

# Formal Languages

## Context-Free Grammars

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21 June 2023

## ① Chomsky Normal Form

*Definition:* A grammar is said to be in CNF if all its productions are in one of two forms:  $A \rightarrow BC$  or  $A \rightarrow a$ .

Linz 6th, §6.2, definition 6.4, page 171

HUM 3rd, §7.1.5, page 272

Kozen, Lecture 21, definition 21.1, page 140

Sudkamp 3rd, §4.5, definition 4.5.1, page 122

## ② Greibach Normal Form

*Definition:* A grammar is said to be in GNF if all its productions are in the form:  $A \rightarrow aX$  where  $a \in T$  and  $X \in V^*$ .

Linz 6th §6.2, definition 6.5, page 174

HMU 3rd, §7.1, page 277

Kozen, Lecture 21, definition 21.1, page 140

Sudkamp 3rd, §4.8, definition 4.8.1, page 131

Hein 4th, §11.7, page 824

NB  $\epsilon \notin L(G)$ .

First we consider Greibach Normal Form, which can be skipped as we do not depend on it later.

### Definition (Greibach Normal Form)

A context-free grammar with start nonterminal  $S$  is said to be in Greibach normal form (GNF) if the productions have one of the following forms:

$$\begin{aligned} S &\rightarrow \epsilon, \\ A &\rightarrow a, \\ A &\rightarrow aX_1 \cdots X_n, \end{aligned}$$

where  $a \in T$  and each  $X_i \in V$  is not equal to  $S$ .

Susan A. Greibach (1965). “A new normal-form theorem for context-free phrase structure grammars”. In: *Journal of the ACM* 12.1, pages 31–42.

“*Abstract.* A context-free phrase structure generator is in *standard form* if and only if all of its rules are of the form:  $Z \rightarrow aY_1, \dots, Y_m$  where  $Z$  and  $Y_i$  are intermediate symbols and  $a$  is a terminal symbol, so that one input symbol is processed at each step.”

Except for the special production  $S \rightarrow \epsilon$ , every production produces some terminal symbol, and so every step of a derivation in a grammar in GNF produces some terminal.

# Terminology

It is *implied* by the term *normal form* that every grammar (and hence every CFL) can be put in Chomsky Normal Form and into Greibach Normal Form.

Indeed, we shall give a construction for the former and state the latter without proof.

*There is another interesting normal form for grammars. This form, called Greibach Normal Form, after Sheila Greibach, has several interesting consequences. Since each use of a production introduces exactly one terminal into a sentential form, a string of length  $n$  has a derivation of exactly  $n$  steps. Also, if we apply the PDA construction to a Greibach-Normal grammar, then we get a PDA with no  $\epsilon$ -rules, thus showing that it is always possible to eliminate such transitions of a PDA.*

HMU 3rd, section 7.1, page 277

## Sheila Adele Greibach (b. 1939)



Sheila Greibach was born in New York City and received the A.B. degree from Radcliffe College in Linguistics and Applied Mathematics *summa cum laude* in 1960. She received the Ph.D. in Applied Mathematics from Harvard University in 1963. She joined the UCLA Faculty in 1969 and the Computer Science Department in 1970 and is now Emeritus Professor.

An important consequence of GNF is that every nonterminal is not left recursive. A nonterminal  $A$  is *left recursive* iff  $A \xRightarrow{*+} A\alpha$  for some  $\alpha \in (\Sigma \cup T)^*$ . Left recursive nonterminals cause top-down deterministic parsers to loop. Every context-free language can be accepted by a (non-deterministic) pushdown automaton that reads a letter from its input every step.



Neither the procedure to convert a grammar to Greibach normal form nor the proof that this can always be done is a simple matter. (Not in HMU either.)

### Theorem (Theorem 6.7, page 176)

*For every context-free grammar  $G$  with  $\epsilon \notin L(G)$ , there exists an equivalent grammar  $\hat{G}$  in Greibach normal form.*

We return to CNF, it is easy to put a grammar in Chomsky Normal Form. And, because of that, there is a very convenient and efficient algorithm to determine membership in any CFL. This algorithm is known as the CYK Algorithm and has many uses, e.g., in natural language processing.

# Chomsky Normal Form

## Definition (Chomsky Normal Form)

A grammar is said to be in Chomsky Normal Form (CNF) if all its productions are in one of two forms:  $A \rightarrow BC$  or  $A \rightarrow a$ .

# Chomsky Normal Form

## Theorem

*Chomsky Normal Form Every context free grammar can be put in Chomsky Normal Form.*

In other words, for every context grammar  $G$ , there is a context free grammar  $G'$  in Chomsky Normal Form such that  $\mathcal{L}(G) = \mathcal{L}(G')$

Linz 6th, §6.2, Theorem 6.6, page 172.

Linz 7th, §6.2, Theorem 6.6, page 182.

HMU 3rd, §7.1.5, page 272.

Martin, §6.6, Theorem 6.6, page 190.

Kozen, Lecture 21, page 140.

## Algorithm to Put a Grammar in CNF

First, eliminate  $\epsilon$  productions and unit productions. One might as well eliminate useless nonterminals (and their productions).

- 1 See that all RHS of length two or more consist only of nonterminals by introducing a nonterminal  $T_a$  and adding the production  $T_a \rightarrow a$  to the grammar for each  $a \in T$ .
- 2 For productions of length  $k$  greater than two, add a cascade of  $k - 2$  new nonterminals and productions.

## Algorithm to Put a Grammar in CNF

For productions of length  $k$  greater than two  $X \rightarrow A_1A_2 \dots A_k$ , add a cascade of  $k - 2$  new nonterminals and productions.

$$\begin{aligned} X &\rightarrow A_1X_1 \\ X_1 &\rightarrow A_2X_2 \\ X_2 &\rightarrow A_3X_3 \\ &\vdots \rightarrow \vdots \\ X_{k-3} &\rightarrow A_{k-2}X_{k-2} \\ X_{k-2} &\rightarrow A_{k-1}A_k \end{aligned}$$

$$X \xRightarrow{1} A_1X_1 \xRightarrow{1} A_1A_2X_2 \xRightarrow{1} A_1A_2A_3X_3 \xRightarrow{*} A_1A_2 \dots A_{k-2}X_{k-2} \xRightarrow{1} A_1A_2 \dots A_{k-2}A_{k-1}A_k$$

## Converting to CNF

Linz, §6.2, example 6.8, page 173. Convert  $S \rightarrow ABa, A \rightarrow aab, B \rightarrow Ac$  to CNF.

$$\begin{array}{ll} S \rightarrow ABT_a & S \rightarrow AV \\ A \rightarrow T_a T_a T_b & V \rightarrow BT_a \\ B \rightarrow AT_c & A \rightarrow T_a W \\ T_a \rightarrow a & W \rightarrow T_a T_b \\ T_b \rightarrow b & B \rightarrow AT_c \\ T_c \rightarrow c & T_a \rightarrow a \\ & T_b \rightarrow b \\ & T_c \rightarrow c \end{array}$$

## Kozen, Example 21.4.

Kozen, example 21.4.  $L = \{a^n b^n \mid n \geq 1\}$



## Kozen, Example 21.5. Balanced parentheses.

Convert the following grammar to CNF.

- 1 Run procedure 'e'; remove  $\epsilon$  productions.
- 2 Add productions for terminals '[' and ']'.  
*(Note: The original image contains a typo ']' which has been corrected to '['.)*
- 3 Factor out the chain of three terminals  $ASB$ .

$$S \rightarrow [S] \mid SS \mid \epsilon$$

$$S \rightarrow [S] \mid [] \mid SS$$

$$S \rightarrow ASB \mid SS \mid AB, \quad A \rightarrow [, \quad B \rightarrow ]$$

$$S \rightarrow AC \mid SS \mid AB, \quad A \rightarrow [, \quad B \rightarrow ], \quad C \rightarrow SB$$

Note that  $S$  is productive even when we eliminate  $S \rightarrow \epsilon$

## Example ???

Put in Chomsky Normal Form:

$$S \rightarrow ASA \mid aB$$

$$A \rightarrow B \mid S$$

$$B \rightarrow b \mid \epsilon$$

Make  $B \rightarrow \epsilon$  redundant

$$S \rightarrow ASA \mid aB \mid a$$

$$A \rightarrow B \mid S \mid \epsilon$$

$$B \rightarrow b \mid \epsilon$$

Make  $A \rightarrow \epsilon$  reduncant

$$S \rightarrow ASA \mid aB \mid a \mid AS \mid SA$$

$$A \rightarrow B \mid S \mid \epsilon$$

$$B \rightarrow b \mid \epsilon$$

Drop  $\epsilon$ -productions.

$$\begin{aligned} S &\rightarrow ASA \mid aB \mid a \mid AS \mid SA \\ A &\rightarrow B \mid S \\ B &\rightarrow b \mid \end{aligned}$$

Make unit productions  $A \rightarrow B$  and  $A \rightarrow S$  redundant

$$\begin{aligned} S &\rightarrow ASA \mid aB \mid a \mid AS \mid SA \\ S &\rightarrow BSB \mid aB \mid a \mid BS \mid SB \\ S &\rightarrow SSS \mid aB \mid a \mid SS \mid SS \\ A &\rightarrow B \mid S \\ B &\rightarrow b \end{aligned}$$

Drop unit productions.

$$\begin{aligned}S &\rightarrow ASA \mid aB \mid a \mid AS \mid SA \\S &\rightarrow BSB \mid aB \mid a \mid BS \mid SB \\S &\rightarrow SSS \mid aB \mid a \mid SS \\B &\rightarrow b\end{aligned}$$

A is now unproductive. Drop RHSs with A.

$$\begin{aligned}S &\rightarrow aB \mid a \\S &\rightarrow BSB \mid aB \mid a \mid BS \mid SB \\S &\rightarrow SSS \mid aB \mid a \mid SS \\B &\rightarrow b\end{aligned}$$

Drop duplicates.

$$\begin{aligned}S &\rightarrow aB \mid a \\S &\rightarrow BSB \mid BS \mid SB \\S &\rightarrow SSS \mid SS \\B &\rightarrow b\end{aligned}$$

Continuing ... The grammar has been cleaned-up:  $\epsilon$ -productions and unit productions are gone. Now complete the conversions to CNF by making the all productions have terminals or non-terminals, but not both, on the RHS.

$$\begin{aligned} S &\rightarrow T_a B \mid a \\ S &\rightarrow BSB \mid BS \mid SB \\ S &\rightarrow SSS \mid SS \\ B &\rightarrow b \\ T_a &\rightarrow a \end{aligned}$$

Finish by shortening sequences of non-terminals.

$$\begin{aligned} S &\rightarrow T_a B \mid a \\ S &\rightarrow BV \mid BS \mid SB \\ S &\rightarrow SW \mid SS \\ B &\rightarrow b \\ T_a &\rightarrow a \\ V &\rightarrow SB \\ W &\rightarrow SS \end{aligned}$$

# Exhaustive Search

Linz 6th, Section 5.2, page 141.

*Algorithm.* Given a grammar  $G$  without  $\epsilon$ -productions and unit productions and a string  $w$ , systematically construct all possible leftmost (or rightmost) derivations and see if the sentential forms are consistent with  $w$ .

It is easy (but exhausting) to try each of productions and all of the sentential forms derived so far and test the result with the given input string  $w$ . But how do we know when to stop and declare that the string is *not* derivable?

