## Computer Science Comprehensive Exam—Fall 2004 Compiler Construction

Instructions: Do not put your name on the exam, please answer all the questions directly on the exam itself. Answer all the questions. You have 90 minutes. Explain answers as fully as possible, give examples or define terms, if appropriate.

1. What is the relationship between the set of languages generated by ambiguous grammars and the set of languages recognized by $\mathrm{LR}(1)$ parsers? Circle the best response.
(a) ambiguous grammars are a subset of LR(1)
(b) $\mathrm{LR}(1)$ is a subset of ambiguous grammars
(c) Their intersection is non-empty
(d) Their intersection is empty
(e) They are the same
2. What is the relationship between the set of languages recognized by $\operatorname{LR}(1)$ parsers and the set of languages recognized by LL(1) parsers? Circle the best response.
(a) $\operatorname{LR}(1)$ is a subset of $\mathrm{LL}(1)$
(b) $\mathrm{LL}(1)$ is a subset of $\mathrm{LR}(1)$
(c) Their intersection is non-empty
(d) Their intersection is empty
(e) They are the same
3. What is the relationship between the set of languages recognized by $\operatorname{LR}(1)$ parsers and the set of languages recognized by LALR(1) parsers? Circle the best response.
(a) $\operatorname{LR}(1)$ is a subset of $\operatorname{LALR}(1)$
(b) $\operatorname{LALR}(1)$ is a subset of $\operatorname{LR}(1)$
(c) Their intersection is non-empty
(d) Their intersection is empty
(e) They are the same
4. Name one specific use for dataflow analysis in a compiler.
5. Consider the following augmented grammar over the alphabet $\{\mathbf{i d},()$,$\} .$

$$
\begin{array}{ll}
0 & S \rightarrow E \$ \\
1 & E \rightarrow T \$ \\
2 & E \rightarrow E+T \\
3 & T \rightarrow \mathbf{i d} \\
4 & T \rightarrow(E)
\end{array}
$$

(a) Compute nullable, FIRST, and FOLLOW for all nonterminals of the grammar.

|  | nullable FIRST FOLLOW |
| :--- | :--- |
| $S$ |  |
| $E$ |  |
| $T$ |  |

(b) Give the $\operatorname{LR}(0)$ automaton for the grammar. Give the $\operatorname{LR}(0)$ parse table. Is the grammar LR(0)?
6. Consider the algorithm to compute $\operatorname{CLOSE}[I]$ for the set of $\operatorname{LR}(1)$ items $I$ for some grammar. Suppose the grammar contains the production $X \rightarrow \gamma$ where $X$ is some non-terminal and $\gamma$ is some string of terminals and non-terminals. Answer the following questions assuming $y$ and $z$ are terminal symbols.
(a) If $A \rightarrow \alpha \bullet X, z$ is in $I$, which item or items (if any) would be added to CLOSE[I]?
(b) If $A \rightarrow \alpha \bullet X y, z$ is in $I$, which item or items (if any) would be added to CLOSE[I]?
(c) If $A \rightarrow \alpha \bullet X \beta, z$ is in $I$, which item or items (if any) would be added to CLOSE[I]?
7. Using Sethi-Ullman register allocation, generate code for the following expression:

$$
(a+b \times c+d) \times(e+a)
$$

First, draw a tree for the given expression and label each node with the number of registers it needs during evaluation. Then give the generated code, use instructions like $r_{1} \leftarrow M[a]$ (load a into register 1 ), $r_{2} \leftarrow r_{3} \times r_{4}, r_{1} \leftarrow r_{3}-r_{2}$, etc.

