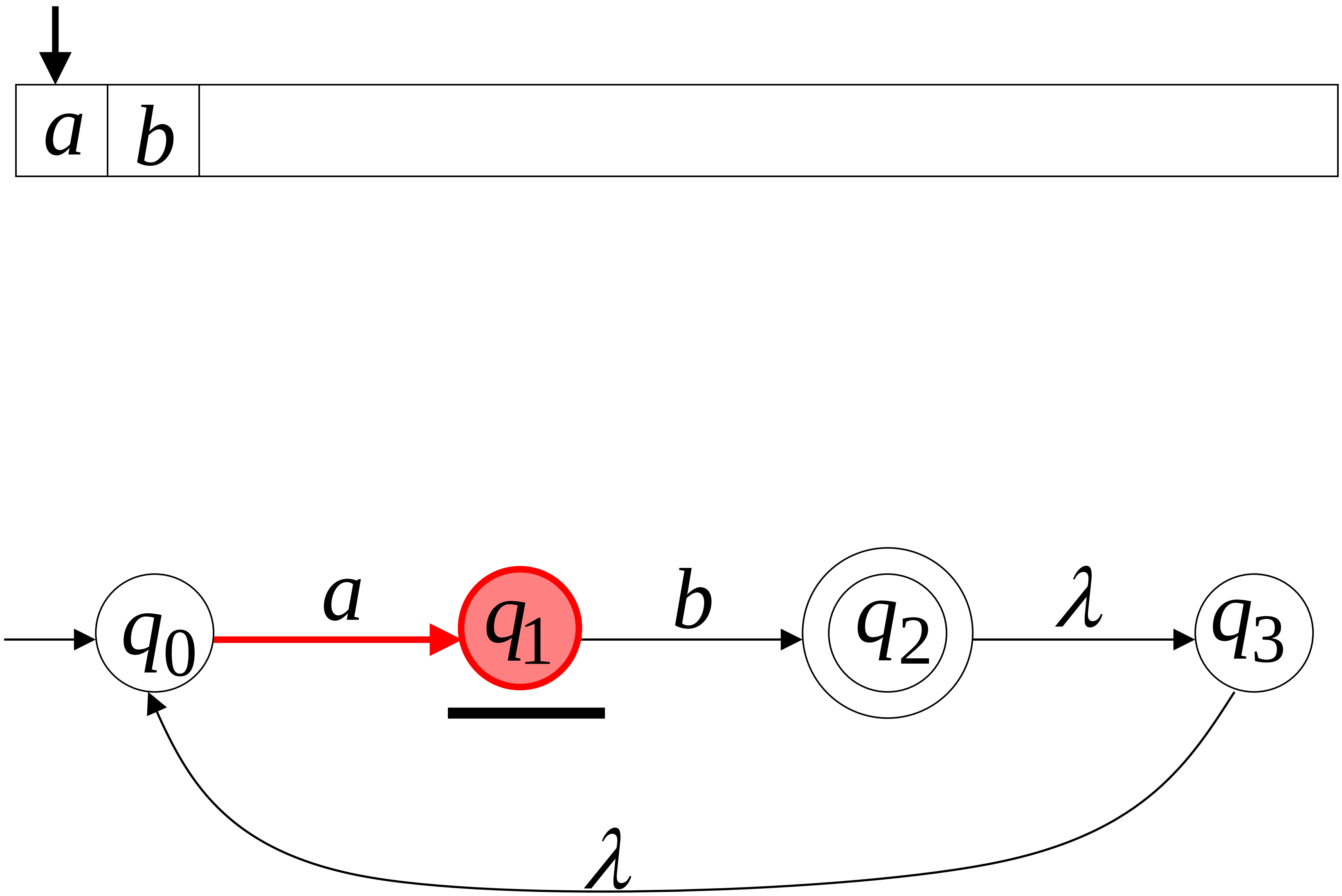
Language accepted: $L = \{aa\}$

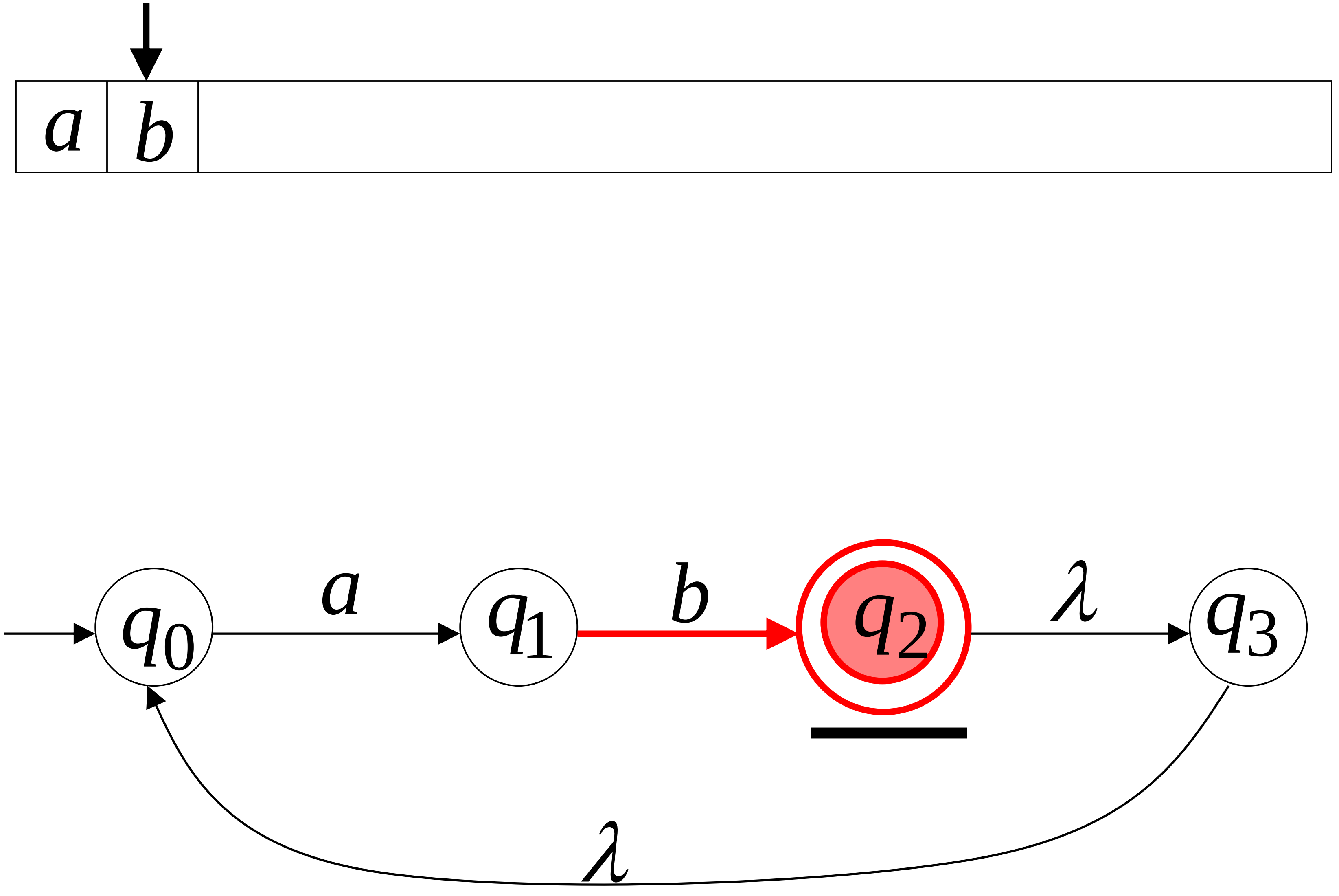


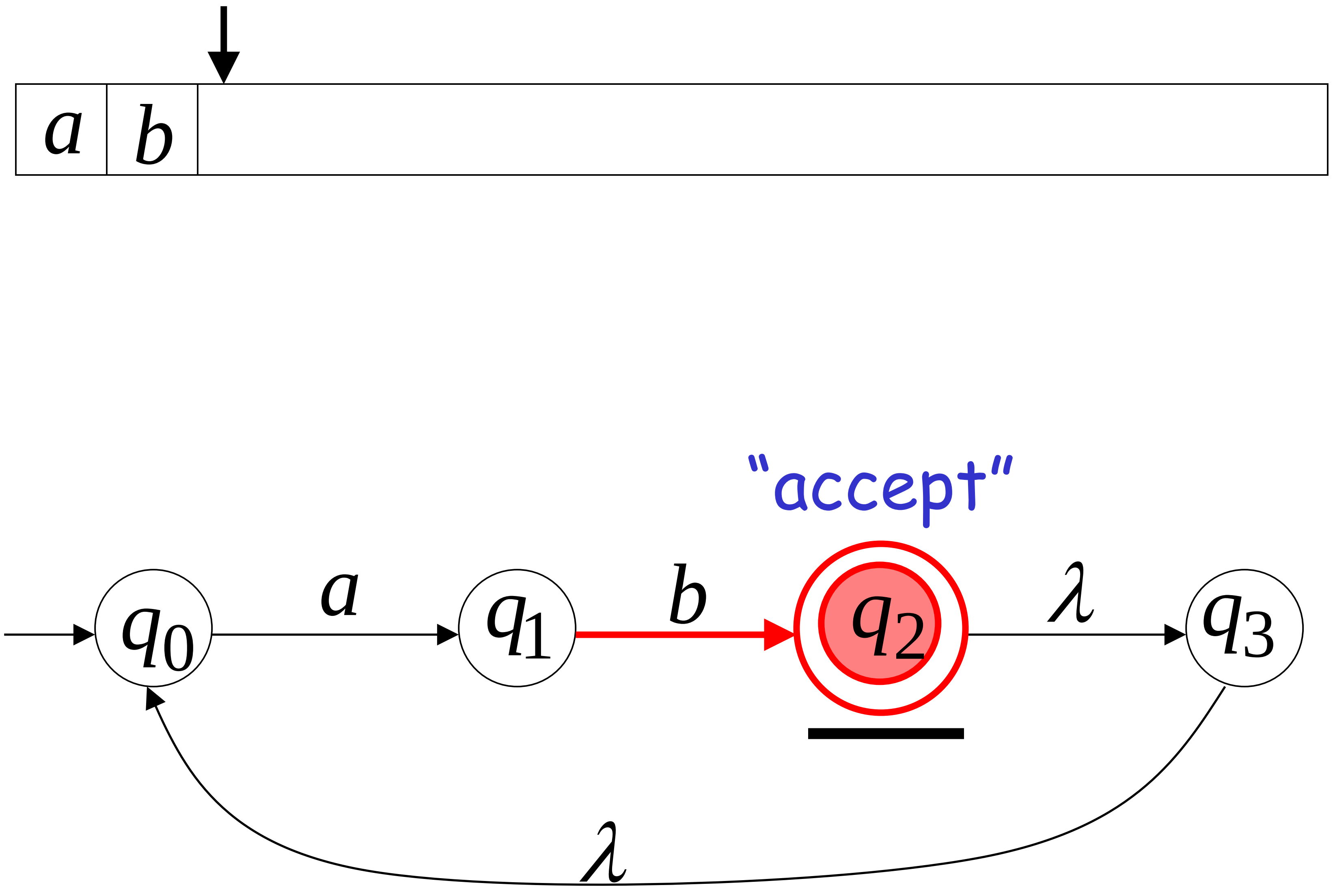
Another NFA Example



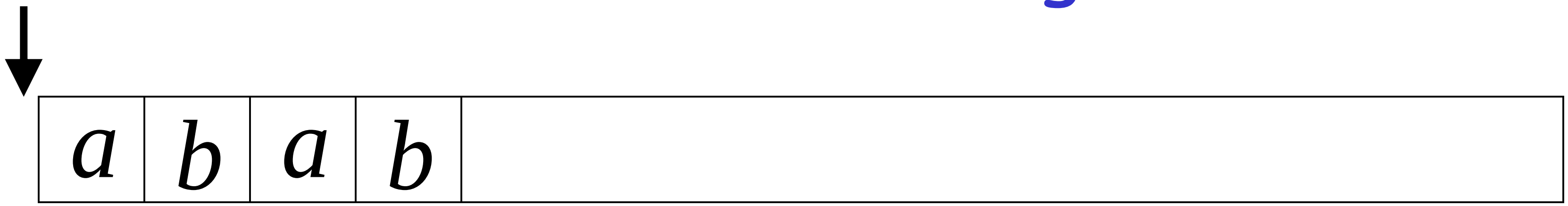

a
b
a
b
λ
λ
q_0
q_1
q_2
q_3

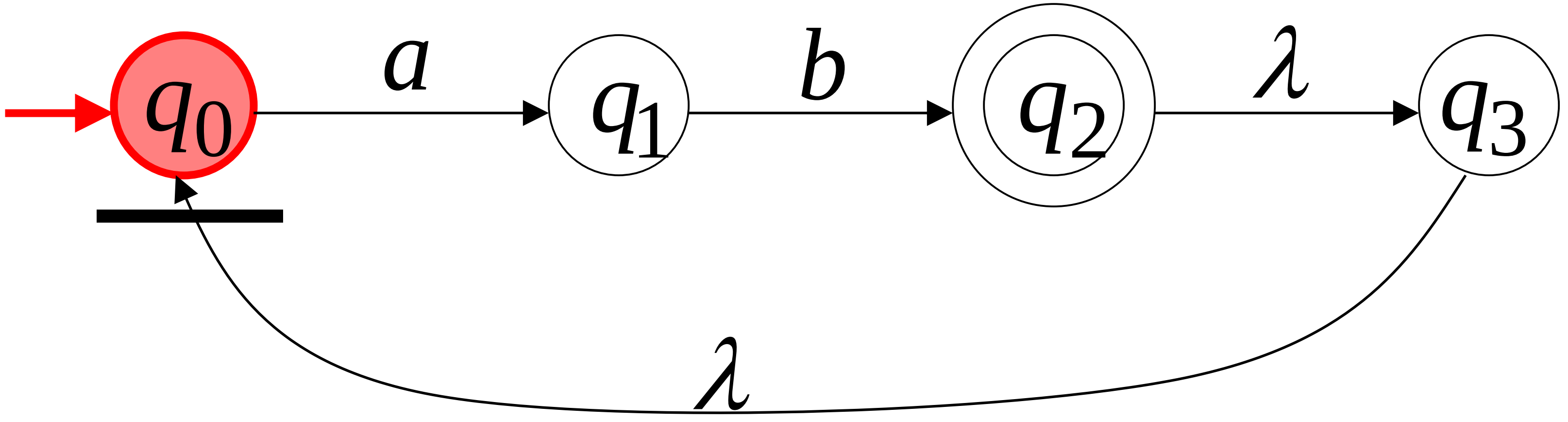

a
b
q0
a
q1
b
q2
λ
q3
λ

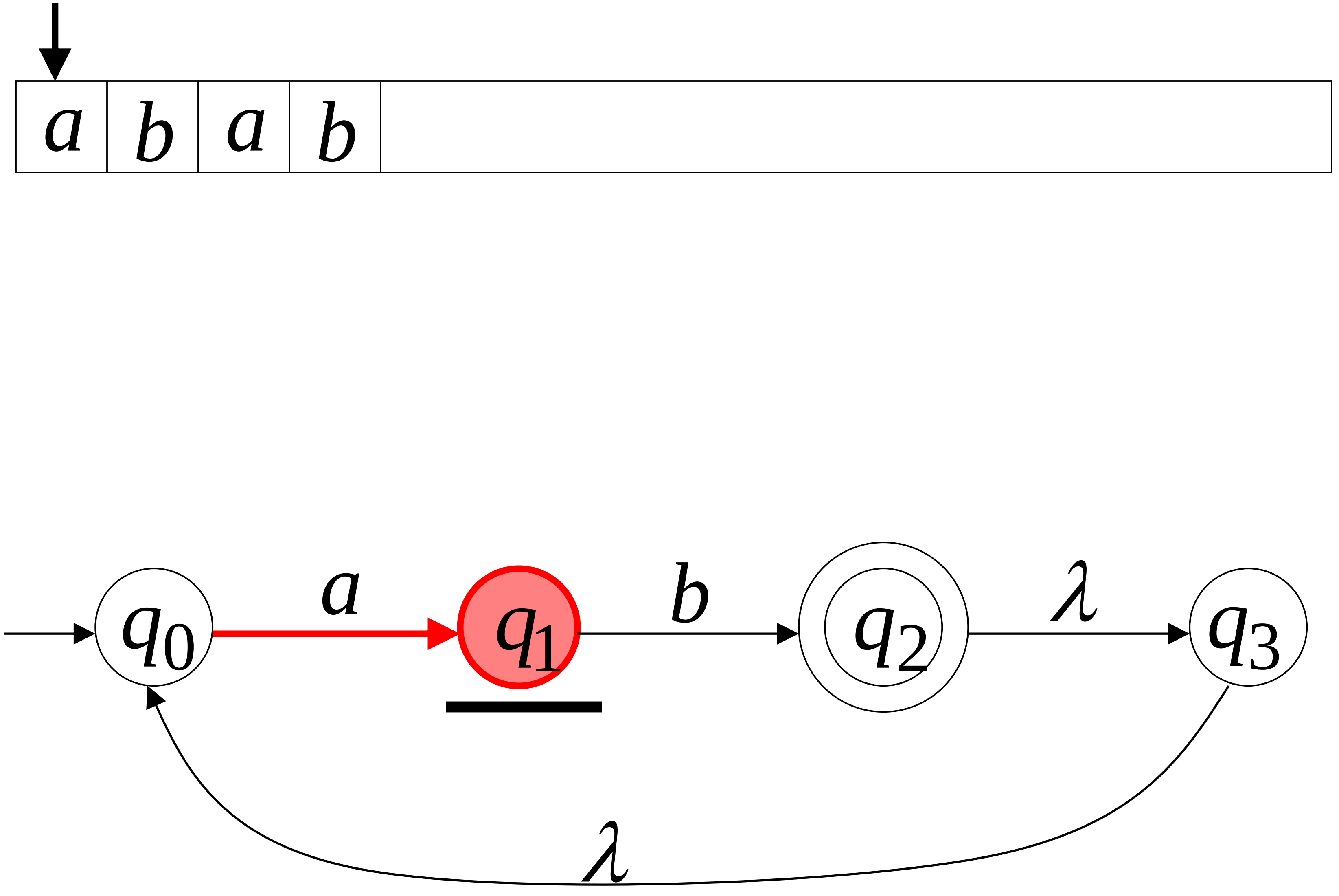

a
b
a
b
λ
λ
q0
q1
q2
q3

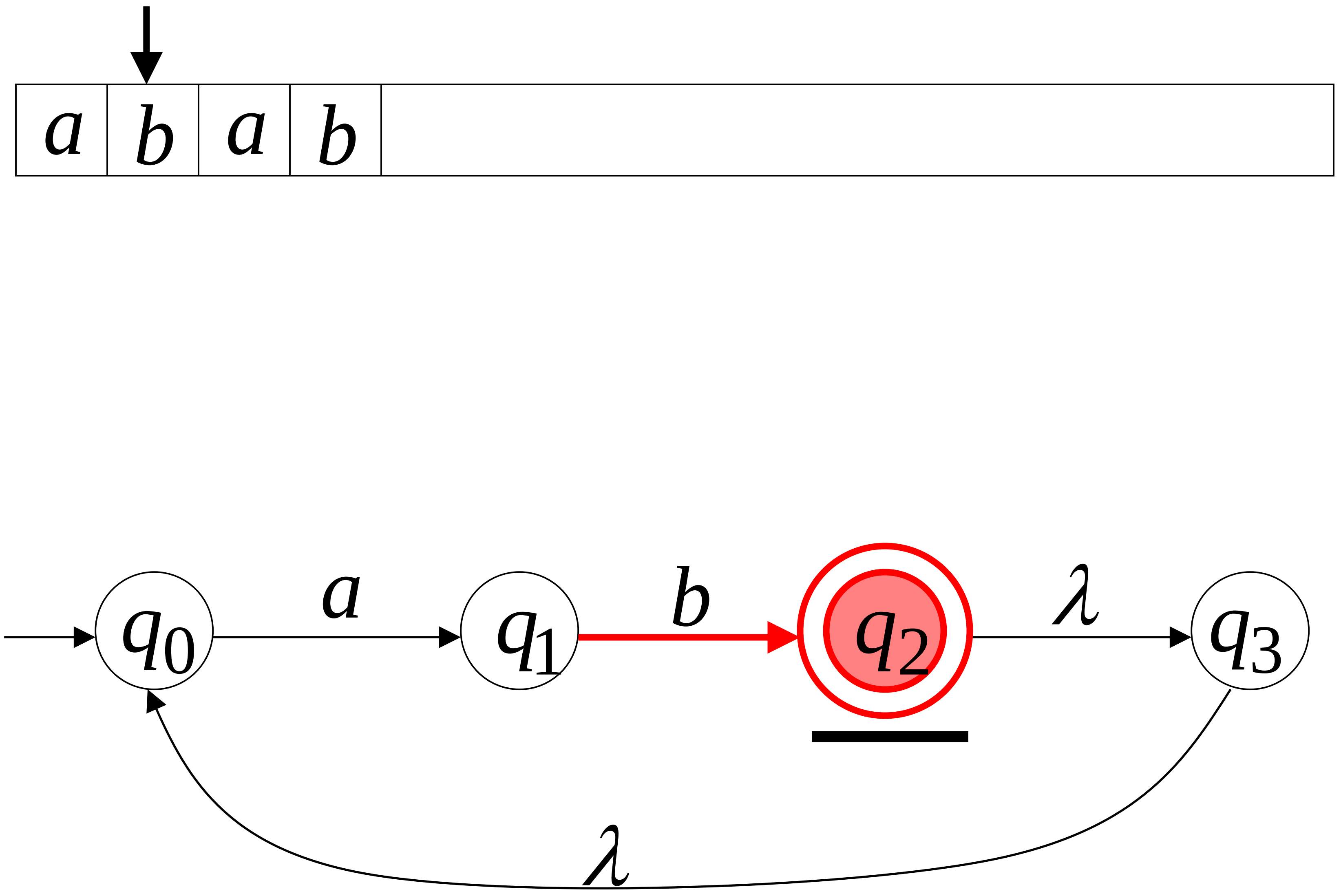

a
b
"accept"