# Exhaustive Search

Lemma. Given a grammar  $G$  without  $\epsilon$ -productions. If  $S \xRightarrow{*}_G \alpha \xRightarrow{*}_G x$ , then the  $|\alpha| \leq |x|$ . Proof. For all nonterminals  $N$ , if  $N \xRightarrow{*}_G x$ , then  $1 \leq |x|$ .

But the length  $\alpha \xRightarrow{*}_G \beta$ , may be such that  $|\alpha| = |\beta|$  For example, on productions  $A \rightarrow B$  or  $A \rightarrow a$ .

So a derivation could go on and on, without getting any longer

# Exhaustive Search

## Lemma

Given a grammar  $G$  without  $\epsilon$ -productions and unit productions. Let  $\#_+(\alpha)$  be the number of terminal symbols in  $\alpha$  plus the length of  $\alpha$ . If  $\beta \xrightarrow{1}_G \gamma$ , then  $\#_+(\beta) < \#_+(\gamma)$ . And, hence, if  $\beta \xrightarrow{*}_G \gamma$ , then  $\#_+(\beta) < \#_+(\gamma)$ .

For example:  $\#(S) = 1$  and  $\#_+(w) = 2|W|$  for all  $w$  in  $\Sigma^*$ . Also,  $xAy \xrightarrow{1}_G xay$  has  $\#_+(xAy) = 2 + 3 = 5 < 6 = 3 + 3 = \#_+(xay)$ . And,  $xAy \xrightarrow{1}_G xaBy$  has  $\#_+(xAy) = 2 + 3 = 5 < 7 = 3 + 4 = \#_+(xaBy)$ .

## Corollary

If  $S \xrightarrow{*}_G w$  is a derivation, then the number of steps is bounded by  $2|w|$ .

Total search cost:  $2|w|$  stages with outdegree  $|P|$  equals  $|P|^{2|w|}$ .



## Example of Exhaustive Search

Linz 6th, Section 5.2, Example 5.7, page 141.

Busch's notes, [class09](#), page 12ff.

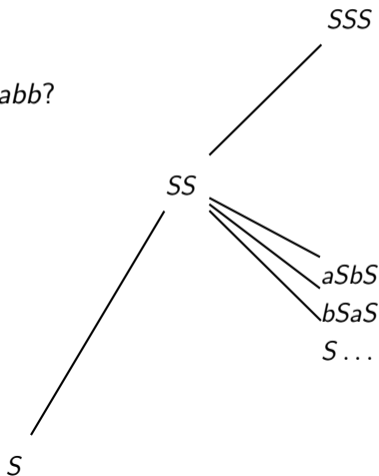
Related Linz 7th, Section 5.2, Example 5.8, page 151.

Find a derivation of  $aabb$  in the grammar  $S \rightarrow SS \mid aSb \mid bSa \mid \epsilon$ .

Find a derivation of  $aabb$  in the grammar  $S \rightarrow SS \mid aSb \mid bSa \mid ab \mid ba$ .

- 1  $S \rightarrow SS$
- 2  $S \rightarrow aSb$
- 3  $S \rightarrow bSa$
- 4  $S \rightarrow ab$
- 5  $S \rightarrow ba$

Does  $S$  derive  $aabb$ ?



$SSSS$

$aSbSS$

$bSaSS$

$SS$

$aSSbS$

$aaSbSbS$

$abSaSbS$

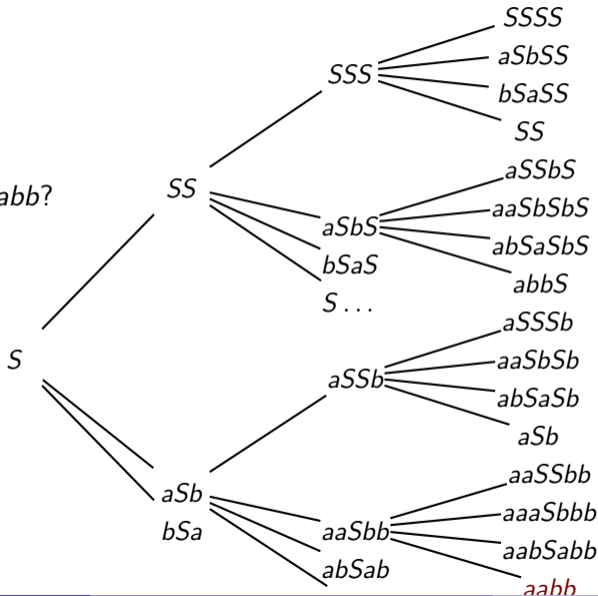
$abbS$

$aSSSb$

now expand

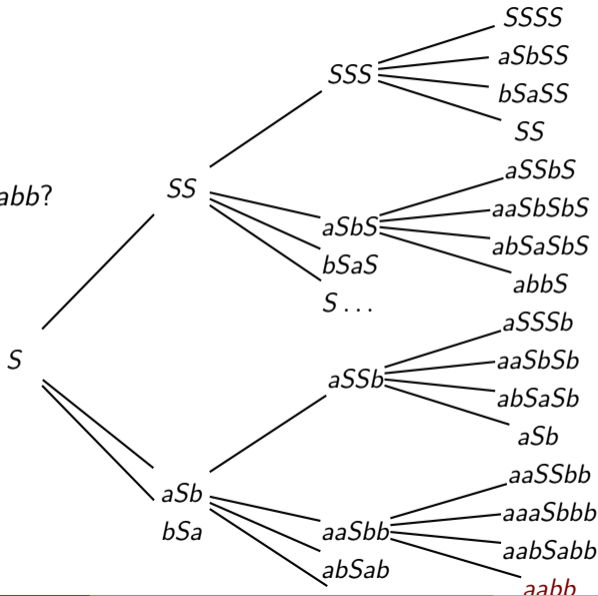
- 1  $S \rightarrow SS$
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Does  $S$  derive  $aabb$ ?



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- 5  $S \rightarrow ba$

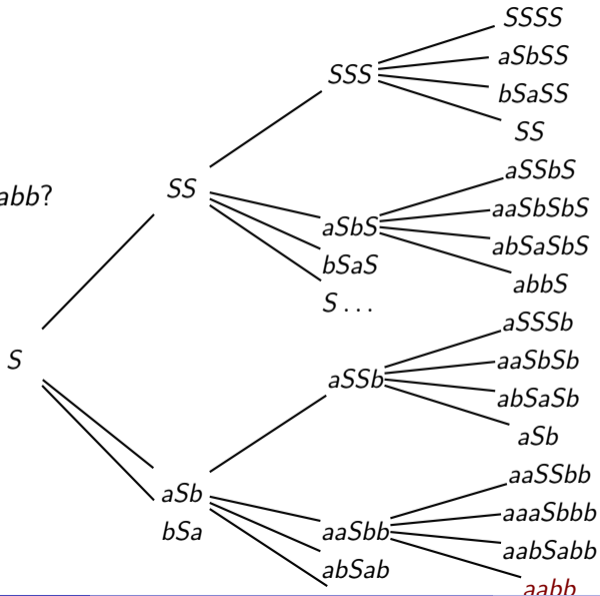
Does  $S$  derive  $aabb$ ?



now expand

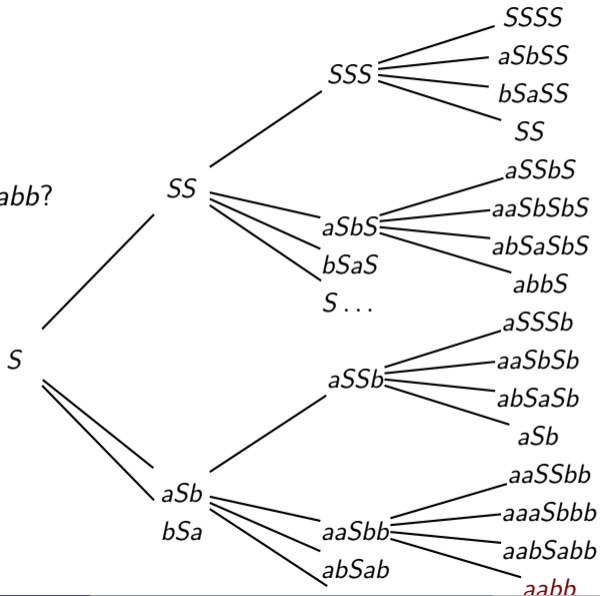
- 1  $S \rightarrow SS$
- 2  $S \rightarrow aSb$
- 3  $S \rightarrow bSa$
- 4  $S \rightarrow ab$
- 5  $S \rightarrow ba$

Does  $S$  derive  $aabb$ ?



- 1  $S \rightarrow SS$
- 2  $S \rightarrow aSb$
- 3  $S \rightarrow bSa$
- 4  $S \rightarrow ab$
- 5  $S \rightarrow ba$

Does  $S$  derive  $aabb$ ?



now expand

# Earley's Algorithm

Earley parser [↗](#) at Wiki

# Earley's Algorithm

Another general algorithm for parsing strings in a given context-free language was invented by Jay Earley.

It uses so-called items,  $N \rightarrow \alpha \bullet \beta$ , to mark the progress of the parsing. Information about partial parses is stored in an array,



*Earley's algorithm and CYK both run in time  $O(n^3)$  on general grammars. However, Earley's algorithm does much better on grammars of practical interest. Although we will not prove it, Earley's algorithm runs in time  $O(n^2)$  on unambiguous grammars and in time  $O(n)$  on LR(1) grammars. (We will not define LR(1) grammars, but they are equivalent to DCFLs. Furthermore, the syntax of almost every programming language is essentially an LR(1) grammar. For more information, see a textbook on compiler design.) In contrast, the CYK algorithm always takes time bounded above and below by multiples of  $n^3$ , regardless of the grammar.*

Floyd, Section 5.12, page 392–393.

Donald Ervin Knuth (Dec. 1965). “On the Translation of Languages from Left to Right”. In: *Information and Control* 8.6, pages 607–639

### Theorem

*Every DCFL has a LR(1) grammar.*

(Aho, Ullman 1972, exercise 5.2.26, page 398.)

### Theorem

*Every language accepted by an LR(1) is a DCFL.*

LL(k) grammars introduced by Stearns and Lewis (1969)  
Information and Control

Aho, Ullman 1972, exercise 5.2.25, page 398. Every LL(k) grammar is a LR(k) grammar. Hence, the LL(1) languages are a subset of the LR(1) languages.

### Theorem

*The LL(1) languages are a proper subset of the LR(1) languages*

## Cocke-Younger-Kasami (CYK) Algorithm

Linz 6th, section 6.3, page 178

Busch's notes, [class10](#), page 21ff

HMU 3rd, section 7.4.4, page 303

Sudkamp 3rd, Section 4.6: The CYK Algorithm; pages 124–128

Kozen, Lecture 27, page 191

Floyd, Section 5.11, Figure 5.20, page 390

[cyk](#) at Wikipedia

John Cocke and Jacob T. Schwartz (1970). *Programming Languages and their Compilers*. Technical report. New York: NYU, Courant Institute

T. Kasami (1965). *An efficient recognition and syntax analysis algorithm for context-free languages*. Technical report AFCRL-65-758. Air Force Cambridge Res. Lab., Bedford Mass.

D. H. Younger (1967). "Recognition and Parsing of Context-Free Languages in  $n^3$ ". In: *Information and Control* 10, pages 189–208

# CYK Algorithm

Let  $G = \langle T, V, P, S \rangle$  be a CFG in Chomsky normal form. The questions:

Does  
 $S \xRightarrow{*} w$

is transformed to many questions:

For which  $A \in V$  does  
 $A \xRightarrow{*} w[i : j]$   
for each substring  $w[i : j]$

# CYK Algorithm

Let  $G = \langle T, V, P, S \rangle$  be a CFG in Chomsky normal form.

