

# Formal Languages and Automata Theory

## Homework Set #4

William Shoaff

Summer 2018 (June 5, 2018)

Most languages are not regular. The pumping lemma can sometimes be used to prove a language  $L$  is not regular.

Theorem: Pumping Lemma for Regular Languages Let  $L$  be a regular language.

- There exists a natural number  $n$  (depending on  $L$ ) such that
- For all  $x \in L$  with  $|x| \geq n$
- There exists strings  $u, v, w$  such that
  1.  $x = uvw$ ,
  2.  $v \neq \epsilon$  ( $|v| \geq 1$ ),
  3.  $|uv| \leq n$ ,
  4. and for all  $i \geq 0$ ,  $uv^i w \in L$ .

The pumping lemma follows from the pigeonhole principle: If  $x$  is a sufficiently long string in  $L$ , then the path from  $q_0$  to a final state must contain a cycle that can be pumped (traversed) an arbitrary number of times.

To use the pumping lemma, assume, by way of contradiction, that a language  $L$  is regular and construct a string that cannot be pumped deriving a contradiction to  $L$  is regular.

Show the following languages are not regular.

1. (2 pts)  $L = \{ww : w \in \{a, b\}^*\}$  Hint: Picking the right string in the language that cannot be pumped is critical.
  - (a) Show that  $0^n 0^n \in L$  and can be pumped using  $u = \epsilon$ ,  $v = 00$ , and  $w = 0^{2n-2}$ .
  - (b) Show that  $(01)^n (01)^n \in L$  and can be pumped and can be pumped using  $u = \epsilon$ ,  $v = 0101$ , and  $w = (01)^{2n-2}$ .
  - (c) Show that  $(0)^n 10^n 1 \in L$  and cannot be pumped.
2. (2 pts) Let  $L = \{w : w = w^r\}$  be the set of palindromes over  $\{a, b\}$ .
3. (2 pts) Let  $L$  be the set of strings of digits and the symbols  $+$  and  $=$  that represent equations that are True. For example,  $1 + 1 = 2 \in L$ , while  $3 + 5 = 9 \notin L$ .

4. (2 pts) Consider the right-linear grammar

$$\begin{aligned} S &\rightarrow 0A \\ A &\rightarrow 10A|\epsilon \end{aligned}$$

Construct an NFA for the grammar and identify the regular expression defined by the grammar.

5. (2 pts) Now consider the left-linear grammar

$$S \rightarrow S10|0$$

(a) Reverse its productions to get

$$S \rightarrow 10S|0$$

and construct an NFA for this grammar.

(b) Reverse the edges of the NFA constructed in problem 5 and exchange initial and final states. What regular expression does this NFA accept?

Total Points: 10