a
b
λ
λ
q_0
q_1
q_2
q_3

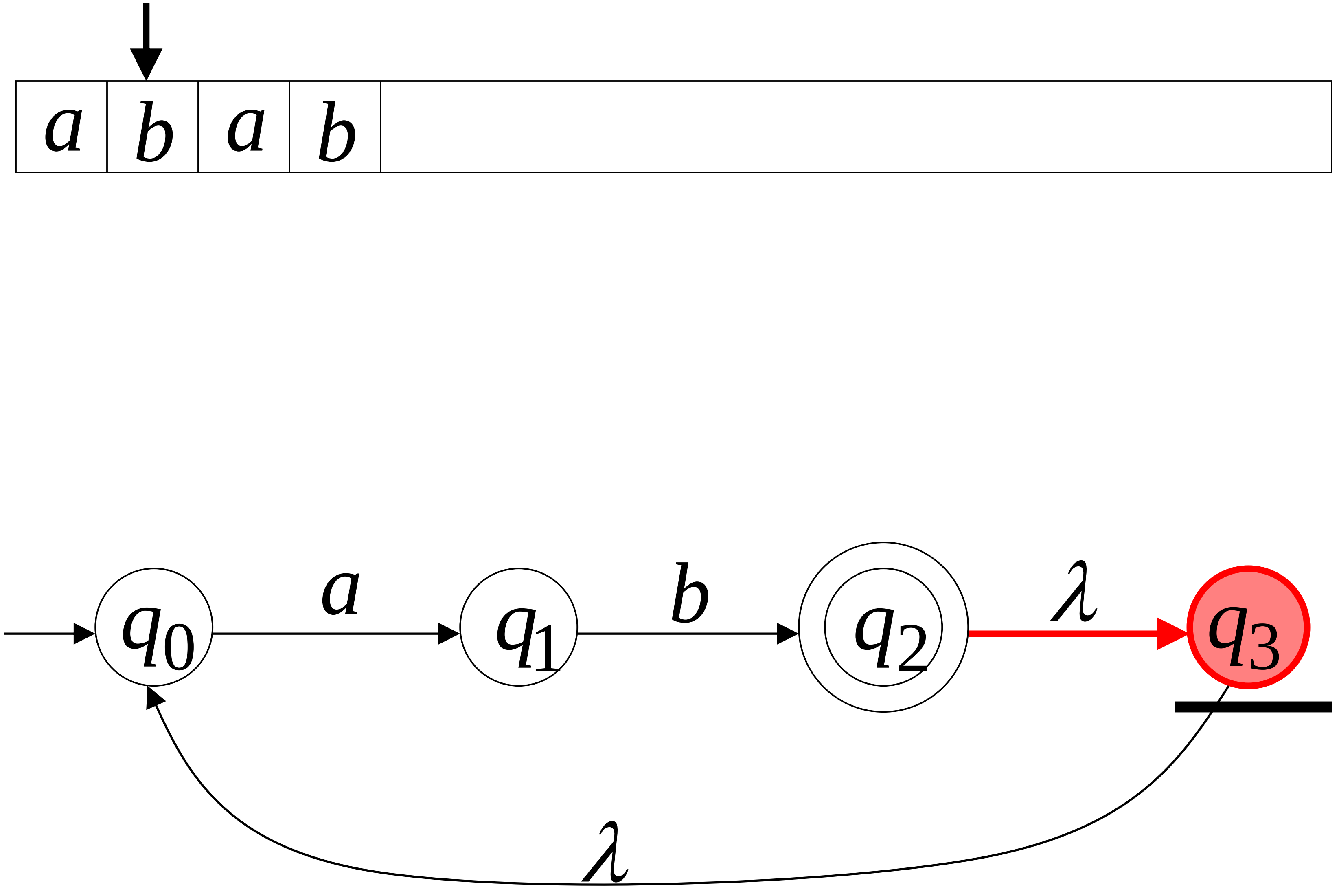
Another String

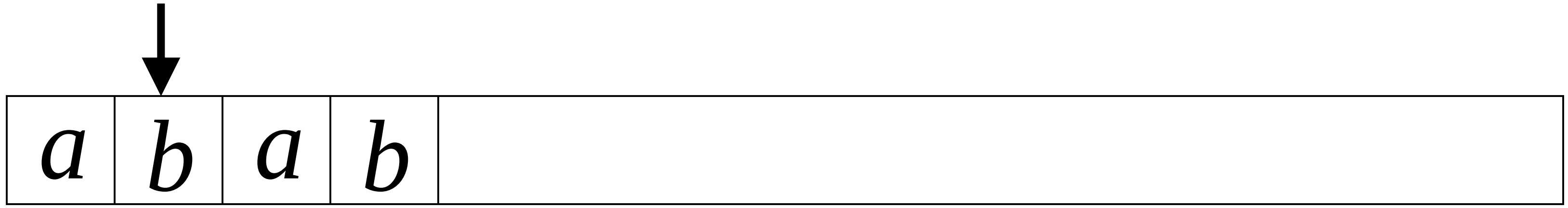


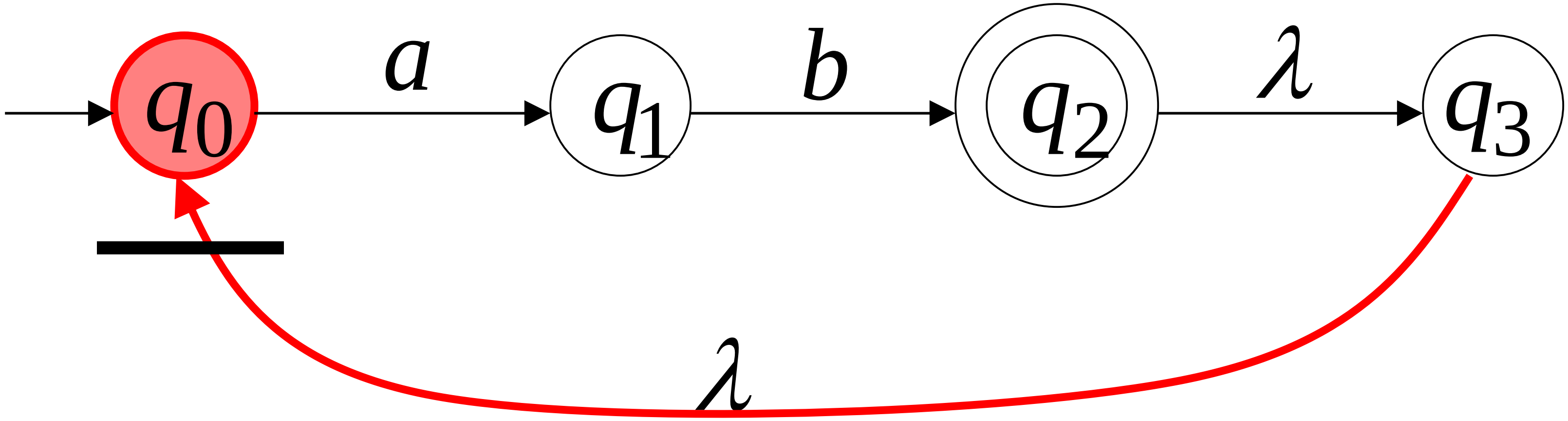


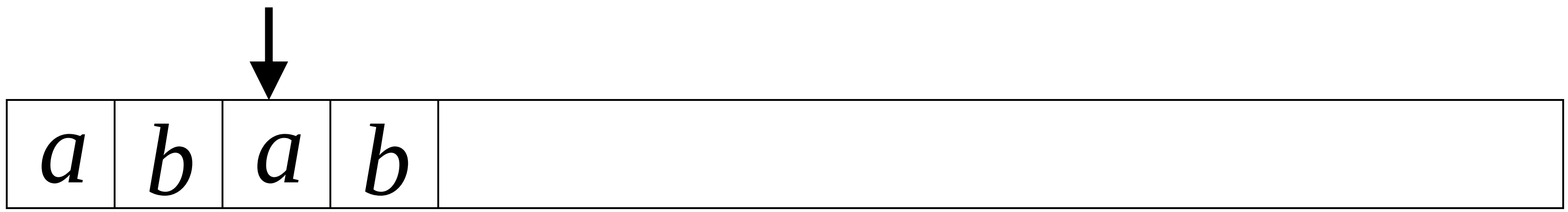

a
b
a
b
a
b
λ
λ
q_0
q_1
q_2
q_3

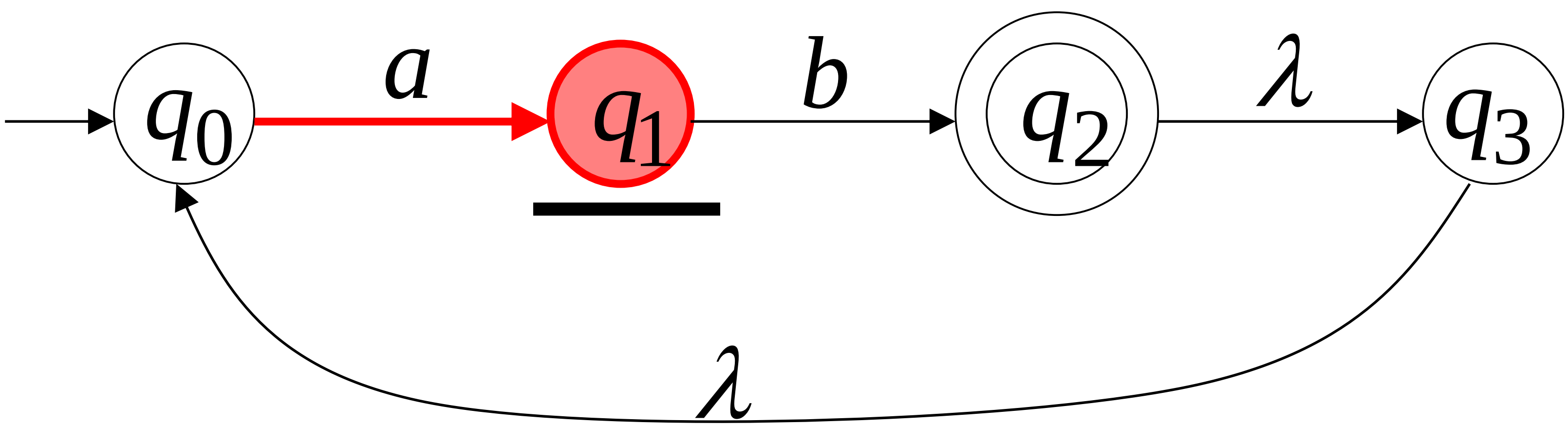

a
b
a
b
a
b
λ
λ
q_0
q_1
q_2
q_3

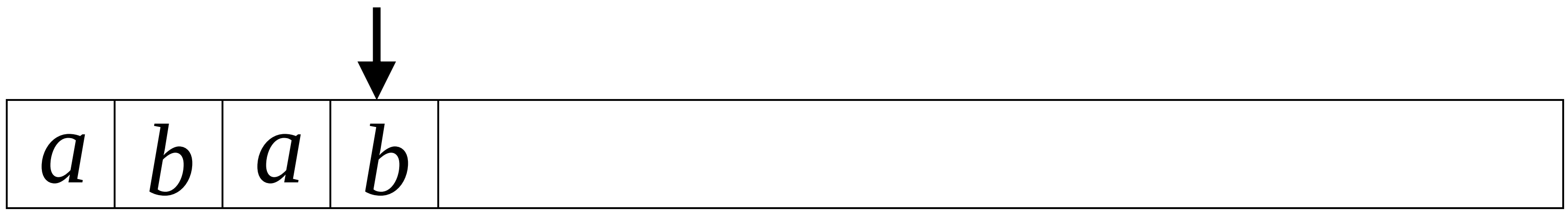

a
b
a
b
a
b
λ
λ
q_0
q_1
q_2
q_3

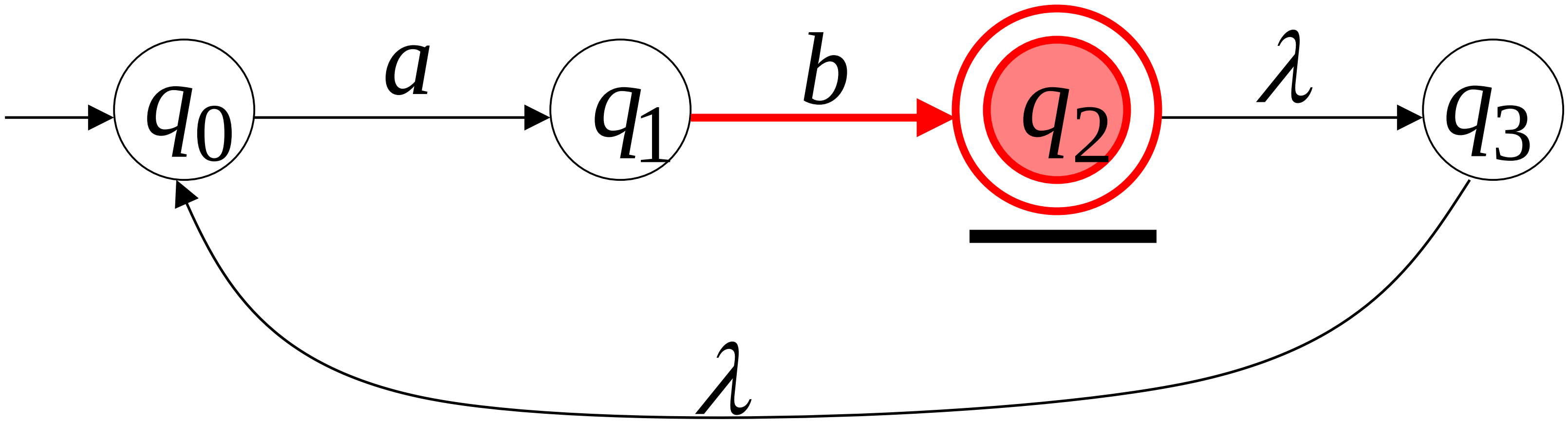