We can determine if a string  $s$  is in  $L(G)$ . Suppose  $s = a_0a_1, \dots, a_{n-1}$ . We write  $s[i : j]$  where  $0 \leq i \leq j \leq n - 1$  for the substring of  $s$  of length  $j - i + 1$  beginning at position  $i$  and ending at position  $j$ .

A classic, dynamic programming algorithm ensues. We require a two-dimensional, triangular array  $M$  containing sets of nonterminals. The nonterminal  $A$  is in the cell  $A \in M[i, j]$  iff  $A \Rightarrow^* s[i : j]$  for  $0 \leq i \leq j < n$ .

# CYK Algorithm

$A \in M[i, j]$  iff  $A \Rightarrow^* s[i : j]$ ; starting/ending index

	0	1	2	...	j	...	n-1
0	$s[0 : 0]$	$s[0 : 1]$	$s[0 : 2]$		$s[0 : j]$		$s[0 : n - 1]$
1		$s[1 : 1]$	$s[1 : 2]$		$s[1 : j]$		$s[1 : n - 1]$
2			$s[2 : 2]$		$s[2 : j]$		$s[2 : n - 1]$
⋮							
i							$s[i : n - 1]$
⋮							
n-1							$s[n - 1 : n - 1]$

There are different ways of organizing the matrix, making it essential to get the indices correct.

They are all equivalent, but the one that works best conceptionally is having the rows belong to substrings of increasing lengths.

So we code the algorithm that way.

On paper, label the cells with the appropriate substring to keep the meaning of the cell straight.



# CYK Algorithm

$A \in M[l, i]$  iff  $A \Rightarrow^* s[l : i + l]$ ; length/starting index

$A \in M[j, k - j]$  iff  $A \Rightarrow^* s[j : k]$ ; length/starting index

	0	1	2	...	i	...	n-1
0	$s[0 : 0]$	$s[1 : 1]$	$s[2 : 2]$		$s[i : i]$		$s[n - 1 : n - 1]$
1	$s[0 : 1]$	$s[1 : 2]$	$s[2 : 3]$		$s[i : i + 1]$		
2	$s[0 : 2]$	$s[1 : 3]$	$s[2 : 4]$		$s[i : i + 2]$		
⋮							
l	$s[0 : l]$	$s[1 : l + 1]$	$s[2 : l + 2]$		$s[i : i + l]$		
⋮							
n-1	$s[0 : n - 1]$						

# Key Recurrence

A nonterminal  $N \Rightarrow^* s[r : r + g]$  iff for some  $r \leq m \leq r + g$

- ①  $N \rightarrow AB$  is a production of the grammar, and
- ②  $A \Rightarrow^* s[r : m]$ , and
- ③  $B \Rightarrow^* s[m + 1 : r + g]$ .

## CYK Algorithm – Initialization

```
-- Given the string  $s = a_0a_1 \cdots a_{n-1}$ 
for  $i$  in  $0, \dots, n-1$  loop
   $M[i, i] := \emptyset$ 
  for  $N \rightarrow a$  in  $P$  loop
    add  $N$  to  $M[i, i]$  if  $a = a_i$  --  $N \Rightarrow_G^* s[i : i]$ 
  end loop
end loop
```

# CYK Algorithm

```
for g in 1, ..., n-1 loop -- every substring of length 2, 3, ..., n
  for r in 0, ..., n-g-1 loop -- substring starting at index r
    M[r, r+g] := ∅
    for m in r, ..., r+g-1 loop -- every midpoint of substring
      -- s[r:r+g] = s[r:m] ++ s[m+1:r+g]
      -- the set of nonterminals generating s[r:m]
      L = M[r, m]
      -- the set of nonterminals generating s[m+1:r+g]
      R = M[m+1, r+g]
      for A → BC in P loop -- some productions
        -- add A if A ⇒G* s[r:r+g]
        add A to M[r, r+g] if B ∈ L and C ∈ R
      end loop
    end loop
  end loop
end loop
return S ∈ M[0, n-1]
```

```

for  $i := 0$  to  $n - 1$  do                                /* strings of length 1 first */
  begin
     $T_{i,i+1} := \emptyset$ ;                                /* initialize to  $\emptyset$  */
    for  $A \rightarrow a$  a production of  $G$  do
      if  $a = x_{i,i+1}$  then  $T_{i,i+1} := T_{i,i+1} \cup \{A\}$ 
    end;
  for  $m := 2$  to  $n$  do                                /* for each length  $m \geq 2$  */
    for  $i := 0$  to  $n - m$  do                            /* for each substring */
      begin                                              /* of length  $m$  */
         $T_{i,i+m} := \emptyset$ ;                            /* initialize to  $\emptyset$  */
        for  $j := i + 1$  to  $i + m - 1$  do            /* for all ways to break */
          for  $A \rightarrow BC$  a production of  $G$  do    /* up the string */
            if  $B \in T_{i,j} \wedge C \in T_{j,i+m}$ 
              then  $T_{i,i+m} := T_{i,i+m} \cup \{A\}$ 
          end;
        end;
      end;
    end;
  end;

```

Kozen  
1997

```

 $n := |s|$ ; (* initialization *)
for every variable  $X$  do begin
  for  $i := 1$  to  $n$  do
    for  $k := i$  to  $n$  do
       $T[i, k, X] := \text{false}$ ;
  for  $i := 1$  to  $n$  do
    if  $X \rightarrow s_{ij}$  is a production then
       $T[i, i, X] := \text{true}$ ;
end;

for  $k := 2$  to  $n$  do
  for  $i := k - 1$  down to  $1$  do
    for all productions of the form  $X \rightarrow YZ$  do
      for  $j := i$  to  $k - 1$  do
        if  $T[i, j, Y]$  and  $T[j + 1, k, Z]$  then
           $T[i, k, X] := \text{true}$ ;

```

Floyd & Beigel  
1994

**FIGURE 5.20:** The CYK algorithm. The string  $s$  belongs to  $L(G)$  if and only if  $T[1, n, S] = \text{true}$ , where  $n = |s|$  and  $S$  is  $G$ 's start variable.



Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$S \rightarrow AB$$

$$A \rightarrow BB \mid a$$

$$B \rightarrow AB \mid b$$



Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

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Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned} S &\rightarrow AB \\ A &\rightarrow BB \mid a \\ B &\rightarrow AB \mid b \end{aligned}$$

$a$	$a$	$b$	$b$	$b$
$aa$	$ab$	$bb$	$bb$	
$aab$	$abb$	$bbb$		
$aabb$	$abbb$			
$aabbb$				

$$A \xRightarrow{1} a$$

Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned}
 S &\rightarrow AB \\
 A &\rightarrow BB \mid a \\
 B &\rightarrow AB \mid b
 \end{aligned}$$

$a$	$a$	$b$	$b$	$b$
$\{A\}$				
$aa$	$ab$	$bb$	$bb$	
$aab$	$abb$	$bbb$		
$aabb$	$abbb$			
$aabbb$				

$$A \xrightarrow{1} a$$

Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned}
 S &\rightarrow AB \\
 A &\rightarrow BB \mid a \\
 B &\rightarrow AB \mid b
 \end{aligned}$$

$a$	$a$	$b$	$b$	$b$
$\{A\}$	$\{A\}$			
$aa$	$ab$	$bb$	$bb$	
$aab$	$abb$	$bbb$		
$aabb$	$abbb$			
$aabbb$				

$$B \xRightarrow{1} b$$

Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned}
 S &\rightarrow AB \\
 A &\rightarrow BB \mid a \\
 B &\rightarrow AB \mid b
 \end{aligned}$$

$a$	$a$	$b$	$b$	$b$
$\{A\}$	$\{A\}$	$\{B\}$		
$aa$	$ab$	$bb$	$bb$	
$aab$	$abb$	$bbb$		
$aabb$	$abbb$			
$aabbb$				

$$B \xRightarrow{1} b$$

Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned}
 S &\rightarrow AB \\
 A &\rightarrow BB \mid a \\
 B &\rightarrow AB \mid b
 \end{aligned}$$

$a$	$a$	$b$	$b$	$b$
$\{A\}$	$\{A\}$	$\{B\}$	$\{B\}$	
$aa$	$ab$	$bb$	$bb$	
$aab$	$abb$	$bbb$		
$aabb$	$abbb$			
$aabbb$				

$$B \stackrel{1}{\Rightarrow} b$$

Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned}
 S &\rightarrow AB \\
 A &\rightarrow BB \mid a \\
 B &\rightarrow AB \mid b
 \end{aligned}$$

$a$	$a$	$b$	$b$	$b$
$\{A\}$	$\{A\}$	$\{B\}$	$\{B\}$	$\{B\}$
$aa$	$ab$	$bb$	$bb$	
$aab$	$abb$	$bbb$		
$aabb$	$abbb$			
$aabbb$				

Nothing derives  $aa$ .

Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned} S &\rightarrow AB \\ A &\rightarrow BB \mid a \\ B &\rightarrow AB \mid b \end{aligned}$$

$a$	$a$	$b$	$b$	$b$
$\{A\}$	$\{A\}$	$\{B\}$	$\{B\}$	$\{B\}$
$aa$	$ab$	$bb$	$bb$	
$\emptyset$				
$aab$	$abb$	$bbb$		
$aabb$	$abbb$			
$aabbb$				

$$S \xrightarrow{1} AB \xrightarrow{*} ab \text{ and also } B \xrightarrow{1} AB \xrightarrow{*} ab.$$



Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned} S &\rightarrow AB \\ A &\rightarrow BB \mid a \\ B &\rightarrow AB \mid b \end{aligned}$$

	$a$	$a$	$b$	$b$	$b$
	$\{A\}$	$\{A\}$	$\{B\}$	$\{B\}$	$\{B\}$
	$aa$	$ab$	$bb$	$bb$	
	$\emptyset$	$\{S, B\}$			
	$aab$	$abb$	$bbb$		
	$aabb$	$abbb$			
	$aabbb$				

$$A \xrightarrow{1} BB \xrightarrow{*} bb.$$

Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned}
 S &\rightarrow AB \\
 A &\rightarrow BB \mid a \\
 B &\rightarrow AB \mid b
 \end{aligned}$$

$a$	$a$	$b$	$b$	$b$
$\{A\}$	$\{A\}$	$\{B\}$	$\{B\}$	$\{B\}$
$aa$	$ab$	$bb$	$bb$	
$\emptyset$	$\{S, B\}$	$\{A\}$		
$aab$	$abb$	$bbb$		
$aabb$	$abbb$			
				$aabbb$

$$A \xrightarrow{1} BB \xrightarrow{*} bb.$$

Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned} S &\rightarrow AB \\ A &\rightarrow BB \mid a \\ B &\rightarrow AB \mid b \end{aligned}$$

	$a$	$a$	$b$	$b$	$b$
	$\{A\}$	$\{A\}$	$\{B\}$	$\{B\}$	$\{B\}$
	$aa$	$ab$	$bb$	$bb$	
	$\emptyset$	$\{S, B\}$	$\{A\}$	$\{A\}$	
	$aab$	$abb$	$bbb$		
	$aabb$	$abbb$			
	$aabbb$				

$$S \xrightarrow{1} AB \xrightarrow{*} a \cdot ab \text{ and also } B \xrightarrow{1} AB \xrightarrow{*} a \cdot ab.$$

Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned}
 S &\rightarrow AB \\
 A &\rightarrow BB \mid a \\
 B &\rightarrow AB \mid b
 \end{aligned}$$

	$a$	$a$	$b$	$b$	$b$
	$\{A\}$	$\{A\}$	$\{B\}$	$\{B\}$	$\{B\}$
	$aa$	$ab$	$bb$	$bb$	
	$\emptyset$	$\{S, B\}$	$\{A\}$	$\{A\}$	
	$aab$	$abb$	$bbb$		
	$\{S, B\}$				
	$aabb$	$abbb$			
	$aabbb$				

$$A \xrightarrow{1} BB \xrightarrow{*} ab \cdot b.$$

Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned} S &\rightarrow AB \\ A &\rightarrow BB \mid a \\ B &\rightarrow AB \mid b \end{aligned}$$

$a$	$a$	$b$	$b$	$b$
$\{A\}$	$\{A\}$	$\{B\}$	$\{B\}$	$\{B\}$
$aa$	$ab$	$bb$	$bb$	
$\emptyset$	$\{S, B\}$	$\{A\}$	$\{A\}$	
$aab$	$abb$	$bbb$		
$\{S, B\}$	$\{A\}$			
$aabb$	$abbb$			
$aabbb$				

$$S \xrightarrow{1} AB \xrightarrow{*} bb \cdot b \text{ and also } B \xrightarrow{1} AB \xrightarrow{*} bb \cdot b.$$



Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned} S &\rightarrow AB \\ A &\rightarrow BB \mid a \\ B &\rightarrow AB \mid b \end{aligned}$$

$a$	$a$	$b$	$b$	$b$
$\{A\}$	$\{A\}$	$\{B\}$	$\{B\}$	$\{B\}$
$aa$	$ab$	$bb$	$bb$	
$\emptyset$	$\{S, B\}$	$\{A\}$	$\{A\}$	
$aab$	$abb$	$bbb$		
$\{S, B\}$	$\{A\}$	$\{S, B\}$		
$aabb$	$abbb$			
$\{A\}$				
$aabbb$				

$S \xrightarrow{1} AB \xrightarrow{*} a \cdot bbb$  or  $abb \cdot b$ , and also  $B \xrightarrow{1} AB \xrightarrow{*} a \cdot bbb$  or  $abb \cdot b$ .

Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned} S &\rightarrow AB \\ A &\rightarrow BB \mid a \\ B &\rightarrow AB \mid b \end{aligned}$$

$a$	$a$	$b$	$b$	$b$
$\{A\}$	$\{A\}$	$\{B\}$	$\{B\}$	$\{B\}$
$aa$	$ab$	$bb$	$bb$	
$\emptyset$	$\{S, B\}$	$\{A\}$	$\{A\}$	
$aab$	$abb$	$bbb$		
$\{S, B\}$	$\{A\}$	$\{S, B\}$		
$aabb$	$abbb$			
$\{A\}$	$\{S, B\}$			
$aabbb$				

$$S \xrightarrow{1} AB \xrightarrow{*} a \cdot abbb \text{ or } aabb \cdot b, \text{ and also } B \xrightarrow{1} AB \xrightarrow{*} a \cdot abbb \text{ or } aabb \cdot b$$



Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned}
 S &\rightarrow AB \\
 A &\rightarrow BB \mid a \\
 B &\rightarrow AB \mid b
 \end{aligned}$$

$a$	$a$	$b$	$b$	$b$
$\{A\}$	$\{A\}$	$\{B\}$	$\{B\}$	$\{B\}$
$aa$	$ab$	$bb$	$bb$	
$\emptyset$	$\{S, B\}$	$\{A\}$	$\{A\}$	
$aab$	$abb$	$bbb$		
$\{S, B\}$	$\{A\}$	$\{S, B\}$		
$aabb$	$abbb$			
$\{A\}$	$\{S, B\}$			
$aabbb$				
$\{S, B\}$				

Linz 6th, §6.3, example 6.11, page 179. Use the CYK method to determine if the string  $w = aabbb$  is in the language generated by the following grammar  $G$  in Chomsky Normal Form.

$$\begin{aligned} S &\rightarrow AB \\ A &\rightarrow BB \mid a \\ B &\rightarrow AB \mid b \end{aligned}$$

	$a$	$a$	$b$	$b$	$b$
	$\{A\}$	$\{A\}$	$\{B\}$	$\{B\}$	$\{B\}$
	$aa$	$ab$	$bb$	$bb$	
	$\emptyset$	$\{S, B\}$	$\{A\}$	$\{A\}$	
	$aab$	$abb$	$bbb$		
	$\{S, B\}$	$\{A\}$	$\{S, B\}$		
	$aabb$	$abbb$			
	$\{A\}$	$\{S, B\}$			
	$aabbb$				
	$\{S, B\}$				

Since  $S$  is found in the cell labeled  $aabbb$ ,  $S \xRightarrow{*} w$  and the string  $w = aabbb$  is in the language  $L(G)$ .

## Example

Linz 6th, §6.3, exercise 4, page 180 (solution page 421). Use the CYK method to determine if the string  $w = aaabbbb$  is in the language generated by the grammar  $S \rightarrow aSb \mid b$ .

First, convert the grammar to CNF yielding the following:

$$S \rightarrow AC \mid b$$

$$C \rightarrow SB$$

$$A \rightarrow a$$

$$B \rightarrow b$$

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

*a a a b b b b*

*aa aa ab bb bb bb*

$S \rightarrow AC \mid b$

*aaa aab abb bbb bbb*

$C \rightarrow SB$

$A \rightarrow a$

*aaab aabb abbb bbbb*

$B \rightarrow b$

*aaabb aabbb abbbb*

*aaabbb aabbbb*

*aaabbbb*

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

*a*      *a*      *a*      *b*      *b*      *b*      *b*  
*A*

*aa*      *aa*      *ab*      *bb*      *bb*      *bb*

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

*aaa*      *aab*      *abb*      *bbb*      *bbb*

*aaab*      *aabb*      *abbb*      *bbbb*

*aaabb*      *aabbb*      *abbbb*

$A \xRightarrow{*} a$

*aaabbb*      *aabbbb*

*aaabbbb*

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

*a*      *a*      *a*      *b*      *b*      *b*      *b*  
*A*      *A*  
*aa*      *aa*      *ab*      *bb*      *bb*      *bb*

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

*aaa*      *aab*      *abb*      *bbb*      *bbb*

*aaab*      *aabb*      *abbb*      *bbbb*

*aaabb*      *aabbb*      *abbbb*

$A \xRightarrow{*} a$

*aaabbb*      *aabbbb*

*aaabbbb*

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>				
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>
------------	------------	------------	------------	------------

<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>
-------------	-------------	-------------	-------------

<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>
--------------	--------------	--------------

$A \xRightarrow{*} a$

<i>aaabbb</i>	<i>aabbbb</i>
---------------	---------------

*aaabbbb*

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>			
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>
------------	------------	------------	------------	------------

<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>
-------------	-------------	-------------	-------------

<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>
--------------	--------------	--------------

$S \xRightarrow{*} b$  and

$B \xRightarrow{*} b$

<i>aaabbb</i>	<i>aabbbb</i>
---------------	---------------

*aaabbbb*



Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>		
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>
------------	------------	------------	------------	------------

<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>
-------------	-------------	-------------	-------------

*aaabb aabbb abbbb*

$S \xRightarrow{*} b$  and

$B \xRightarrow{*} b$

*aaabbb aabbbb*

*aaabbbb*

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>
------------	------------	------------	------------	------------

<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>
-------------	-------------	-------------	-------------

<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>
--------------	--------------	--------------

$S \xRightarrow{*} b$  and

$B \xRightarrow{*} b$

<i>aaabbb</i>	<i>aabbbb</i>
---------------	---------------

*aaabbbb*

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>
------------	------------	------------	------------	------------

<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>
-------------	-------------	-------------	-------------

<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>
--------------	--------------	--------------

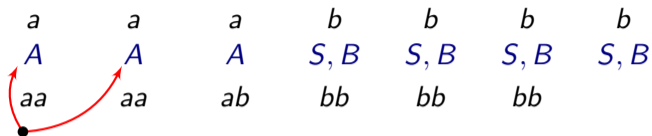
$S \xRightarrow{*} b$  and

$B \xRightarrow{*} b$

<i>aaabbb</i>	<i>aabbbb</i>
---------------	---------------

*aaabbbb*

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)



$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

aaa aab abb bbb bbb

aaab aabb abbb bbbb

aaabb aabbb abbbb

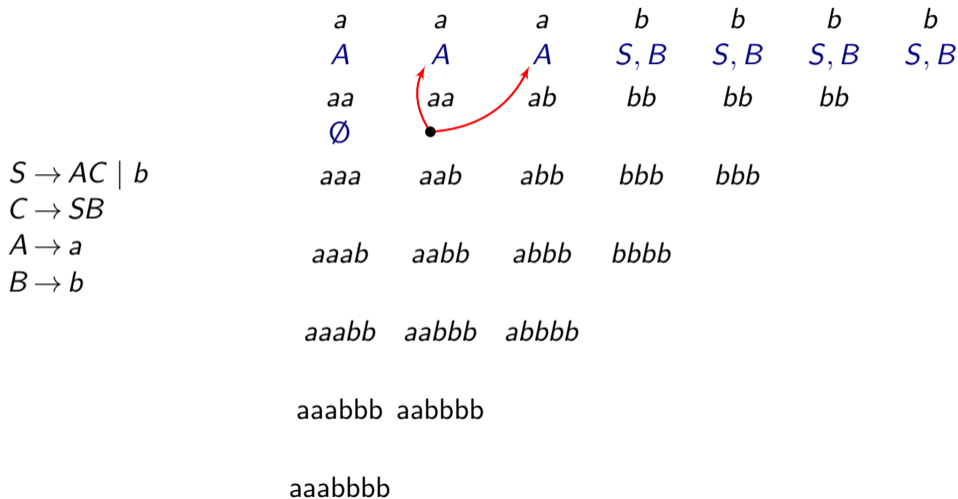
aaabbb aabbbb

aaabbbb

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$						
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$							
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$							
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)



Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$					
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$							
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$							
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

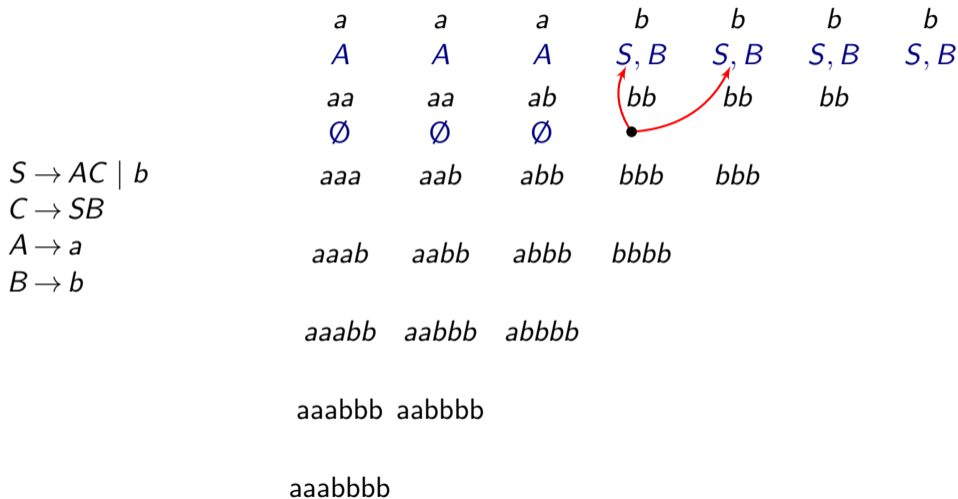
	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	•				
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$							
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$							
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						



Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$				
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$							
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$							
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)



Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>			
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$							
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$							
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	•		
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$							
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$							
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>		
<i>S</i> → <i>AC</i>   <i>b</i>	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
<i>C</i> → <i>SB</i>							
<i>A</i> → <i>a</i>	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
<i>B</i> → <i>b</i>							
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

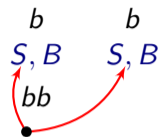
<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	<i>C</i>
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>	<i>bbb</i>	<i>bbb</i>

*aaab aabb abbb bbbb*

*aaabb aabbb abbbb*

*aaabbb aabbbb*

*aaabbbb*



Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$							
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$							
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						



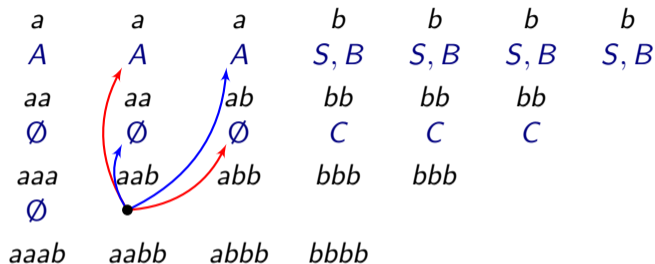


Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$	$\emptyset$						
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$							
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$



*aaabb* *aaabbb* *abbbb*

*aaabbb* *aaabbbb*

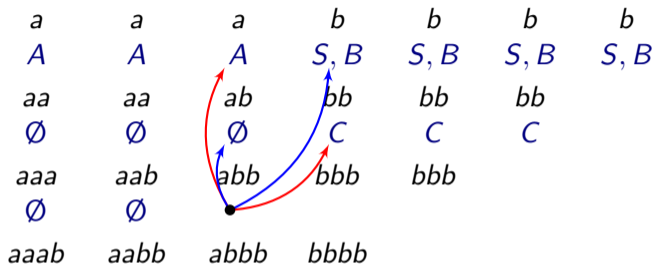
*aaabbbb*

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$	$\emptyset$	$\emptyset$					
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$							
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$



*aaabb* *aabbb* *abbbb*

*aaabbb* *aabbbb*

*aaabbbb*

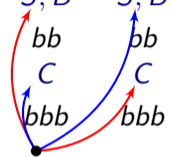
Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>S</i> → <i>AC</i>   <i>b</i>	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
<i>C</i> → <i>SB</i>	$\emptyset$	$\emptyset$	<i>S</i>				
<i>A</i> → <i>a</i>	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
<i>B</i> → <i>b</i>							
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$	$\emptyset$	<i>S</i>				
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			



*aaabb*   *aabbb*   *abbbb*

*aaabbb*   *aabbbb*

*aaabbbb*

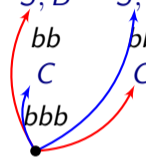
Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>S</i> → <i>AC</i>   <i>b</i>	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
<i>C</i> → <i>SB</i>	$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$			
<i>A</i> → <i>a</i>	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
<i>B</i> → <i>b</i>							
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$			
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			



*aaabb*   *aabbb*   *abbbb*

*aaabbb*   *aabbbb*

*aaabbbb*



Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

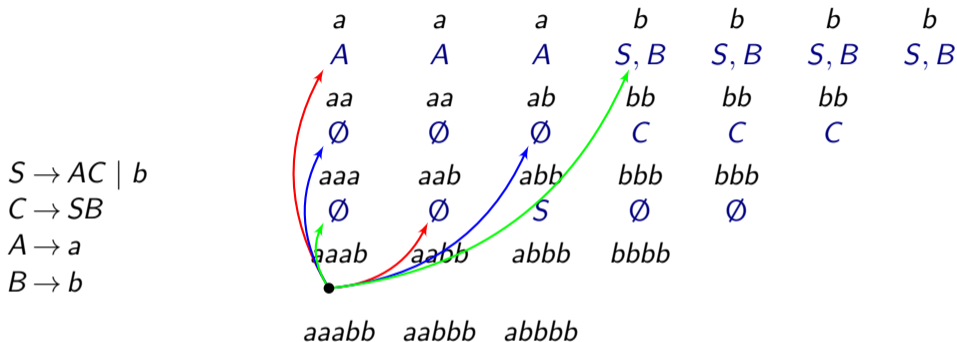
	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$	$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$							

*aaabb aabbb abbbb*

*aaabbb aabbbb*

*aaabbbb*

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)



aaabbb aabbbb

aaabbbb

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

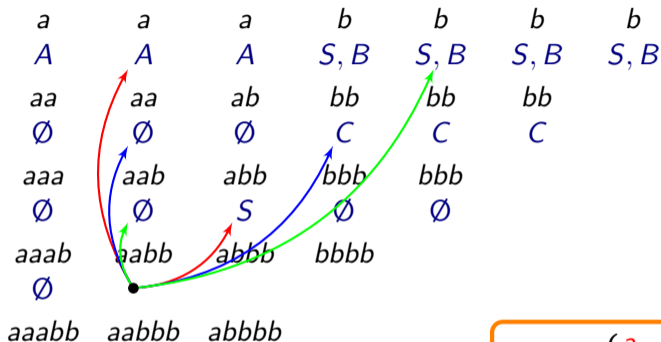
	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$	$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$	$\emptyset$						
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				

aaabbb aabbbb

aaabbbb

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$



aaabb    aabbb    abbbb

aaabbb    aabbbb

aaabbbb

$$aabb = \begin{cases} a \cdot abb \\ aa \cdot bb \\ aab \cdot b \end{cases}$$

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
	$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
	$\emptyset$	$\emptyset$					
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						

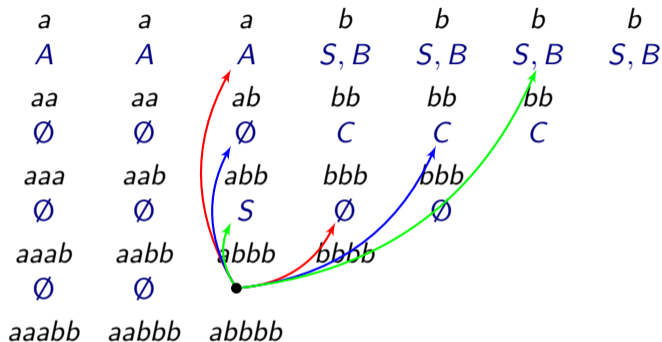
For no  $N \in V$   
does  $N \rightarrow AS$

*aabb*

$\{A\} \times \{S\} \cup$   
 $\emptyset \times \{C\} \cup$   
 $\emptyset \times \{S, B\} =$   
 $\{(A, S)\}$

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$



aaabbb aabbbb

aaabbbb

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$	$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$	$\emptyset$	$\emptyset$	<i>C</i>				
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				

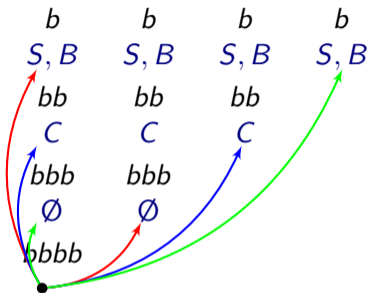
aaabbb aabbbb

aaabbbb

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$\emptyset$	$\emptyset$	<i>C</i>				
<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				



$$bbbb = \begin{cases} b \cdot bbb \\ bb \cdot bb \\ bbb \cdot b \end{cases}$$

aaabbb aabbbb

aaabbbb



Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

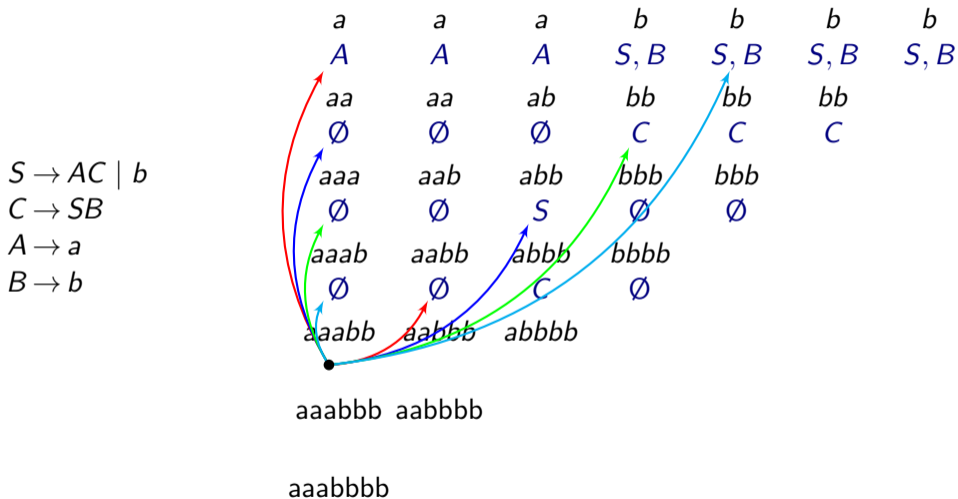
<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$\emptyset$	$\emptyset$	<i>C</i>	$\emptyset$			
<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				

aaabbb aabbbb

aaabbbb

$$bbbb = \begin{cases} b \cdot bbb \\ bb \cdot bb \\ bbb \cdot b \end{cases}$$

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

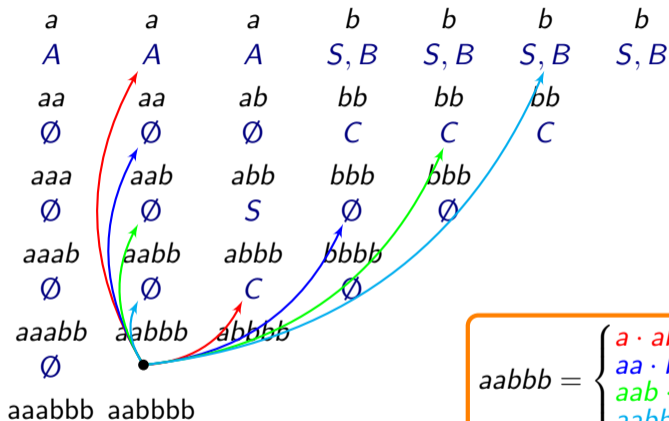


Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$	$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$	$\emptyset$	$\emptyset$	<i>C</i>	$\emptyset$			
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	$\emptyset$						
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$



$$aabbbb = \begin{cases} a \cdot abbb \\ aa \cdot bbb \\ aab \cdot bb \\ aabb \cdot b \end{cases}$$

aaabbbb

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

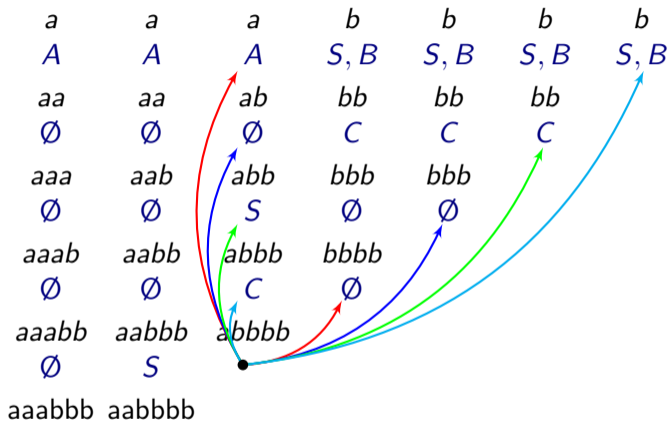
$a$	$a$	$a$	$b$	$b$	$b$	$b$
$A$	$A$	$A$	$S, B$	$S, B$	$S, B$	$S, B$
$aa$	$aa$	$ab$	$bb$	$bb$	$bb$	
$\emptyset$	$\emptyset$	$\emptyset$	$C$	$C$	$C$	
$aaa$	$aab$	$abb$	$bbb$	$bbb$		
$\emptyset$	$\emptyset$	$S$	$\emptyset$	$\emptyset$		
$aaab$	$aabb$	$abbb$	$bbbb$			
$\emptyset$	$\emptyset$	$C$	$\emptyset$			
$aaabb$	$aabbb$	$abbbb$				
$\emptyset$	$S$					
$aaabbb$	$aabbbb$					

$$aabbbb = \begin{cases} a \cdot abbb \\ aa \cdot bbb \\ aab \cdot bb \\ aabb \cdot b \end{cases}$$

aaabbbb

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$



aaabbbb

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$	$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$	$\emptyset$	$\emptyset$	<i>C</i>	$\emptyset$			
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	$\emptyset$	<i>S</i>	$\emptyset$				
	<i>aaabbb</i>	<i>aabbbb</i>					
	<i>aaabbbb</i>						





