Name:

## Fall 2017 Comprehensive Examination

## Analysis of Algorithms

1. (10 pts) Show that  $\sum_{0 \le i < j < n} 1 = \binom{n}{2}$ . Note, this sum can also be written as  $\sum_{i=0}^{n-2} \sum_{j=i+1}^{n-1} 1$ .

2. (10 pts) Let A be an  $n \times n$  matrix. Use summation notation to compute the big-O time complexity of the code:

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3. (10 pts) Assume n is a power of 2, that is,  $n = 2^p$  for some natural number p > 0. Solve the recurrence

$$T(n) = 3T(n/2) + n,$$
  $T(1) = 0$ 

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4. (10 pts) The recursive cutRod algorithm below takes an array of prices p[] and a rod length n. It determines how to cut the rod to maximize the total value of the pieces.

```
int cutRod(int p[], int n) {
    <u>if</u> (n == 0) { <u>return</u> 0; }
    q = -\infty;
    for (int k = 1; k <= n; k++) { q = max {q, p[k] + cutRod(p, n - k)};}
    <u>return</u> q
    }
```

(a) What recurrence equation describes the time complexity T(n) of cutRod?

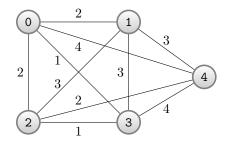
(b) Show that  $T(n) = 2^n$  satisfies the recurrence from part 4a.

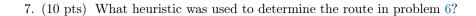
5. (10 pts) Here is a dynamic programming algorithm for the rod cutting problem. What is its time complexity?

```
int cutRod(int p[], n) {
1
2
      int r[n];
      r[0] = 0;
3
      for (int j=1; j <= n; j++) {
4
        q = -\infty;
\mathbf{5}
         for (int i=1; i <= j; i++) { q = \max \{q, p[i] + r[j-1]\}; \}
6
         r[j] = q;
\overline{7}
      }
8
9
      return r[n];
   }
10
```

6. (10 pts) Given a complete weighted graph, the traveling salesman problem (TSP) asks: What is the shortest possible route that visits each city exactly once and returns to the origin city?

Starting from city 0, construct a route for the graph below and determine its cost.





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8. (10 pts) If the heuristic was developed into an algorithm, explain what its time complexity would be. Would the algorithm always produce an optimal route?

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- 9. (20 pts) Answer True or False for the following sentences, or explain that the answer is unknown.(a) All NP-hard problems are NP-complete.
  - (b) The set of NP-complete problems is a subset of NP.
  - (c) NP-complete problems cannot have polynomial algorithms.
  - (d) In order to prove a problem X to be NP-complete one needs to develop a polynomial transformation from a known NP-complete problem to X.
  - (e) 2-SAT is an NP-hard problem.
  - (f) A tree with n nodes has n-1 edges.
  - (g) There exists a polynomial-time algorithm for finding the maximum spanning tree in a undirected and weighted graph.
  - (h) QuickSort has time complexity  $O(n \lg n)$  when executed on a sorted array of size n where the pivot is always the head of the array.
  - (i) There is a non-deterministic polynomial time algorithm that decides the halting problem.
  - (j) This sentence is False.

Thursday, November 9, 2017