a
b
a
b

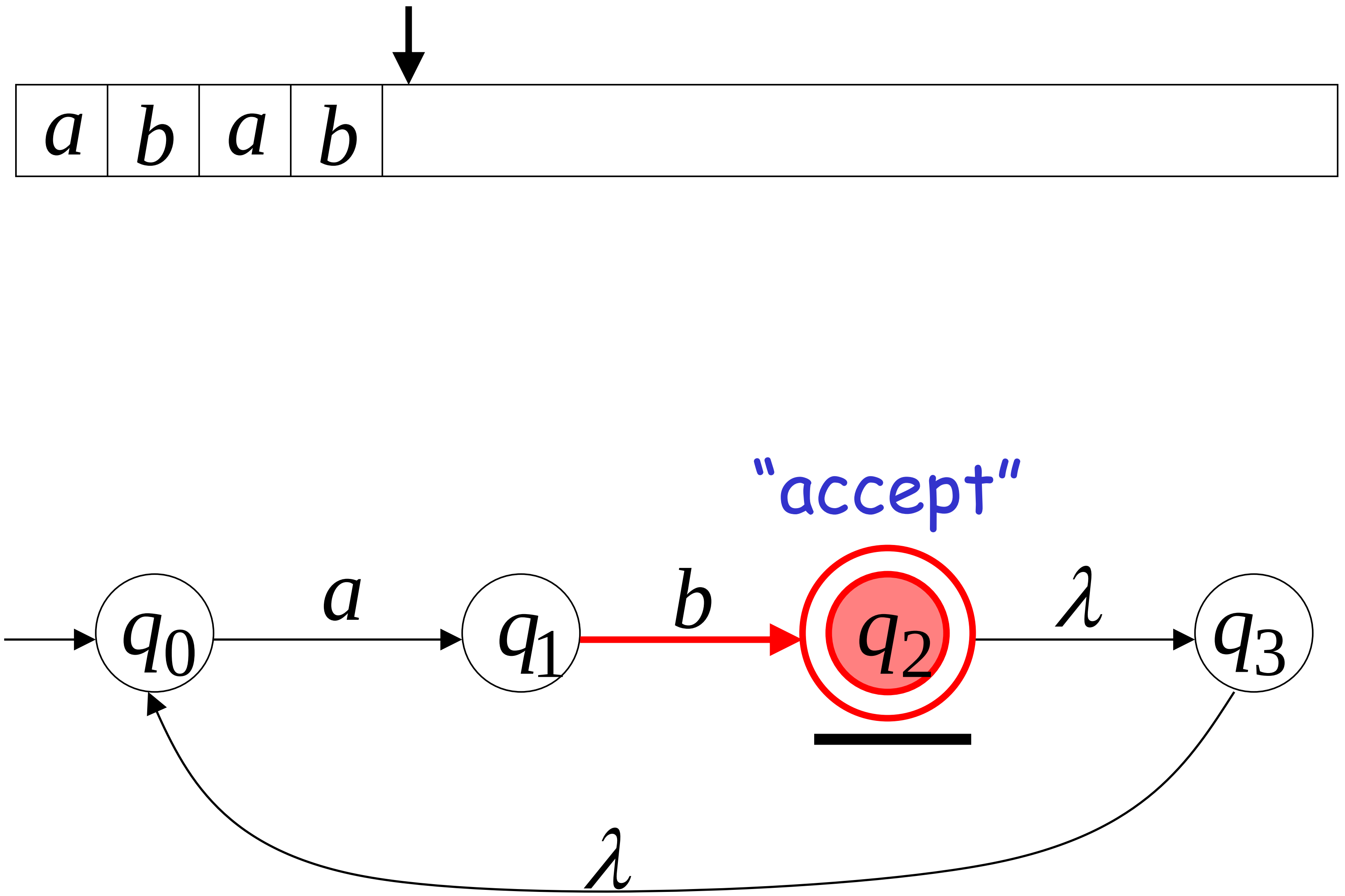

q0
a
q1
b
q2
λ
q3
λ


a
b
a
b


q_0
a
q_1
b
q_2
λ
q_3
λ


a
b
a
b

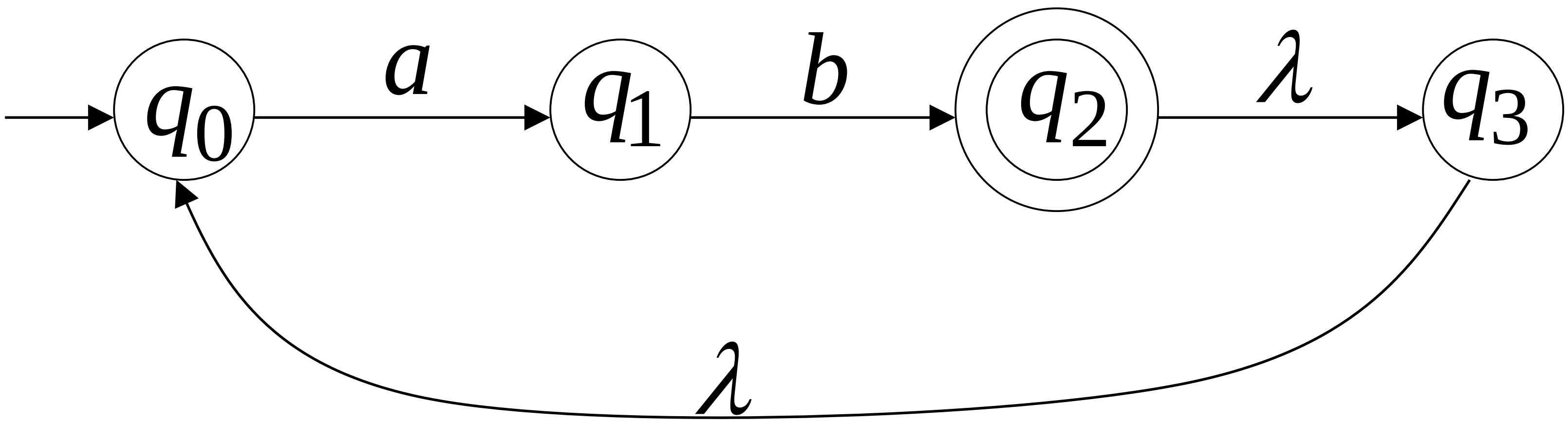

q_0
a
q_1
b
q_2
λ
q_3
λ


a
b
a
b
"accept"
a
b
λ
λ
q_0
q_1
q_2
q_3

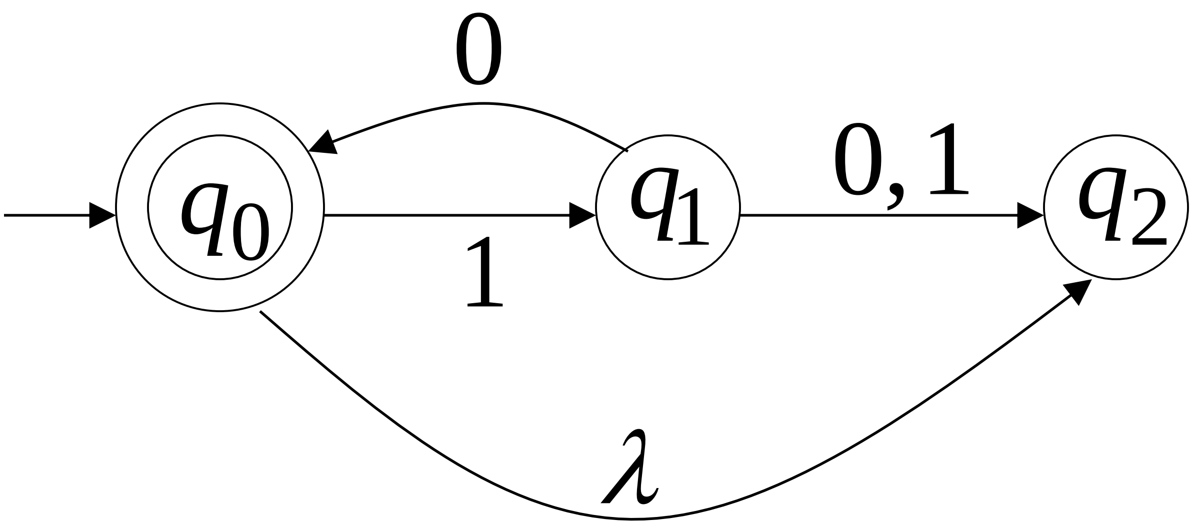
Language accepted

$$L = \{ab,\ abab,\ ababab,\ \ldots\}$$

$$= \{ab\}^+$$



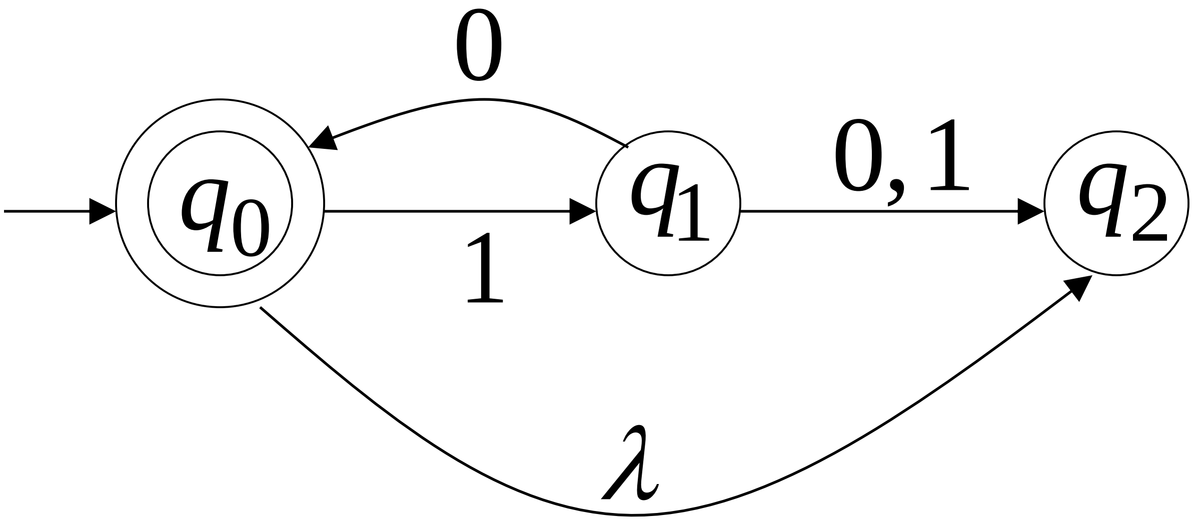
Another NFA Example



Language accepted