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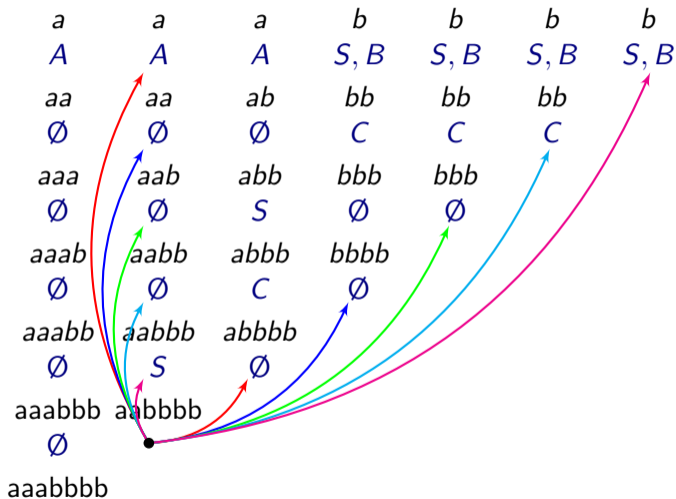
$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
	$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
	$\emptyset$	$\emptyset$	<i>C</i>	$\emptyset$			
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	$\emptyset$	<i>S</i>	$\emptyset$				
	<i>aaabbb</i>	<i>aabbbb</i>					
	$\emptyset$						
	<i>aaabbbb</i>						

$$aaabbbb = \begin{cases} a \cdot aabbb \\ aa \cdot abbb \\ aaa \cdot bbb \\ aaab \cdot bb \\ aabb \cdot b \end{cases}$$

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

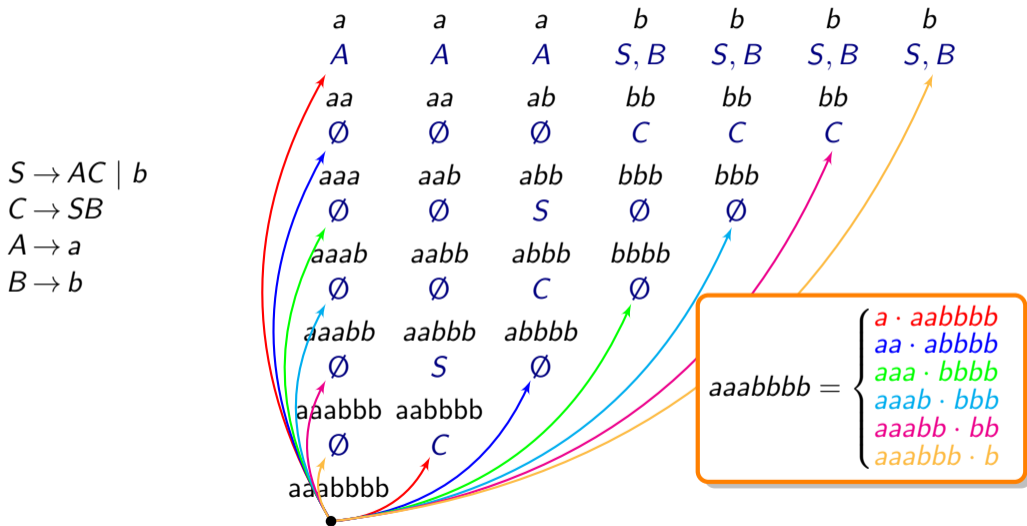
$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$



Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
$S \rightarrow AC \mid b$	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$C \rightarrow SB$	$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
$A \rightarrow a$	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$B \rightarrow b$	$\emptyset$	$\emptyset$	<i>C</i>	$\emptyset$			
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	$\emptyset$	<i>S</i>	$\emptyset$				
	<i>aaabbb</i>	<i>aabbbb</i>					
	$\emptyset$	<i>C</i>					
	<i>aaabbbb</i>						

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Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

	<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
	<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
	<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
	$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
	<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
	$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
	<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
	$\emptyset$	$\emptyset$	<i>C</i>	$\emptyset$			
	<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	$\emptyset$	<i>S</i>	$\emptyset$				
	<i>aaabbb</i>	<i>aabbbb</i>					
	$\emptyset$	<i>C</i>					
	<i>aaabbbb</i>						
	<i>S</i>						

$$aaabbbb = \begin{cases} a \cdot aabbbb \\ aa \cdot abbbb \\ aaa \cdot bbbb \\ aaab \cdot bbb \\ aabb \cdot bb \\ aabbb \cdot b \end{cases}$$

*a a a b b b b*

*A*

*aa aa ab bb bb bb*

*aaa aab abb bbb bbb*

$S \rightarrow AC \mid b$

$C \rightarrow SB$

*aaab aabb abbb bbbb*

$A \rightarrow a$

$B \rightarrow b$

*aaabb aabbb abbbb*

*aaabbb aabbbb*

*aaabbbb*

*a a a b b b b*

*A A*

*aa aa ab bb bb bb*

*aaa aab abb bbb bbb*

*aaab aabb abbb bbbb*

*aaabb aabbb abbbb*

*aaabbb aabbbb*

*aaabbbb*

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

*a a a b b b b*

*A A A*

*aa aa ab bb bb bb*

*aaa aab abb bbb bbb*

*aaab aabb abbb bbbb*

*aaabb aabbb abbbb*

*aaabbb aabbbb*

*aaabbbb*

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$



*a a a b b b b*

*A A A S, B*

*aa aa ab bb bb bb*

*aaa aab abb bbb bbb*

*aaab aabb abbb bbbb*

*aaabb aabbb abbbb*

*aaabbb aabbbb*

*aaabbbb*

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

*a a a b b b b*

*A A A S, B S, B*

*aa aa ab bb bb bb*

*aaa aab abb bbb bbb*

*aaab aabb abbb bbbb*

*aaabb aabbb abbbb*

*aaabbb aabbbb*

*aaabbbb*

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	

*aaa aab abb bbb bbb*

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

*aaab aabb abbb bbbb*

*aaabb aabbb abbbb*

*aaabbb aabbbb*

*aaabbbb*

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	

*aaa      aab      abb      bbb      bbb*

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

*aaab      aabb      abbb      bbbb*

*aaabb      aabbb      abbbb*

*aaabbb      aabbbb*

*aaabbbb*

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>

<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>
-----------	-----------	-----------	-----------	-----------	-----------

$\emptyset$

<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>
------------	------------	------------	------------	------------

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>
-------------	-------------	-------------	-------------

<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>
--------------	--------------	--------------

<i>aaabbb</i>	<i>aabbbb</i>
---------------	---------------

*aaabbbb*

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$					
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

*aaab*    *aabb*    *abbb*    *bbbb*

*aaabb*    *aabbb*    *abbbb*

*aaabbb*    *aabbbb*

*aaabbbb*

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$				
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

*aaab*    *aabb*    *abbb*    *bbbb*

*aaabb*    *aabbb*    *abbbb*

*aaabbb*    *aabbbb*

*aaabbbb*

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>			
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

*aaab*    *aabb*    *abbb*    *bbbb*

*aaabb*    *aabbb*    *abbbb*

*aaabbb*    *aabbbb*

*aaabbbb*



<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>		
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

*aaab*    *aabb*    *abbb*    *bbbb*

*aaabb*    *aabbb*    *abbbb*

*aaabbb*    *aabbbb*

*aaabbbb*

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		

$S \rightarrow AC \mid b$

$C \rightarrow SB$

$A \rightarrow a$

$B \rightarrow b$

*aaab*    *aabb*    *abbb*    *bbbb*

*aaabb*    *aabbb*    *abbbb*

*aaabbb*    *aabbbb*

*aaabbbb*

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$						
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>				
			<i>aaabbbb</i>			

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$	$\emptyset$					
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>				
			<i>aaabbbb</i>			

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$	$\emptyset$	<i>S</i>				
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
<i>aaabbb</i>	<i>aabbbb</i>					
<i>aaabbbb</i>						

*abb*  
 $\{A\} \times \{C\} \cup$   
 $\emptyset \times \{S, B\} =$   
 $\{(A, C)\}$

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$			
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
<i>aaabbb</i>	<i>aabbbb</i>					
<i>aaabbbb</i>						

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
<i>aaabbb</i>	<i>aabbbb</i>					
<i>aaabbbb</i>						

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$\emptyset$						
<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
	<i>aaabbb</i>	<i>aabbbb</i>				
			<i>aaabbbb</i>			



$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$\emptyset$	$\emptyset$					
<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
<i>aaabbb</i>	<i>aabbbb</i>					
<i>aaabbbb</i>						

*aabb*  
 $\{A\} \times \{S\} \cup$   
 $\emptyset \times \{C\} \cup$   
 $\emptyset \times \{S, B\} =$   
 $\{(A, S)\}$



$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$\emptyset$	$\emptyset$	<i>C</i>	$\emptyset$			
<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				

*aaabbb aabbbb*

*aaabbbb*

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$\emptyset$	$\emptyset$	<i>C</i>	$\emptyset$			
<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
$\emptyset$						
<i>aaabbb</i>	<i>aabbbb</i>					
<i>aaabbbb</i>						

*aaabb*

$\{A\} \times \emptyset \cup$   
 $\emptyset \times \{S\} \cup$   
 $\emptyset \times \{C\} \cup$   
 $\emptyset \times \{S, B\} =$   
 $\emptyset$

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$\emptyset$	$\emptyset$	<i>C</i>	$\emptyset$			
<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
$\emptyset$	<i>S</i>					
<i>aaabbb</i>	<i>aabbbb</i>					

