

Normal Forms

Assume without loss of generality that every functional dependency in F^+ is of the form $X \Rightarrow A$, where A is a single attribute.

A relational scheme is said to be in **first normal form** (1NF) if and only if each of its domains contains only scalar values (i.e., no “repeating groups” or unbounded lists).

Example:

$R1 = (\underline{SS\#}, \text{Name}, \text{Hobbies})$ is not in 1NF, but $R2 = (\underline{SS\#}, \text{Name}, \underline{\text{Hobby}})$ is.

A functional dependency $X \Rightarrow A$ is a partial dependency if

- X is a proper subset of any candidate key, and
- A is not part of any candidate key.

A relational scheme is said to be in **second normal form** (2NF) if and only if it is in 1NF and contains no partial dependencies.

A functional dependency $X \Rightarrow A$ is a transitive dependency if

- X is not a superset of any candidate key, (i.e., its not a superkey),
- X is not a proper subset of any candidate key, and
- A is not part of any candidate key.

A relational scheme is said to be in **third normal form** (3NF) if and only if it is in 2NF and contains no transitive dependencies.

“Every non-key attribute depends on the (some) key, the whole key, and nothing but the key.”

A relational scheme is said to be in **Boyce/Codd normal form** (BCNF) if every attribute depends on the (some) key, the whole key, and nothing but the key.