$$L(M) = \{\lambda,\ 10,\ 1010,\ 101010,\ \ldots\}$$

$$= \{10\}^*$$

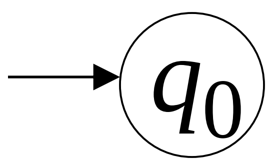


Remarks:

- The λ symbol never appears on the input tape

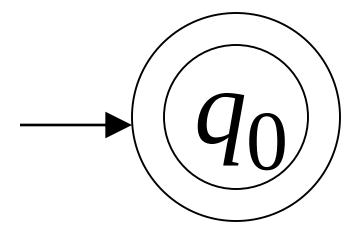
- Extreme automata:

M_1



$$L(M_1) = \{\}$$

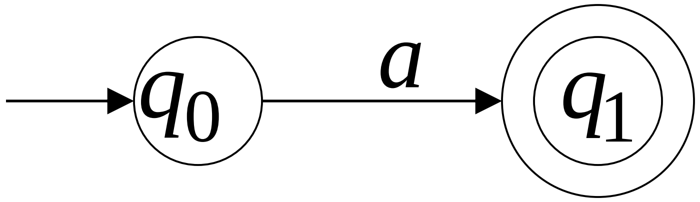
M_2



$$L(M_2) = \{\lambda\}$$

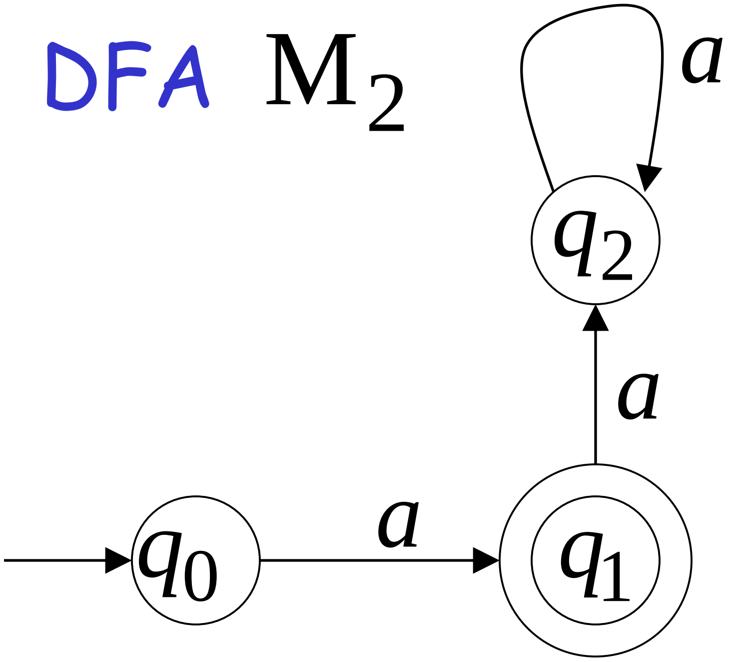
• NFAs are interesting because we can express languages easier than DFAs