*aaabbbb*

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>b</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>A</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>	<i>S, B</i>
<i>aa</i>	<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>bb</i>	<i>bb</i>	
$\emptyset$	$\emptyset$	$\emptyset$	<i>C</i>	<i>C</i>	<i>C</i>	
<i>aaa</i>	<i>aab</i>	<i>abb</i>	<i>bbb</i>	<i>bbb</i>		
$\emptyset$	$\emptyset$	<i>S</i>	$\emptyset$	$\emptyset$		
<i>aaab</i>	<i>aabb</i>	<i>abbb</i>	<i>bbbb</i>			
$\emptyset$	$\emptyset$	<i>C</i>	$\emptyset$			
<i>aaabb</i>	<i>aabbb</i>	<i>abbbb</i>				
$\emptyset$	<i>S</i>	$\emptyset$				
<i>aaabbb</i>	<i>aabbbb</i>					

*aaabbbb*

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

$a$	$a$	$a$	$b$	$b$	$b$	$b$
$A$	$A$	$A$	$S, B$	$S, B$	$S, B$	$S, B$
$aa$	$aa$	$ab$	$bb$	$bb$	$bb$	
$\emptyset$	$\emptyset$	$\emptyset$	$C$	$C$	$C$	
$aaa$	$aab$	$abb$	$bbb$	$bbb$		
$\emptyset$	$\emptyset$	$S$	$\emptyset$	$\emptyset$		
$aaab$	$aabb$	$abbb$	$bbbb$			
$\emptyset$	$\emptyset$	$C$	$\emptyset$			
$aaabb$	$aabbb$	$abbbb$				
$\emptyset$	$S$	$\emptyset$				
$aaabbb$	$aabbbb$					
$\emptyset$						
$aaabbbb$						

$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

$a$	$a$	$a$	$b$	$b$	$b$	$b$
$A$	$A$	$A$	$S, B$	$S, B$	$S, B$	$S, B$
$aa$	$aa$	$ab$	$bb$	$bb$	$bb$	
$\emptyset$	$\emptyset$	$\emptyset$	$C$	$C$	$C$	
$aaa$	$aab$	$abb$	$bbb$	$bbb$		
$\emptyset$	$\emptyset$	$S$	$\emptyset$	$\emptyset$		
$aaab$	$aabb$	$abbb$	$bbbb$			
$\emptyset$	$\emptyset$	$C$	$\emptyset$			
$aaabb$	$aabbb$	$abbbb$				
$\emptyset$	$S$	$\emptyset$				
$aaabbb$	$aabbbb$					
$\emptyset$	$C$					
$aaabbbb$						



$S \rightarrow AC \mid b$   
 $C \rightarrow SB$   
 $A \rightarrow a$   
 $B \rightarrow b$

$a$	$a$	$a$	$b$	$b$	$b$	$b$
$A$	$A$	$A$	$S, B$	$S, B$	$S, B$	$S, B$
$aa$	$aa$	$ab$	$bb$	$bb$	$bb$	
$\emptyset$	$\emptyset$	$\emptyset$	$C$	$C$	$C$	
$aaa$	$aab$	$abb$	$bbb$	$bbb$		
$\emptyset$	$\emptyset$	$S$	$\emptyset$	$\emptyset$		
$aaab$	$aabb$	$abbb$	$bbbb$			
$\emptyset$	$\emptyset$	$C$	$\emptyset$			
$aaabb$	$aabbb$	$abbbb$				
$\emptyset$	$S$	$\emptyset$				
$aaabbb$	$aabbbb$					
$\emptyset$	$C$					
$aaabbbb$						
$S$						

Linz 6th, §6.3, exercise 4, page 180 (solution page 421)

Conclusion

Since the start nonterminal  $S$  is in the cell labeled  $aaabbbb$ , then  $S \xRightarrow{*} aaabbbb$  and so  $aaabbbb \in \mathcal{L}(G)$ .

## Example 7.34 from HMU, 3rd, page 306

Consider the following CFG  $G$  in CNF.

$$S \rightarrow AB \mid BC$$

$$A \rightarrow BA \mid a$$

$$B \rightarrow CC \mid b$$

$$C \rightarrow AB \mid a$$

Let us test  $baaba$  for membership in  $\mathcal{L}(G)$ .

## Example (Continued)

$S \rightarrow AB \mid BC$

$A \rightarrow BA \mid a$

$B \rightarrow CC \mid b$

$C \rightarrow AB \mid a$

$b \quad a \quad a \quad b \quad a$

$B$

$ba \quad aa \quad ab \quad ba$

$baa \quad aab \quad aba$

$baab \quad aaba$

$baaba$

## Example (Continued)

$S \rightarrow AB \mid BC$   
 $A \rightarrow BA \mid a$   
 $B \rightarrow CC \mid b$   
 $C \rightarrow AB \mid a$

$b$        $a$        $a$        $b$        $a$   
 $B$        $A, C$   
 $ba$        $aa$        $ab$        $ba$   
 $baa$        $aab$        $aba$   
 $baab$        $aaba$   
 $baaba$

## Example (Continued)

$S \rightarrow AB \mid BC$   
 $A \rightarrow BA \mid a$   
 $B \rightarrow CC \mid b$   
 $C \rightarrow AB \mid a$

$b$	$a$	$a$	$b$	$a$
$B$	$A, C$	$A, C$		
$ba$	$aa$	$ab$	$ba$	
$baa$	$aab$	$aba$		
$baab$	$aaba$			
$baaba$				

## Example (Continued)

$S \rightarrow AB \mid BC$   
 $A \rightarrow BA \mid a$   
 $B \rightarrow CC \mid b$   
 $C \rightarrow AB \mid a$

$b$	$a$	$a$	$b$	$a$
$B$	$A, C$	$A, C$	$B$	
$ba$	$aa$	$ab$	$ba$	
$baa$	$aab$	$aba$		
$baab$	$aaba$			
$baaba$				

## Example (Continued)

$S \rightarrow AB \mid BC$   
 $A \rightarrow BA \mid a$   
 $B \rightarrow CC \mid b$   
 $C \rightarrow AB \mid a$

$b$	$a$	$a$	$b$	$a$
$B$	$A, C$	$A, C$	$B$	$A, C$
$ba$	$aa$	$ab$	$ba$	
$baa$	$aab$	$aba$		
$baab$	$aaba$			
$baaba$				



## Example (Continued)

$S \rightarrow AB \mid BC$   
 $A \rightarrow BA \mid a$   
 $B \rightarrow CC \mid b$   
 $C \rightarrow AB \mid a$

<i>b</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>a</i>
<i>B</i>	<i>A, C</i>	<i>A, C</i>	<i>B</i>	<i>A, C</i>
<i>ba</i>	<i>aa</i>	<i>ab</i>	<i>ba</i>	
<i>S, A</i>				
<i>baa</i>	<i>aab</i>	<i>aba</i>		
<i>baab</i>	<i>aaba</i>			
<i>baaba</i>				



## Example (Continued)

$S \rightarrow AB \mid BC$   
 $A \rightarrow BA \mid a$   
 $B \rightarrow CC \mid b$   
 $C \rightarrow AB \mid a$

$b$	$a$	$a$	$b$	$a$
$B$	$A, C$	$A, C$	$B$	$A, C$
$ba$	$aa$	$ab$	$ba$	
$S, A$	$B$	$S, C$		
$baa$	$aab$	$aba$		
$baab$	$aaba$			
$baaba$				

## Example (Continued)

$S \rightarrow AB \mid BC$   
 $A \rightarrow BA \mid a$   
 $B \rightarrow CC \mid b$   
 $C \rightarrow AB \mid a$

$b$	$a$	$a$	$b$	$a$
$B$	$A, C$	$A, C$	$B$	$A, C$
$ba$	$aa$	$ab$	$ba$	
$S, A$	$B$	$S, C$	$S, A$	
$baa$	$aab$	$aba$		
$baab$	$aaba$			
$baaba$				



## Example (Continued)

$S \rightarrow AB \mid BC$   
 $A \rightarrow BA \mid a$   
 $B \rightarrow CC \mid b$   
 $C \rightarrow AB \mid a$

<i>b</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>a</i>
<i>B</i>	<i>A, C</i>	<i>A, C</i>	<i>B</i>	<i>A, C</i>
<i>ba</i>	<i>aa</i>	<i>ab</i>	<i>ba</i>	
<i>S, A</i>	<i>B</i>	<i>S, C</i>	<i>S, A</i>	
<i>baa</i>	<i>aab</i>	<i>aba</i>		
$\emptyset$	<i>B</i>			
<i>baab</i>	<i>aaba</i>			
<i>baaba</i>				



## Example (Continued)

$S \rightarrow AB \mid BC$   
 $A \rightarrow BA \mid a$   
 $B \rightarrow CC \mid b$   
 $C \rightarrow AB \mid a$

$b$	$a$	$a$	$b$	$a$
$B$	$A, C$	$A, C$	$B$	$A, C$
$ba$	$aa$	$ab$	$ba$	
$S, A$	$B$	$S, C$	$S, A$	
$baa$	$aab$	$aba$		
$\emptyset$	$B$	$B$		
$baab$	$aaba$			
$\emptyset$				
$baaba$				



## Example (Continued)

$S \rightarrow AB \mid BC$   
 $A \rightarrow BA \mid a$   
 $B \rightarrow CC \mid b$   
 $C \rightarrow AB \mid a$

<i>b</i>	<i>a</i>	<i>a</i>	<i>b</i>	<i>a</i>
<i>B</i>	<i>A, C</i>	<i>A, C</i>	<i>B</i>	<i>A, C</i>
<i>ba</i>	<i>aa</i>	<i>ab</i>	<i>ba</i>	
<i>S, A</i>	<i>B</i>	<i>S, C</i>	<i>S, A</i>	
<i>baa</i>	<i>aab</i>	<i>aba</i>		
$\emptyset$	<i>B</i>	<i>B</i>		
<i>baab</i>	<i>aaba</i>			
$\emptyset$	<i>S, A, C</i>			
<i>baaba</i>				

## Example (Continued)

$S \rightarrow AB \mid BC$

$A \rightarrow BA \mid a$

$B \rightarrow CC \mid b$

$C \rightarrow AB \mid a$

$b$        $a$        $a$        $b$        $a$   
 $B$        $A, C$        $A, C$        $B$        $A, C$

$ba$        $aa$        $ab$        $ba$   
 $S, A$        $B$        $S, C$        $S, A$

$baa$        $aab$        $aba$   
 $\emptyset$        $B$        $B$

$baab$        $aaba$   
 $\emptyset$        $S, A, C$

$baaba$   
 $S, A, C$

## Example

Kozen, Lecture 27, page 192.  $aabbab \in L(G)$ ?

$$S \rightarrow AB \mid BA \mid SS \mid AC \mid BD$$

$$A \rightarrow a$$

$$B \rightarrow b$$

$$C \rightarrow SB$$

$$D \rightarrow SA$$

The grammar is in Chomsky Normal Form.