NFA M_1



$$L(M_1) = \{a\}$$

DFA M_2



$$L(M_2) = \{a\}$$

Formal Definition of NFAs

$$M = (Q,\ \Sigma,\ \delta,\ q_0,\ F)$$

$Q:$ Set of states, i.e. $\{q_0, q_1, q_2\}$

$\Sigma:$ Input aplhabet, i.e. $\{a, b\}$

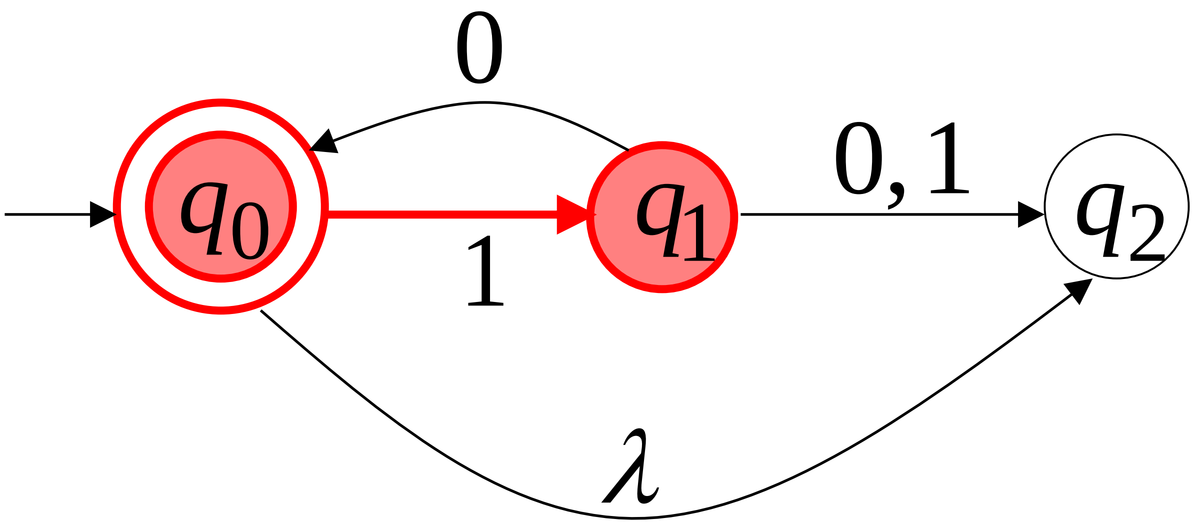
$\delta:$ Transition function

$q_0:$ Initial state

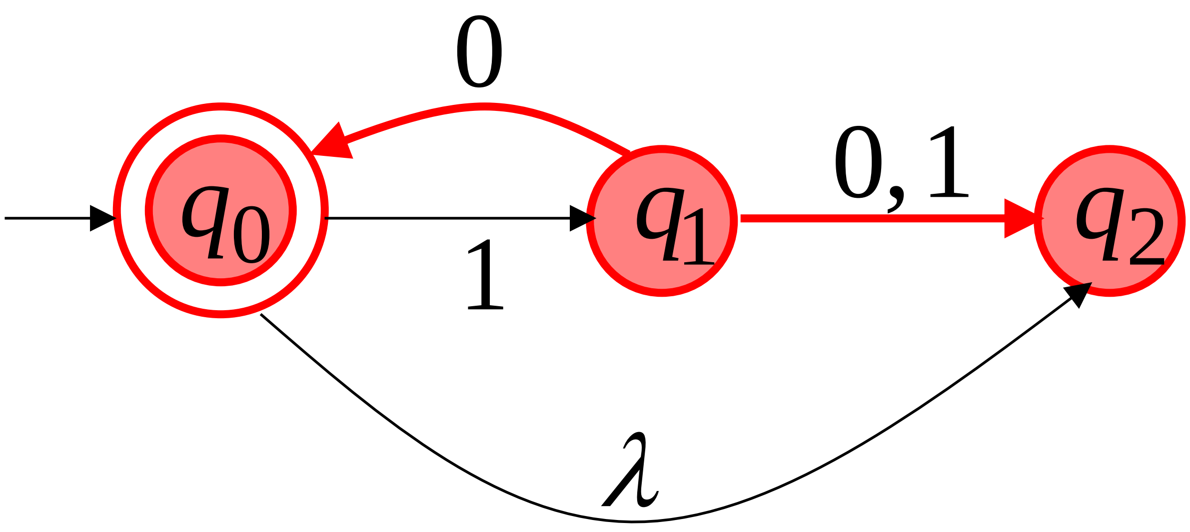
$F:$ Final states

Transition Function δ

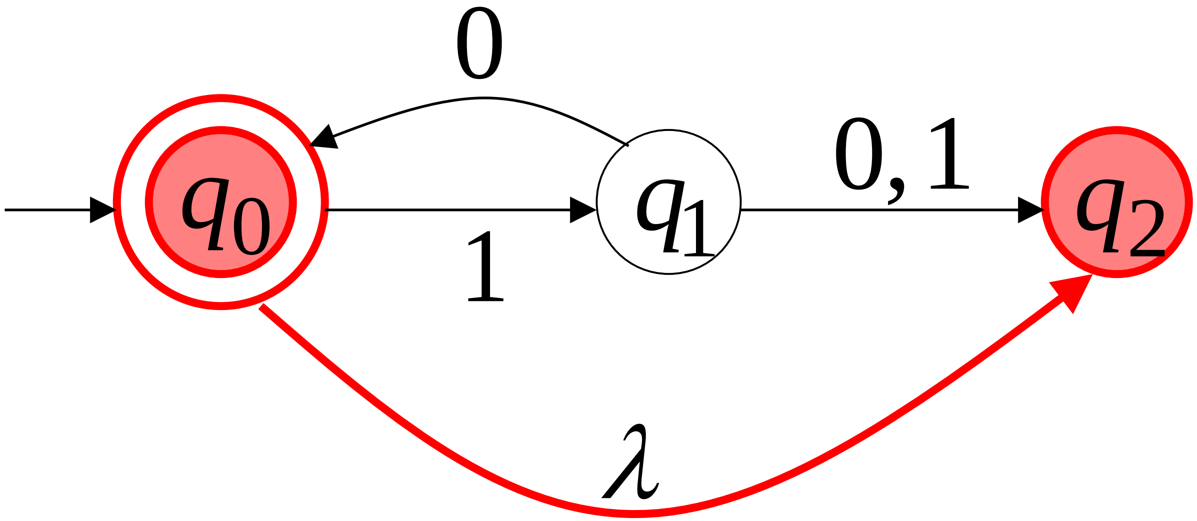
$$\delta(q_0, 1) = \{q_1\}$$



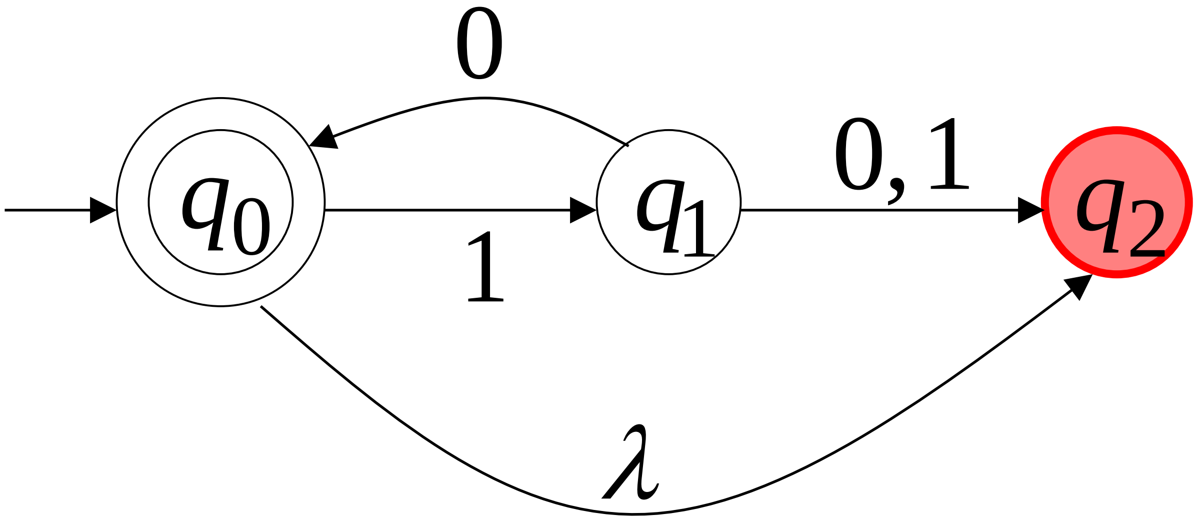
$$\delta(q_1, 0) = \{q_0, q_2\}$$



$$\delta(q_0, \lambda) = \{q_0, q_2\}$$

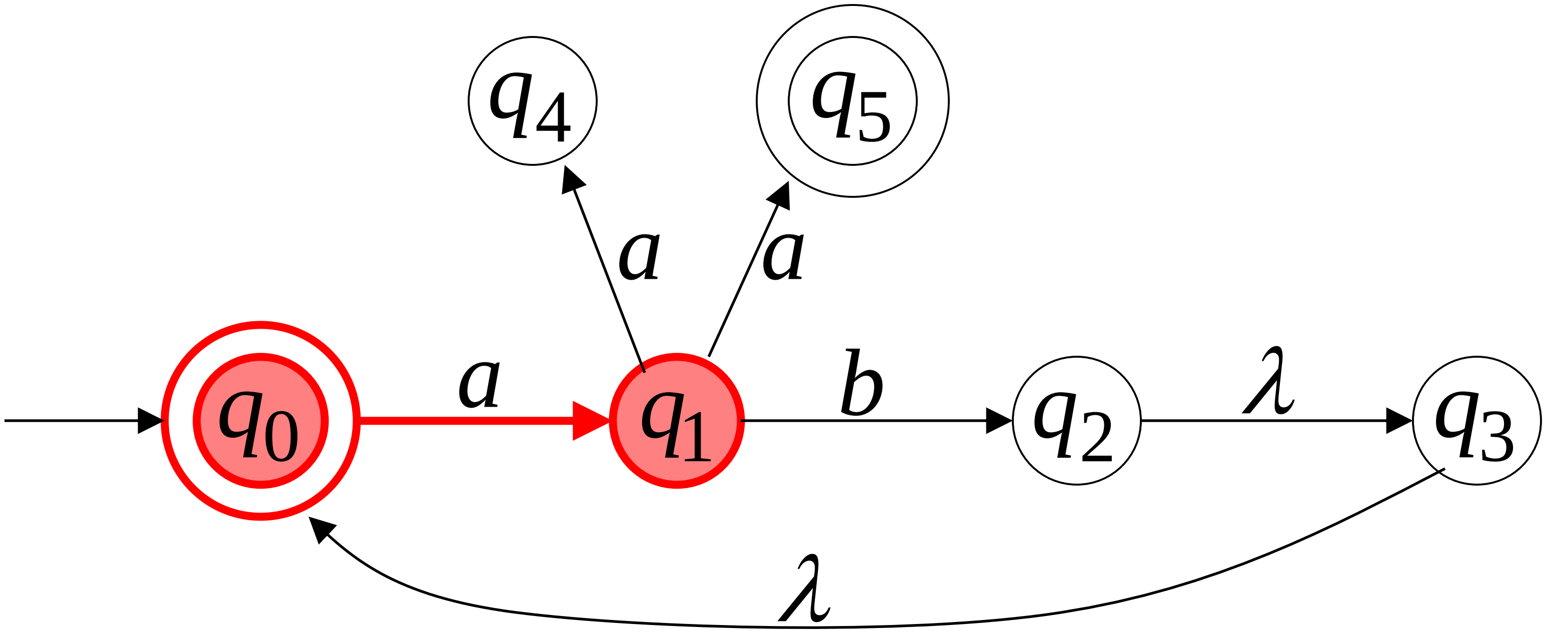


$$\delta(q_2, 1) = \varnothing$$

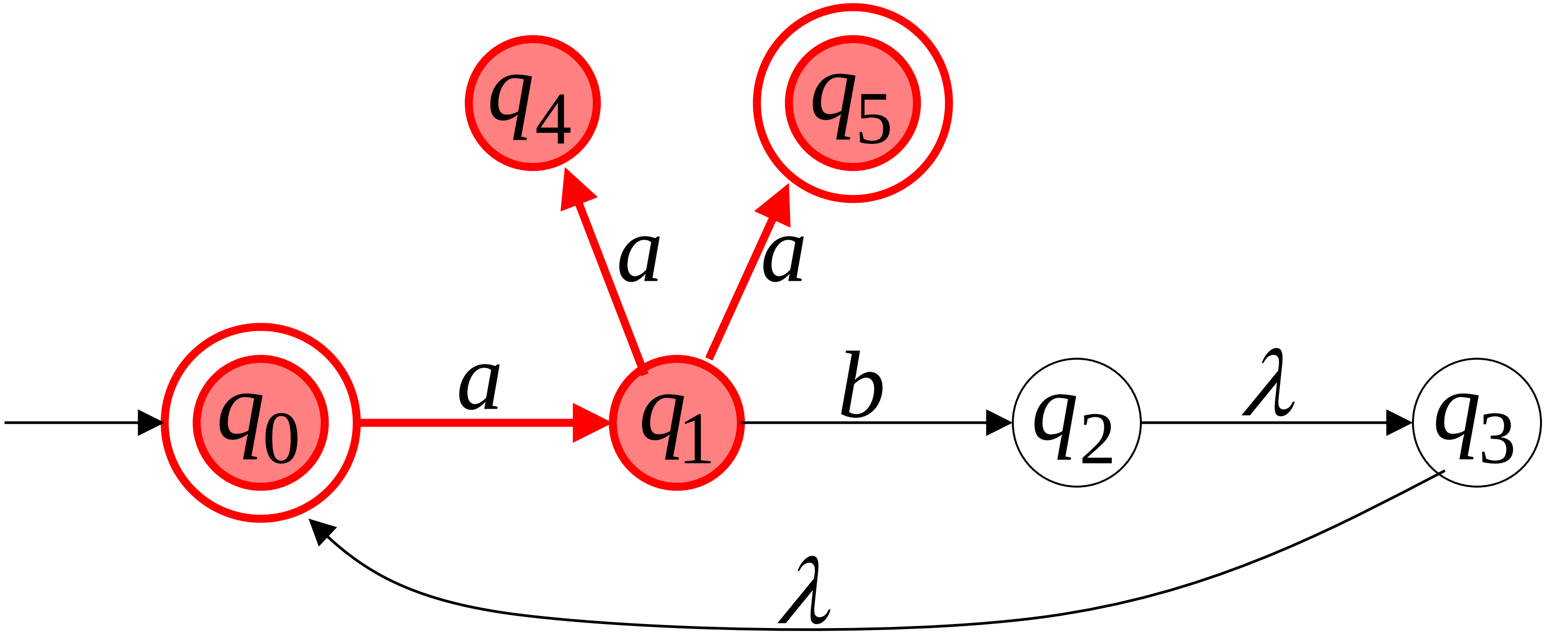


Extended Transition Function δ^*

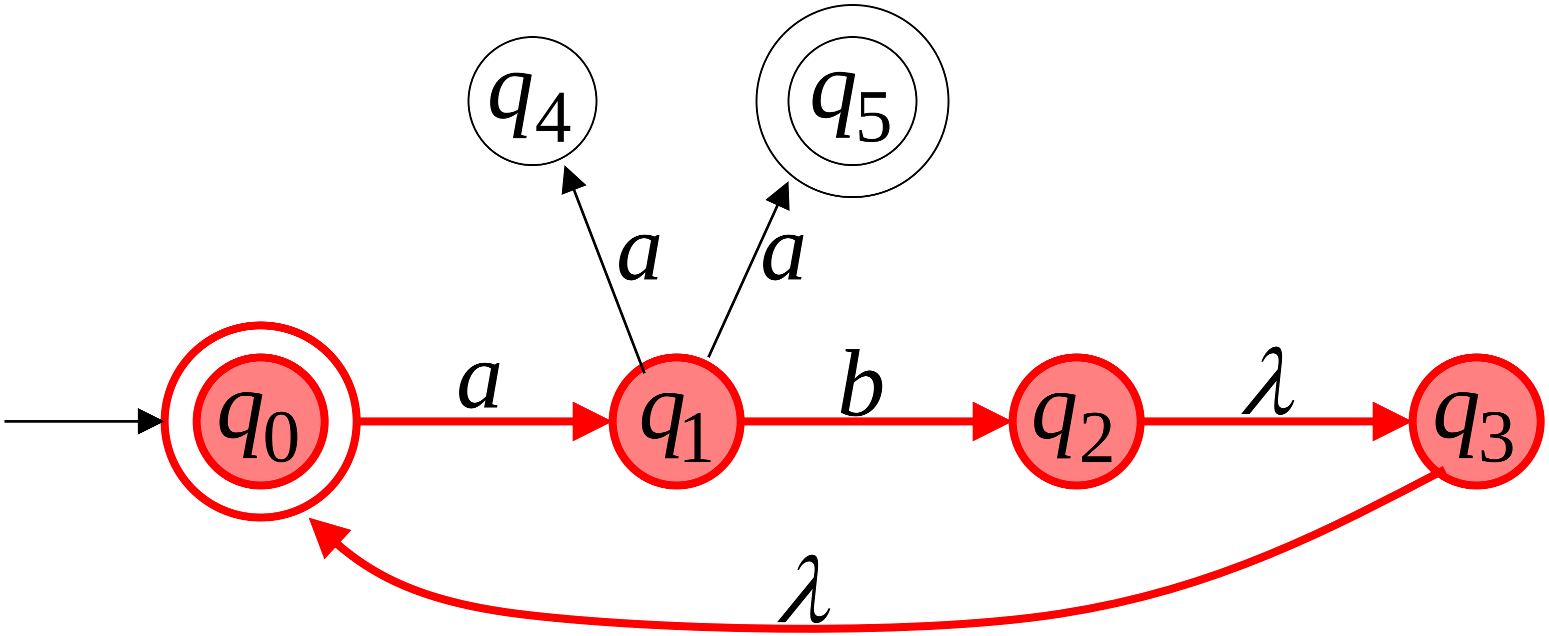
$$\delta^*(q_0, a) = \{q_1\}$$