$S \rightarrow AA \mid AB \mid AC$

$S \rightarrow BA \mid BD$

$C \rightarrow SB$

$D \rightarrow SA$

$A \rightarrow a$

$B \rightarrow b$

$a \quad a \quad b \quad b \quad a \quad b$   
 $A$

$aa \quad ab \quad bb \quad ba \quad ab$

$aab \quad abb \quad bba \quad bab$

$aabb \quad abba \quad bbab$

$aabba \quad abbab$

$aabbab$

$S \rightarrow AA \mid AB \mid AC$   
 $S \rightarrow BA \mid BD$   
 $C \rightarrow SB$   
 $D \rightarrow SA$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>a</i>	<i>b</i>
<i>A</i>	<i>A</i>				
<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>ba</i>	<i>ab</i>	
<i>aab</i>	<i>abb</i>	<i>bba</i>	<i>bab</i>		
<i>aabb</i>	<i>abba</i>	<i>bbab</i>			
<i>aabba</i>	<i>abbab</i>				
<i>aabbab</i>					

*a a b b a b*

*A A B*

*aa ab bb ba ab*

*aab abb bba bab*

*aabb abba bbab*

*aabba abbab*

*aabbab*

$S \rightarrow AA \mid AB \mid AC$

$S \rightarrow BA \mid BD$

$C \rightarrow SB$

$D \rightarrow SA$

$A \rightarrow a$

$B \rightarrow b$

*a a b b a b*

*A A B B*

*aa ab bb ba ab*

*aab abb bba bab*

*aabb abba bbab*

*aabba abbab*

*aabbab*

$S \rightarrow AA \mid AB \mid AC$

$S \rightarrow BA \mid BD$

$C \rightarrow SB$

$D \rightarrow SA$

$A \rightarrow a$

$B \rightarrow b$

<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>a</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>A</i>	
<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>ba</i>	<i>ab</i>	

$S \rightarrow AA \mid AB \mid AC$

$S \rightarrow BA \mid BD$

$C \rightarrow SB$

$D \rightarrow SA$

$A \rightarrow a$

$B \rightarrow b$

*aab*    *abb*    *bba*    *bab*

*aabb*    *abba*    *bbab*

*aabba*    *abbab*

*aabbab*



<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>a</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>A</i>	<i>B</i>
<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>ba</i>	<i>ab</i>	

$S \rightarrow AA \mid AB \mid AC$

$S \rightarrow BA \mid BD$

$C \rightarrow SB$

$D \rightarrow SA$

$A \rightarrow a$

$B \rightarrow b$

*aab*    *abb*    *bba*    *bab*

*aabb*    *abba*    *bbab*

*aabba*    *abbab*

*aabbab*

$S \rightarrow AA \mid AB \mid AC$   
 $S \rightarrow BA \mid BD$   
 $C \rightarrow SB$   
 $D \rightarrow SA$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>a</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>A</i>	<i>B</i>
<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>ba</i>	<i>ab</i>	
$\emptyset$					
<i>aab</i>	<i>abb</i>	<i>bba</i>	<i>bab</i>		
<i>aabb</i>	<i>abba</i>	<i>bbab</i>			
<i>aabba</i>	<i>abbab</i>				
<i>aabbab</i>					

$S \rightarrow AA \mid AB \mid AC$

$S \rightarrow BA \mid BD$

$C \rightarrow SB$

$D \rightarrow SA$

$A \rightarrow a$

$B \rightarrow b$

$a \quad a \quad b \quad b \quad a \quad b$   
 $A \quad A \quad B \quad B \quad A \quad B$

$aa \quad ab \quad bb \quad ba \quad ab$   
 $\emptyset \quad S$

$aab \quad abb \quad bba \quad bab$

$aabb \quad abba \quad bbab$

$aabba \quad abbab$

$aabbab$

$S \rightarrow AA \mid AB \mid AC$   
 $S \rightarrow BA \mid BD$   
 $C \rightarrow SB$   
 $D \rightarrow SA$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>a</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>A</i>	<i>B</i>
<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>ba</i>	<i>ab</i>	
$\emptyset$	<i>S</i>	$\emptyset$			
<i>aab</i>	<i>abb</i>	<i>bba</i>	<i>bab</i>		
<i>aabb</i>	<i>abba</i>	<i>bbab</i>			
<i>aabba</i>	<i>abbab</i>				
<i>aabbab</i>					

$S \rightarrow AA \mid AB \mid AC$   
 $S \rightarrow BA \mid BD$   
 $C \rightarrow SB$   
 $D \rightarrow SA$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>a</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>A</i>	<i>B</i>
<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>ba</i>	<i>ab</i>	
$\emptyset$	<i>S</i>	$\emptyset$	<i>S</i>		
<i>aab</i>	<i>abb</i>	<i>bba</i>	<i>bab</i>		
<i>aabb</i>	<i>abba</i>	<i>bbab</i>			
<i>aabba</i>	<i>abbab</i>				
<i>aabbab</i>					

$S \rightarrow AA \mid AB \mid AC$   
 $S \rightarrow BA \mid BD$   
 $C \rightarrow SB$   
 $D \rightarrow SA$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>a</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>A</i>	<i>B</i>
<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>ba</i>	<i>ab</i>	
$\emptyset$	<i>S</i>	$\emptyset$	<i>S</i>	<i>S</i>	
<i>aab</i>	<i>abb</i>	<i>bba</i>	<i>bab</i>		
<i>aabb</i>	<i>abba</i>	<i>bbab</i>			
<i>aabba</i>	<i>abbab</i>				
<i>aabbab</i>					



$S \rightarrow AA \mid AB \mid AC$   
 $S \rightarrow BA \mid BD$   
 $C \rightarrow SB$   
 $D \rightarrow SA$   
 $A \rightarrow a$   
 $B \rightarrow b$

$a$	$a$	$b$	$b$	$a$	$b$
$A$	$A$	$B$	$B$	$A$	$B$
$aa$	$ab$	$bb$	$ba$	$ab$	
$\emptyset$	$S$	$\emptyset$	$S$	$S$	
$aab$	$abb$	$bba$	$bab$		
$\emptyset$	$C$				
$aabb$	$abba$	$bbab$			
$aabba$	$abbab$				
$aabbab$					



$S \rightarrow AA \mid AB \mid AC$   
 $S \rightarrow BA \mid BD$   
 $C \rightarrow SB$   
 $D \rightarrow SA$   
 $A \rightarrow a$   
 $B \rightarrow b$

$a$	$a$	$b$	$b$	$a$	$b$
$A$	$A$	$B$	$B$	$A$	$B$
$aa$	$ab$	$bb$	$ba$	$ab$	
$\emptyset$	$S$	$\emptyset$	$S$	$S$	
$aab$	$abb$	$bba$	$bab$		
$\emptyset$	$C$	$\emptyset$			
$aabb$	$abba$	$bbab$			
$aabba$	$abbab$				
$aabbab$					

$S \rightarrow AA \mid AB \mid AC$   
 $S \rightarrow BA \mid BD$   
 $C \rightarrow SB$   
 $D \rightarrow SA$   
 $A \rightarrow a$   
 $B \rightarrow b$

$a$	$a$	$b$	$b$	$a$	$b$
$A$	$A$	$B$	$B$	$A$	$B$
$aa$	$ab$	$bb$	$ba$	$ab$	
$\emptyset$	$S$	$\emptyset$	$S$	$S$	
$aab$	$abb$	$bba$	$bab$		
$\emptyset$	$C$	$\emptyset$	$C$		
$aabb$	$abba$	$bbab$			
$aabba$	$abbab$				
$aabbab$					

$S \rightarrow AA \mid AB \mid AC$   
 $S \rightarrow BA \mid BD$   
 $C \rightarrow SB$   
 $D \rightarrow SA$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>a</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>A</i>	<i>B</i>
<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>ba</i>	<i>ab</i>	
$\emptyset$	<i>S</i>	$\emptyset$	<i>S</i>	<i>S</i>	
<i>aab</i>	<i>abb</i>	<i>bba</i>	<i>bab</i>		
$\emptyset$	<i>C</i>	$\emptyset$	<i>C</i>		
<i>aabb</i>	<i>abba</i>	<i>bbab</i>			
<i>S</i>					
<i>aabba</i>	<i>abbab</i>				

*aabbab*

$S \rightarrow AA \mid AB \mid AC$   
 $S \rightarrow BA \mid BD$   
 $C \rightarrow SB$   
 $D \rightarrow SA$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>a</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>A</i>	<i>B</i>
<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>ba</i>	<i>ab</i>	
$\emptyset$	<i>S</i>	$\emptyset$	<i>S</i>	<i>S</i>	
<i>aab</i>	<i>abb</i>	<i>bba</i>	<i>bab</i>		
$\emptyset$	<i>C</i>	$\emptyset$	<i>C</i>		
<i>aabb</i>	<i>abba</i>	<i>bbab</i>			
<i>S</i>	<i>S</i>				
<i>aabba</i>	<i>abbab</i>				

*aabbab*

$S \rightarrow AA \mid AB \mid AC$   
 $S \rightarrow BA \mid BD$   
 $C \rightarrow SB$   
 $D \rightarrow SA$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>a</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>A</i>	<i>B</i>
<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>ba</i>	<i>ab</i>	
$\emptyset$	<i>S</i>	$\emptyset$	<i>S</i>	<i>S</i>	
<i>aab</i>	<i>abb</i>	<i>bba</i>	<i>bab</i>		
$\emptyset$	<i>C</i>	$\emptyset$	<i>C</i>		
<i>aabb</i>	<i>abba</i>	<i>bbab</i>			
<i>S</i>	<i>S</i>	$\emptyset$			
<i>aabba</i>	<i>abbab</i>				

*aabbab*

$S \rightarrow AA \mid AB \mid AC$   
 $S \rightarrow BA \mid BD$   
 $C \rightarrow SB$   
 $D \rightarrow SA$   
 $A \rightarrow a$   
 $B \rightarrow b$

$a$	$a$	$b$	$b$	$a$	$b$
$A$	$A$	$B$	$B$	$A$	$B$
$aa$	$ab$	$bb$	$ba$	$ab$	
$\emptyset$	$S$	$\emptyset$	$S$	$S$	
$aab$	$abb$	$bba$	$bab$		
$\emptyset$	$C$	$\emptyset$	$C$		
$aabb$	$abba$	$bbab$			
$S$	$S$	$\emptyset$			
$aabba$	$abbab$				
$D$					
$aabbab$					

$S \rightarrow AA \mid AB \mid AC$   
 $S \rightarrow BA \mid BD$   
 $C \rightarrow SB$   
 $D \rightarrow SA$   
 $A \rightarrow a$   
 $B \rightarrow b$

<i>a</i>	<i>a</i>	<i>b</i>	<i>b</i>	<i>a</i>	<i>b</i>
<i>A</i>	<i>A</i>	<i>B</i>	<i>B</i>	<i>A</i>	<i>B</i>
<i>aa</i>	<i>ab</i>	<i>bb</i>	<i>ba</i>	<i>ab</i>	
$\emptyset$	<i>S</i>	$\emptyset$	<i>S</i>	<i>S</i>	
<i>aab</i>	<i>abb</i>	<i>bba</i>	<i>bab</i>		
$\emptyset$	<i>C</i>	$\emptyset$	<i>C</i>		
<i>aabb</i>	<i>abba</i>	<i>bbab</i>			
<i>S</i>	<i>S</i>	$\emptyset$			
<i>aabba</i>	<i>abbab</i>				
<i>D</i>	<i>C</i>				
<i>aabbab</i>					

$S \rightarrow AA \mid AB \mid AC$   
 $S \rightarrow BA \mid BD$   
 $C \rightarrow SB$   
 $D \rightarrow SA$   
 $A \rightarrow a$   
 $B \rightarrow b$

$a$	$a$	$b$	$b$	$a$	$b$
$A$	$A$	$B$	$B$	$A$	$B$
$aa$	$ab$	$bb$	$ba$	$ab$	
$\emptyset$	$S$	$\emptyset$	$S$	$S$	
$aab$	$abb$	$bba$	$bab$		
$\emptyset$	$C$	$\emptyset$	$C$		
$aabb$	$abba$	$bbab$			
$S$	$S$	$\emptyset$			
$aabba$	$abbab$				
$D$	$C$				
$aabbab$					
$S$					