$$\delta^*(q_0, aa) = \{q_4, q_5\}$$



$$\delta^*(q_0, ab) = \{q_2, q_3, q_0\}$$



Formally

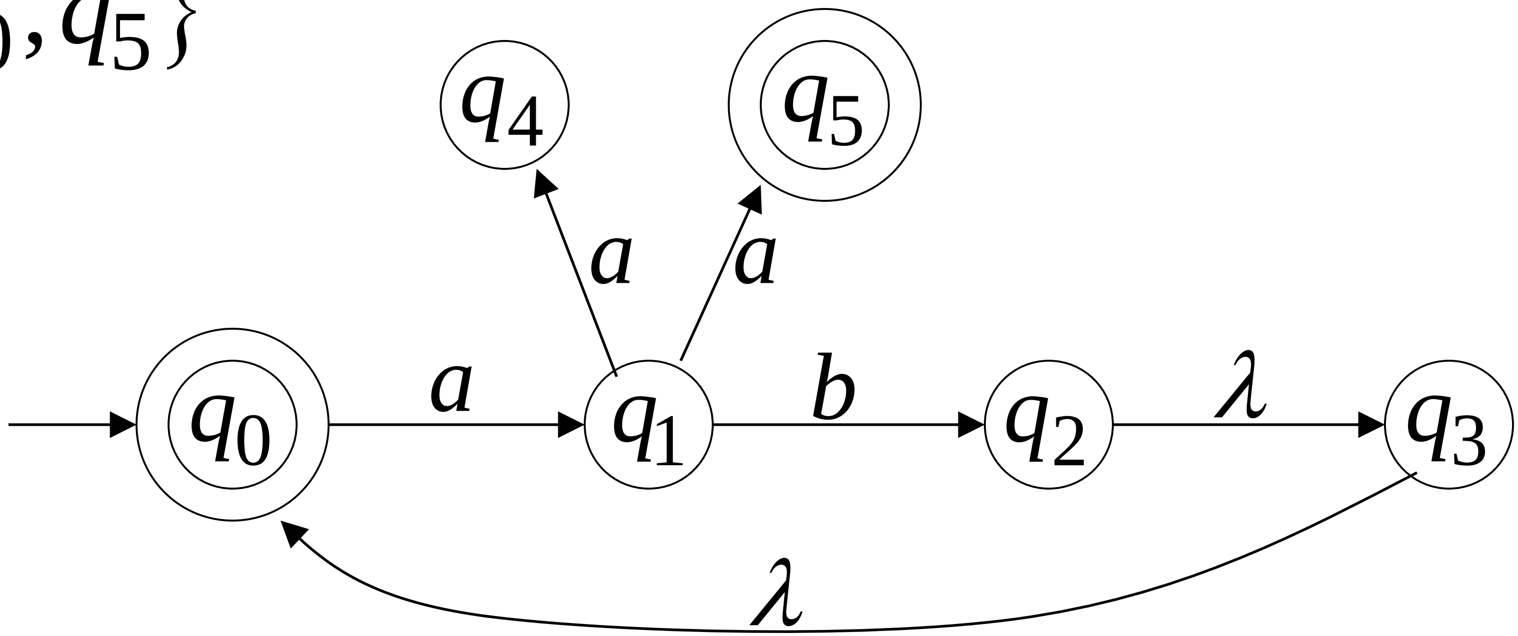
It holds $q_j \in \delta^*(q_i, w)$

if and only if

there is a walk from q_i to q_j with label w

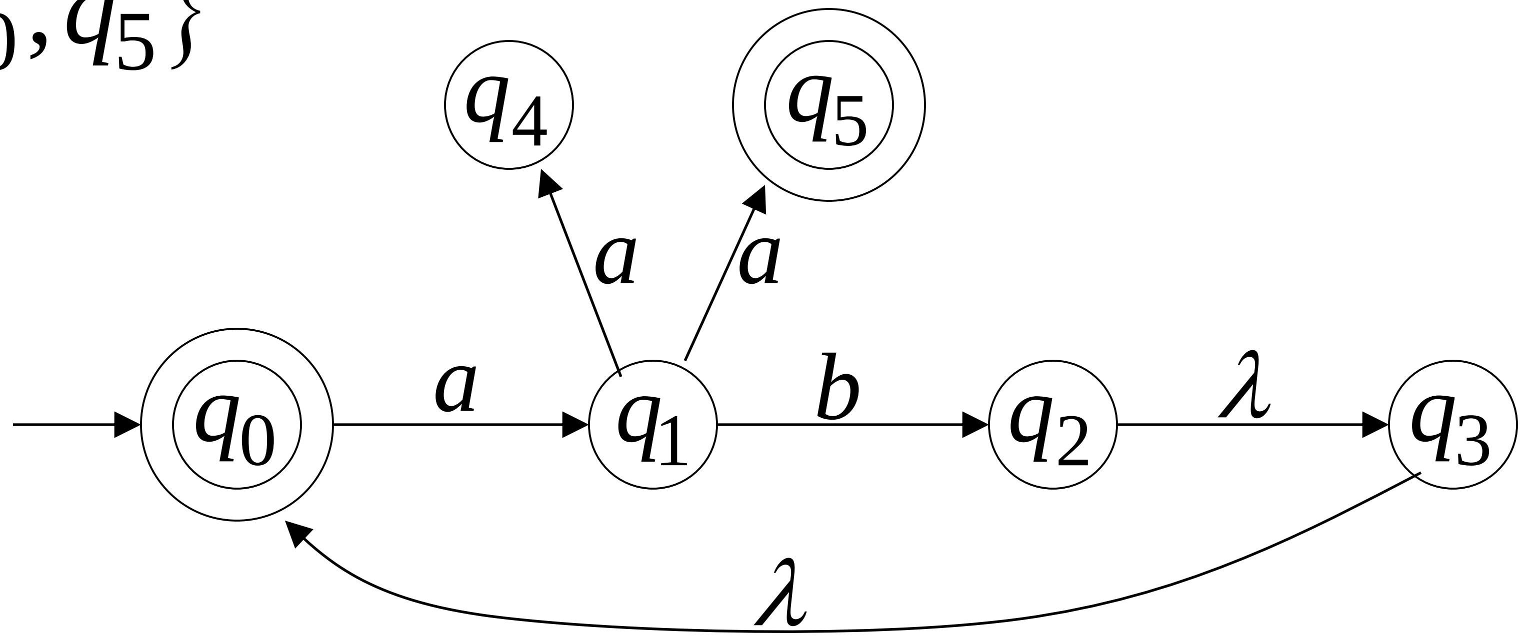
The Language of an NFA M

$F = \{q_0, q_5\}$



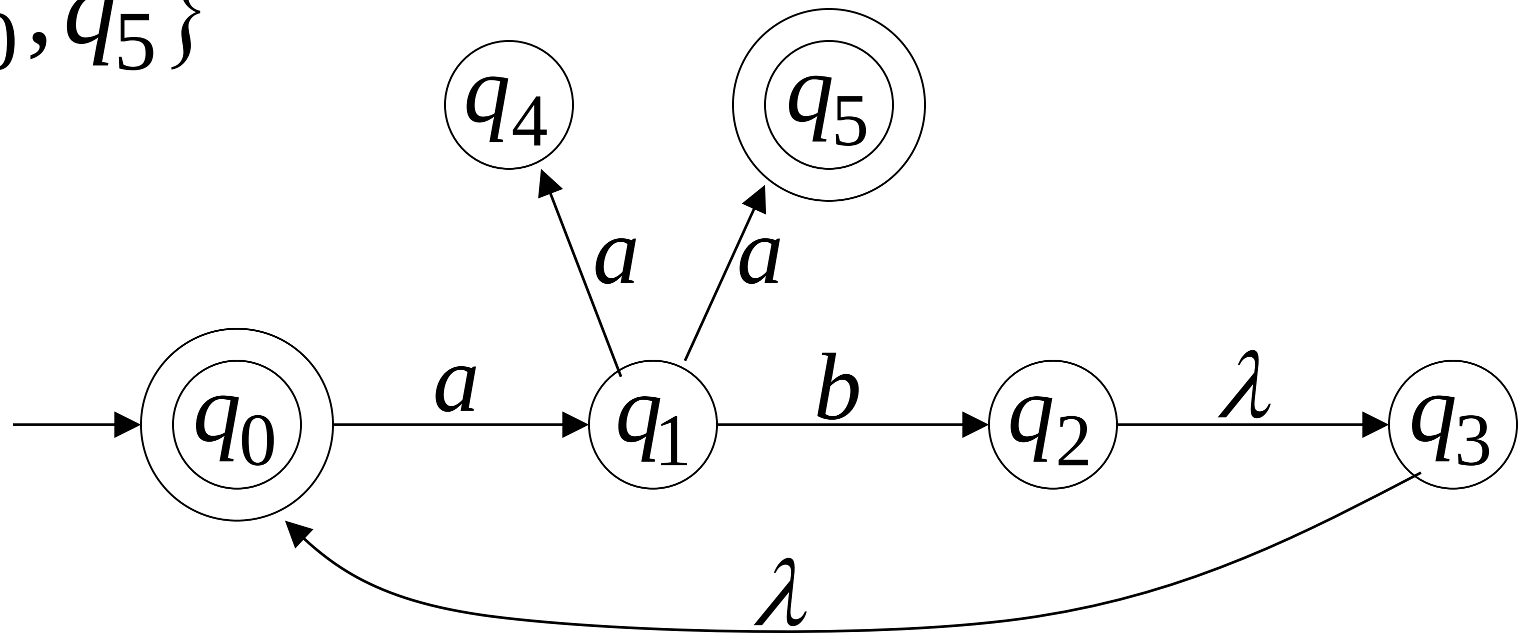
$\delta^*(q_0, aa) = \{q_4, \underline{q_5}\}$ $\qquad$ $aa \in L(M)$

$$F = \{q_0, q_5\}$$



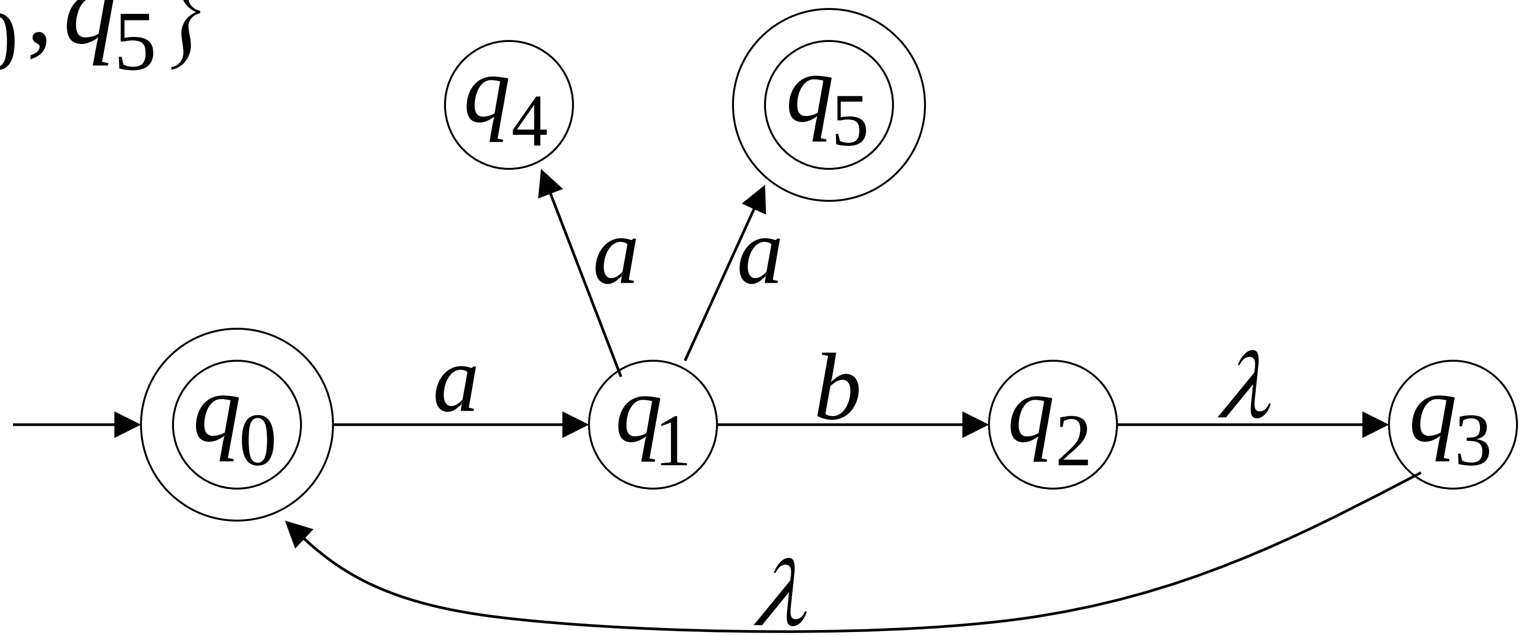
$$\delta^*(q_0, ab) = \{q_2, q_3, \underline{q_0}\} \qquad ab \in L(M)$$

$$F = \{q_0, q_5\}$$

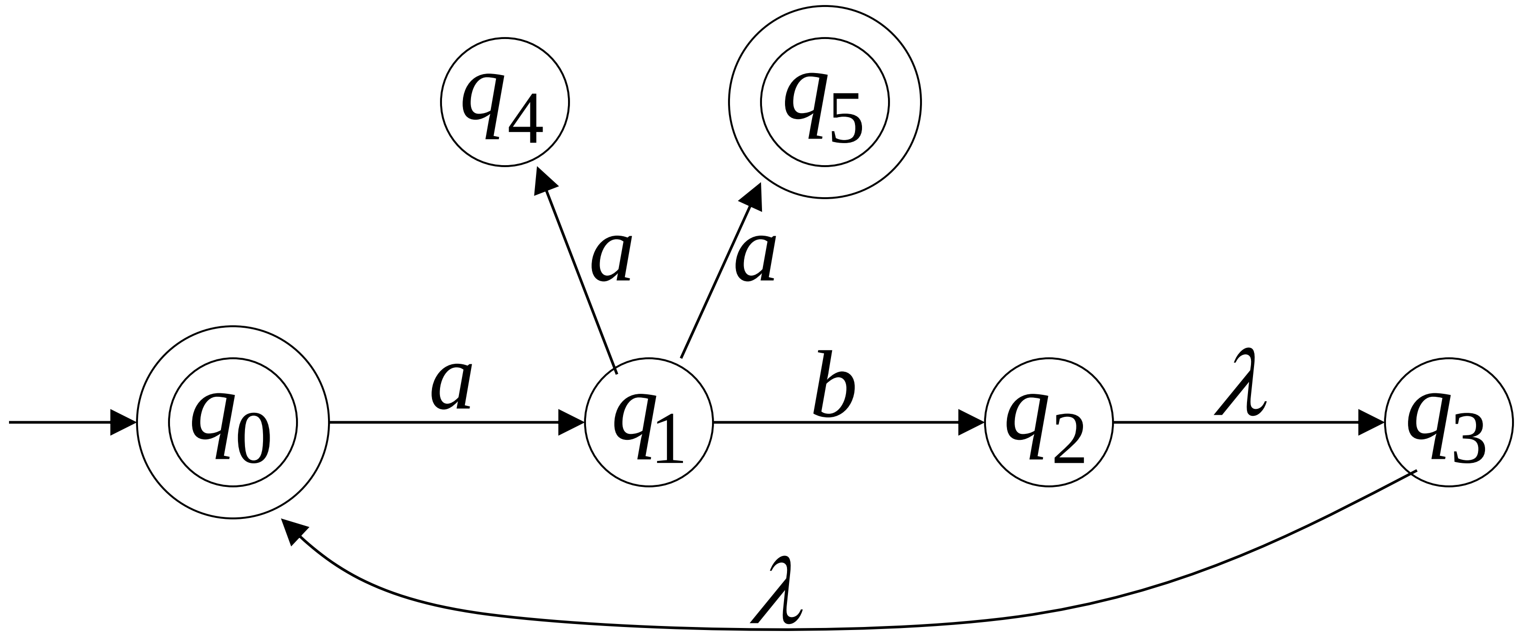


$$\delta^*(q_0, abaa) = \{q_4, \underline{q_5}\} \qquad aaba \in L(M)$$

$$F = \{q_0, q_5\}$$



$$\delta^*(q_0, aba) = \{q_1\} \qquad aba \notin L(M)$$



$$L(M) = \{aa\} \cup \{ab\}^* \cup \{ab\}^+\{aa\}$$

Formally

The language accepted by NFA M is:

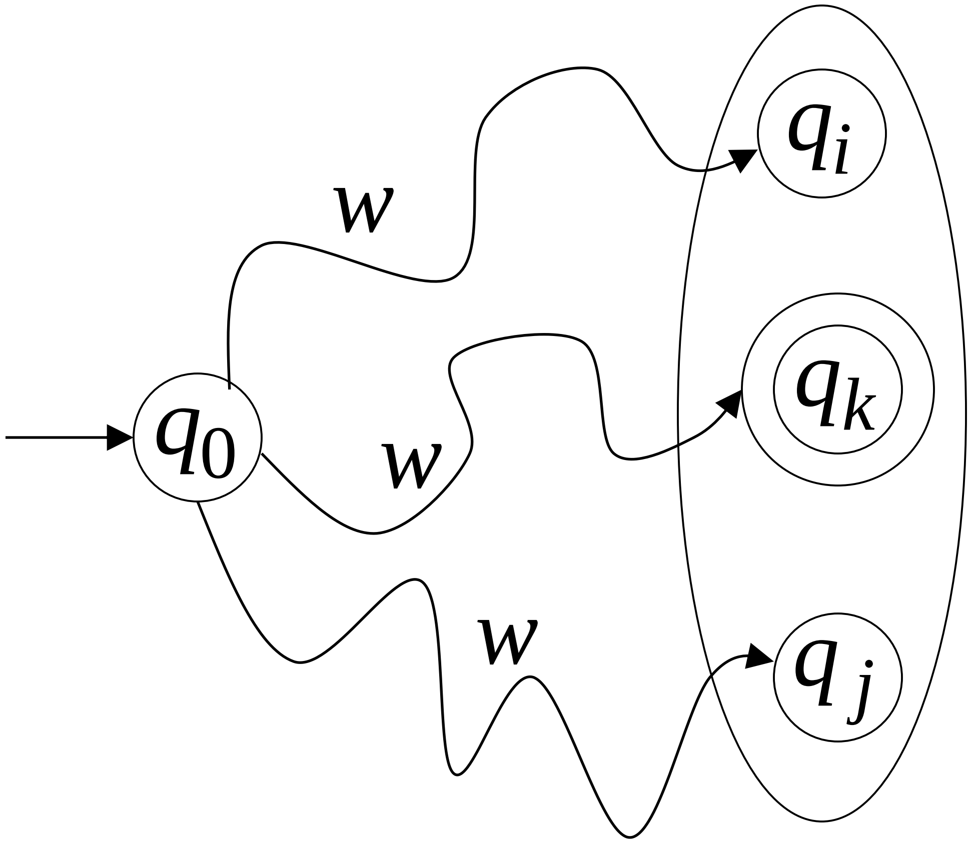
$$L(M) = \{w_1, w_2, w_3, ...\}$$

where $\delta^*(q_0, w_m) = \{q_i, q_j, ...\}$

and there is some $q_k \in F$ (final state)

$$w \in L(M)$$

$$\delta^*(q_0, w)$$



$$q_k \in F$$